

CP/NET Network Operating System Reference Manual

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Foreword

CP/NET®, a network operating system, enables microcomputers to access common resources via a network. CP/NET allows microcomputers to share and transfer disk files, to share printers and consoles, and to share programs and data bases. CP/NET consists of servers running MP/M II® and requesters running CP/M®. The servers are hosts that manage the shared resources that the network requesters can access.

The hardware environment for CP/NET must include two or more microcomputers that can communicate in some way.

One of the microcomputers must execute the MP/M II operating system to provide the CP/NET server facilities. The processor executing MP/M II must be an 8080, 8085, or Z80 CPU with a minimum of 32K bytes of memory, 1 to 16 consoles, 1 to 16 logical or physical disk drives each containing up to eight megabytes, a clock/timer interrupt, and a network interface.

The CP/NET requester microcomputers must have 8080, 8085, or Z80 CPUs with at least 16K bytes of memory, 0 to 16 logical or physical disk drives each containing up to eight megabytes, and a network interface. A console is not absolutely required although it is strongly recommended.

The CP/NET Network Operating System Reference Manual is intended for several different levels of CP/NET users. It contains all the information you need to use CP/M applications programs on a CP/NET requester, to write new application programs under CP/NET, and to customize CP/NET for a specific network.

Section 1, an overview of the CP/NET system, discusses CP/NET features, network topologies, and the principles behind CP/NET operation.

Section 2 contains all the information you need to use the network when executing CP/M application programs. You need no skill level beyond that required for normal CP/M operation.

Section 3 describes the CP/NET interprocessor message format and each of the Network Disk Operating System (NDOS) functions you can invoke from application programs. This section provides the information you need to access the network primitives. Section 3 also discusses the implications of performing CP/M operations on a resource controlled by the MP/M II operating system.

Section 4 provides information for the systems programmer. This section describes how to write a custom Slave Network I/O System (SNIOS) that performs the CP/NET requester network functions. The mechanics of implementing and debugging a custom SNIOS are also discussed. Programmers attempting to develop an SNIOS should be familiar with CP/M and experienced in writing a custom CP/M BIOS. This section also explains how to write a custom Network Interface Process (NETWRKIF) that performs the CP/NET server network functions.

Section 4 also discusses implementing and debugging the NETWRKIF module. You must have a high degree of competence and experience with MP/M II to develop a custom NETWRKIF. You must be familiar with the process and queue descriptor data structures and the MP/M II XDOS primitive functions. Experience with implementing an XIOS for MP/M II might also be necessary.

Appendixes to this manual contain several example network communications packages.

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Section 1 CP/NET Overview

By separating the logical operating system from the hardware environment and placing all hardware-independent code in a separate I/O module, CP/M and MP/M II have gained widespread industry acceptance. The CP/NET operating system uses this same design approach. CP/NET is network independent. The Slave Network I/O System (SNIOS) module contains all network-dependent code for the requester. The Network Interface Process (NETWRKIF) module contains all network-dependent code for the server. Logical messages passed to and from the SNIOS or NETWRKIF are transmitted over an arbitrary network between servers and requesters using an arbitrary network protocol.

CP/NET and CP/NOS can be combined in a composite network consisting of MP/M II servers, CP/M requesters, and diskless CP/NOS requesters.

CP/NET is a bridge between a microcomputer running MP/M II and a microcomputer running CP/M. The MP/M II server manages resources that are considered public to the network. The CP/NET requesters executing CP/M have access to the public resources of the server and to their own local private resources, which cannot be accessed from the network. This architecture permits the server's resources to be shared among the requesters, yet guarantees the security of the requester's resources.

The MP/M II server responds to the network asynchronously in real-time; the CP/M requesters perform sequential I/O and are usually not capable of monitoring a network interface in real-time. Figure 1-1 illustrates the relationship between CP/M, MP/M II, and CP/NET.

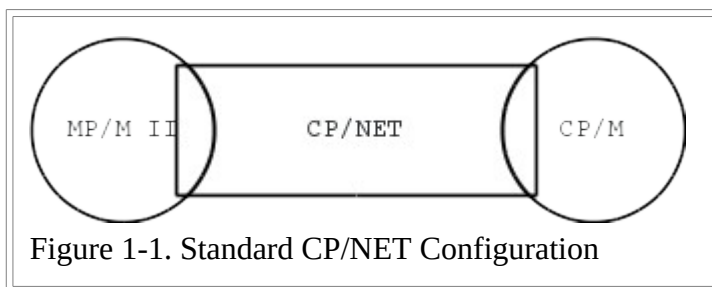
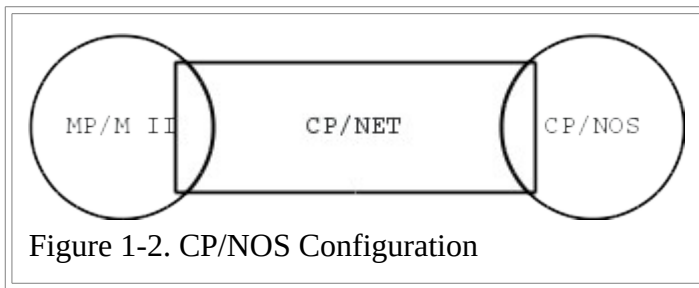


Figure 1-1. Standard CP/NET Configuration

CP/NOS, the second network operating system product, is designed for applications where the requester microcomputer lacks disk resources and is therefore unable to run CP/M. CP/NOS consists of

- a bootstrap loader that can be placed into ROM or PROM
- a skeletal CP/M containing only the console and printer functions
- the logical and physical portions of the CP/NET requester

At the user level, CP/NOS provides a virtual CP/M 2.X system to the requester microcomputer. A requester microcomputer can consist of no more than a processor, memory, and an interface to the network. Thus, a CRT with sufficient RAM can execute CP/M programs, performing its computing locally and depending on the network to provide all disk, printer, and other I/O facilities. Figure 1-2 illustrates the relationship between CP/NOS, MP/M II, and CP/NET.



1.1 CP/NET Features

CP/NET operates in multiple-processor environments ranging from tightly to loosely coupled to networked processors. In this manual, tightly coupled processors are those sharing at least a portion of common memory. Interprocessor messages communicate at memory speed. Loosely coupled processors do not have access to memory that is common or accessible by both processors; they communicate via a short, high-speed bus. Loosely coupled processors usually reside in the same physical box. Networked processors are usually physically separated and communicate over a serial link.

The CP/NET operating system is an upward-compatible version of CP/M 2.2, which provides system I/O facilities to requester microcomputers through a network. Additions to the Basic I/O System (BIOS) called the Slave Network I/O System (SNIOS), and a new Basic Disk Operating System (BDOS) called the Network Disk Operating System (NDOS), provide network access to System I/O facilities. The requester NDOS and NIOS are loaded and executed while running under CP/M 2.2.

In addition to the standard CP/M facilities, CP/NET provides the following capabilities:

- The network can be accessed for system I/O facilities.
- The network environment can be reconfigured to access I/O facilities according to application requirements.
- Messages can be transmitted and received between requesters and servers.
- An electronic mail system allows requesters and servers to send mail to each other.

The MP/M II server is implemented by adding some resident system processes at system generation (GENSYS) time. The resident system processes include server processes (SERVER) that perform the logical message-handling functions for the server and network interface processes (NETWRKIF) that you can customize for a particular hardware network interface.

1.2 CP/NET Configurations

CP/NET supports a number of different network topologies and a variety of system resources. The interprocessor message formats permit a requester to access more than one server for different resources.

Figure 1-3 illustrates an MP/M II system supporting a single CP/NET requester. The requester is a totally independent system, with its own console, printer, and disk resources. The requester can also access the MP/M II system's resources over the network. The MP/M II system also supports other users using local terminals.

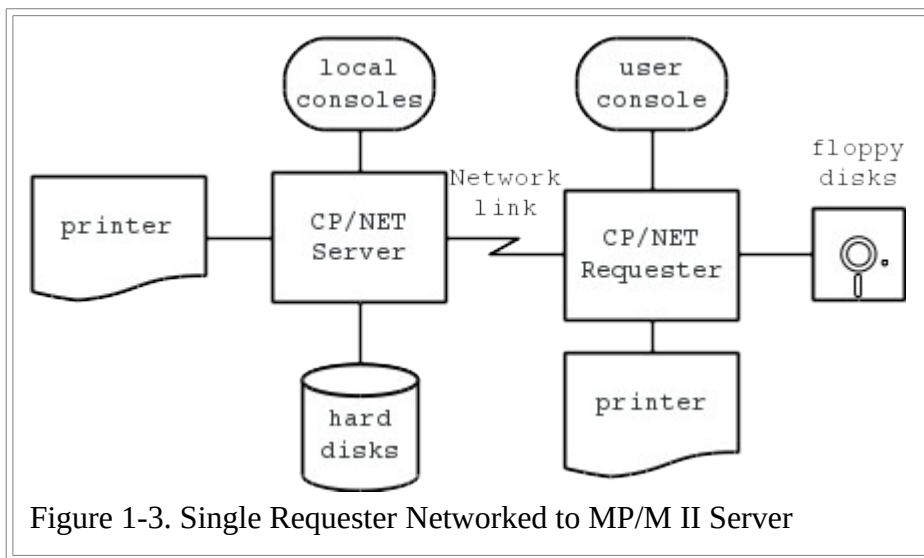


Figure 1-3. Single Requester Networked to MP/M II Server

Figure 1-4 shows an active hub-star network running CP/NET. Each requester is networked to the server through a unique network port. The requesters have their own local resources, but they also share the server's disk and printer resources. This topology is simple to implement because you can adapt the network protocol from the protocol used for RS-232 console drivers. The sample system in [Appendix E](#) uses this topology.

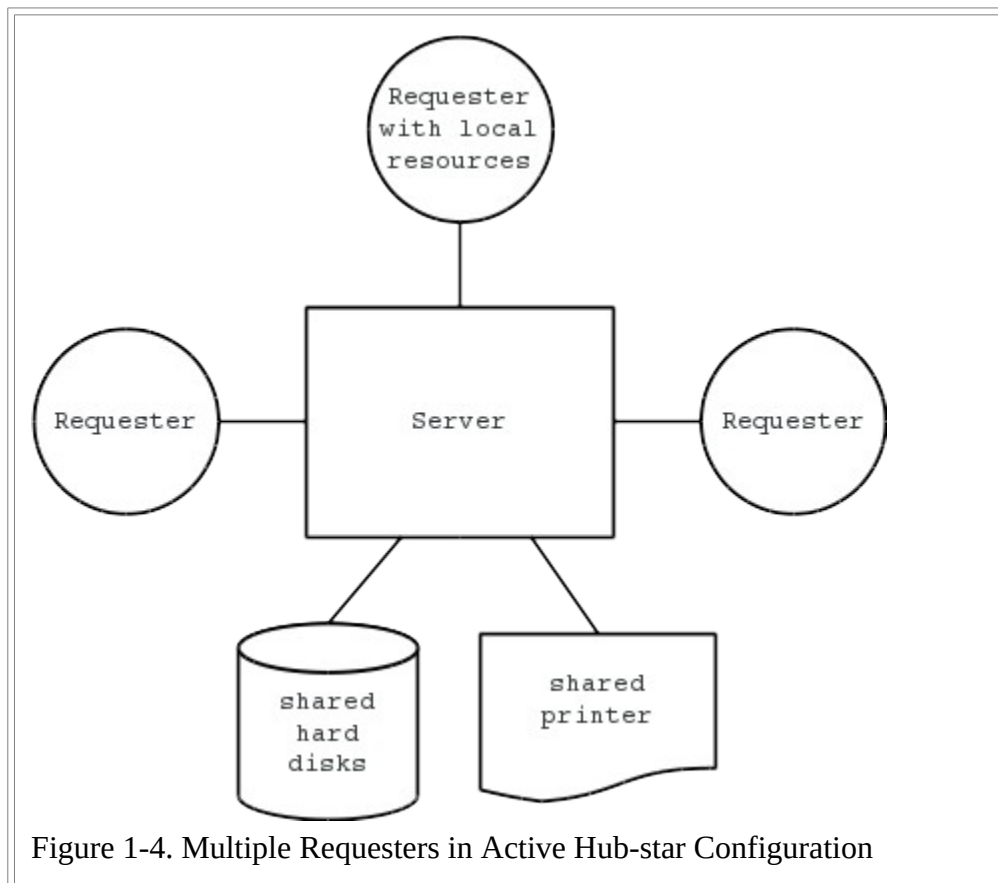
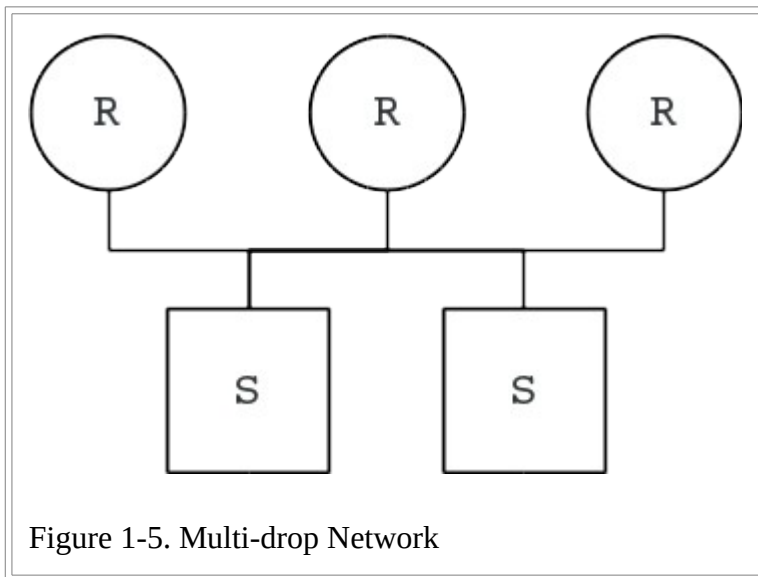
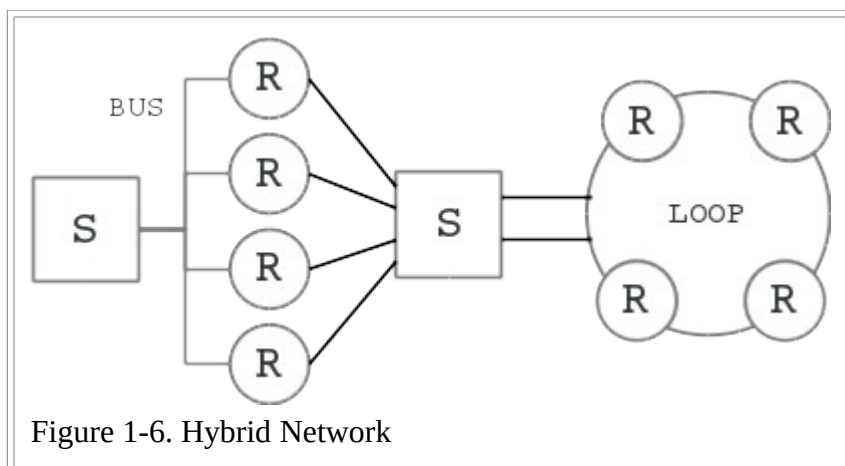


Figure 1-4. Multiple Requesters in Active Hub-star Configuration

Figure 1-5 shows a system of three requesters and two servers networked together in a bus or multi-drop configuration. The network protocol must be capable of resolving conflicts when nodes attempt to use the network simultaneously. Each requester has access to the resources of both servers, in addition to its own local resources. Appendixes [F](#) and [G](#) provide examples of CP/NET systems using this network topology.



Finally, you can combine these topologies, as well as other topologies like loops and trees, into a hybrid network topology. Figure 1-6 depicts such a topology, combining the bus, star, and loop forms.



1.3 How the Requester Works

The CP/NET requester software runs under an unmodified CP/M version 2 operating system. The requester operating system consists of three object modules: NDOS.SPR, SNIOS.SPR, and CCP.SPR. These modules are system page relocatable files that can be loaded directly under the CP/M BDOS and BIOS, regardless of their size or their location in memory.

The module NDOS.SPR contains the Network Disk Operating System (NDOS), the logical portion of the CP/NET system. The NDOS determines whether devices referenced by CP/M function calls are local to the requester or whether they are located on a remote system across a network. If a referenced device is networked, the NDOS, prepares messages to be sent across the network, controls their transmission, and finally reformats the result received from the network into a form usable by the calling application program. NDOS.SPR is distributed in object form by Digital Research. No modification to this module is required to run CP/NET.

The Slave Network I/O System (SNIOS) is contained in the module SNIOS.SPR. The systems implementer must customize this software to run on a particular computer and network system. The SNIOS performs

primitive operations that allow the NDOS to send and receive messages across a network. The SNIOS also provides a number of housekeeping and status functions to the NDOS. Digital Research distributes a number of example SNIOS modules in source form with CP/NET.

The final module, CCP.SPR, is a replacement for the normal CP/M CCP. Like the regular CCP, CCP.SPR is loaded directly below the operating system. However, CCP.SPR performs a number of special network functions that initialize the environment for a program.

The logical origin of SPR files is location zero. Each file has a 256-byte header, with locations 1 and 2 defined as the length of the code in the file. A bit map, appended to the end of the code, identifies bytes of the code that must be relocated when the code is loaded on a particular page (256-byte) boundary.

The CP/NET utility CPNETLDR relocates the bytes defined by the bit map. CPNETLDR loads SNIOS.SPR directly below the CP/M BDOS. NDOS.SPR is loaded directly below the SNIOS. CPNETLDR then passes control to an initialization routine. This routine modifies key areas of the operating system:

1. Location 5, which contains a jump to the BDOS entry point, is saved away by the NDOS.
2. Location 5 is then modified to jump to an entry point in the NDOS. This assures that the NDOS intercepts all CP/M function calls.
3. The BIOS jump vector entries for console status, console in, console out, list status, list out, and warm boot are replaced with entries that jump into special NDOS routines. The NDOS saves the BIOS entry points for these routines, allowing direct BIOS calls to these routines to be intercepted in exactly the same way that CP/M function calls are intercepted.

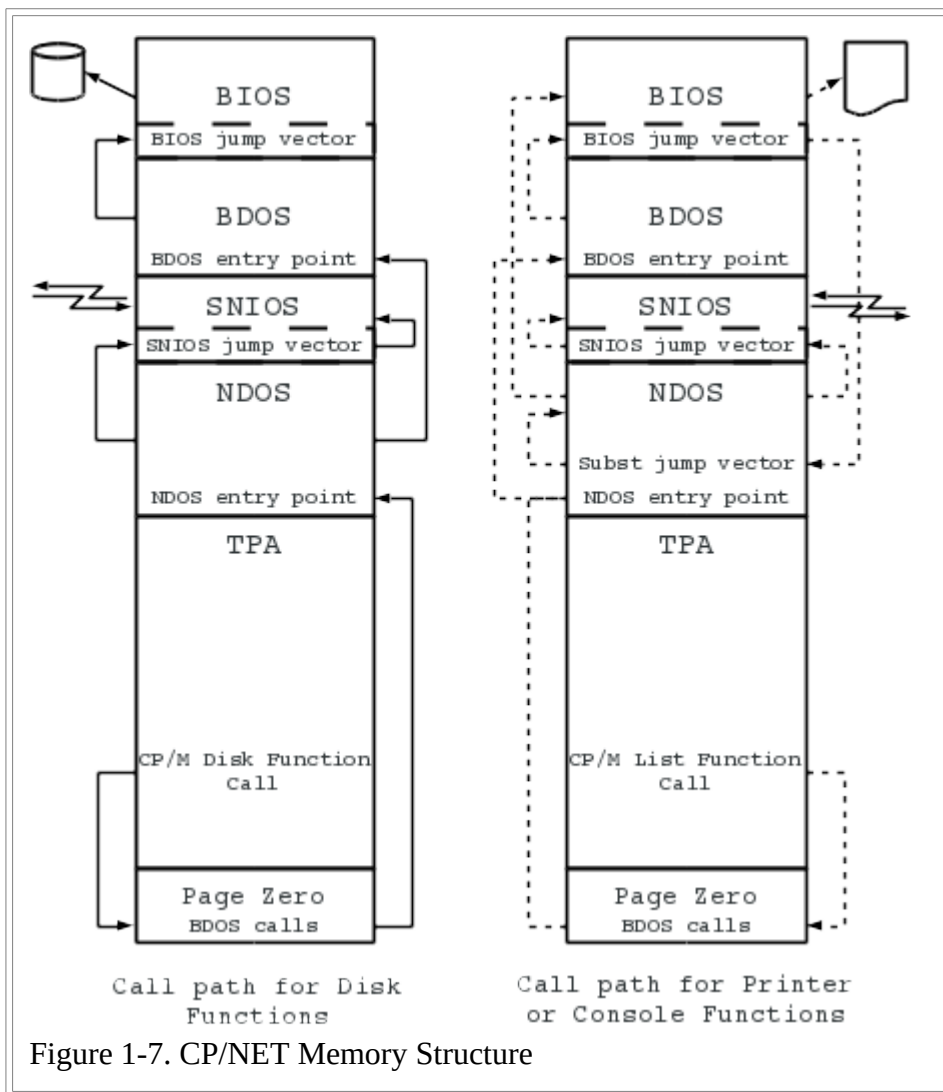
After these modifications have been made, the NDOS calls the SNIOS to initialize the network. The NDOS then jumps to its own warm boot routine, which performs a disk system reset, loads CCP.SPR, and then passes control to the CCP.

When an application program calls the CP/NET operating system via location 5, the NDOS is entered instead of the BDOS. Invalid functions return to the user program immediately as errors. Functions dealing with console or printer I/O immediately pass through to the local BDOS; but these functions are intercepted by the NDOS again when the BDOS calls the BIOS. At this level, the NDOS checks whether the console or printer is a networked device. If so, the NDOS sends a request across the network for the input or output.

Some functions have no meaning when they are sent across the network to a remote server. Examples of these are Function 26 (Set DMA Address), Function 32 (Get/Set User Number), and Function 12 (Return Version Number). The local BDOS always handles these functions. But the NDOS saves certain parameters from these functions for its own use, processing them before allowing them through to the BDOS.

Finally, the NDOS checks most functions that deal with either the disk drive system or the file system to determine whether they reference local devices. If so, these functions pass unmodified to the BDOS. The NDOS also checks whether these functions reference devices that exist somewhere out on the network. If they do, the NDOS constructs a network message to be sent to the system on which the device exists. The network message contains the network function to be performed and the information necessary to perform it.

Figure 1-7 illustrates how the CP/NET operating system is organized. The solid line outlines the function flow of an operation on a networked disk drive. The dotted line traces the flow of an I/O operation to a networked list device or console. Arrows indicate possible function flow.



When an NDOS requester sends a function message out over the network, a response from the addressed server is implied. As soon as the NDOS has successfully called the SNIOS to send the message, the NDOS calls the corresponding message receive routine, also in the SNIOS. This procedure precludes the problem of trying to recover sequencing information from an arbitrary stream of messages.

The NDOS uses the network response to update the application program that made the function call. The NDOS then returns to the application program. If the device referenced was local, then the requester's BDOS updates the application program.

1.4 How the Server Works

Unlike the requester, the server software that runs under MP/M II does not modify the actual operating system. Rather, the operating system is a set of cooperating processes under MP/M II.

In its most basic form, each requester to be attached to a server requires two processes, communicating through two queues. One process, resident in the NETWRKIF.RSP module, performs the physical message transport task. The systems implementer must modify this process to accommodate the network's node-to-node protocol. The process's protocol must be compatible with that of the requester's SNIOS.

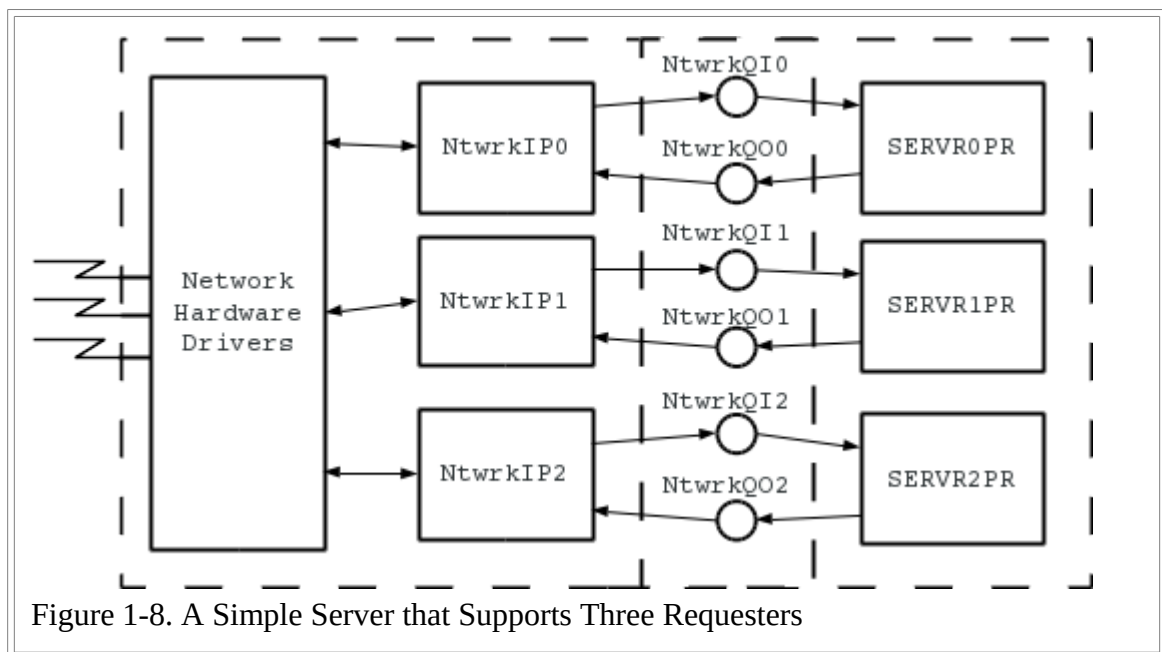
The NETWRKIF must be capable of monitoring one or more network lines in real-time and detecting when a requester is trying to send a message. The NETWRKIF must then receive the message, check it for data

integrity, and send it on to the logical portion of the server, contained in the module SERVER.RSP. When the SERVER module returns its response to the logical message, the NETWRKIF must receive the message and then transmit it across the network back to the requester.

The module SERVER.RSP performs the logical operation the requester specifies. After receiving the message from the NETWRKIF, SERVER.RSP checks to make sure that the requester is logged in properly. Then SERVER.RSP responds to the message by performing a series of MP/M II operating system calls. Using the information returned by those calls, the SERVER constructs a response message and sends it to the NETWRKIF module for transmission.

Both the NETWRKIF and SERVER modules are Resident System Process files (RSPs). RSPs are built into the MP/M II system during its GENSYS operation. When MP/M II is cold started, all RSPs are automatically dispatched. Each RSP module might contain multiple processes, but only one process per RSP is automatically dispatched. Because each requester bound to a server might require one process from the NETWRKIF and one from the SERVER, both RSPs contain initialization code to create additional copies of themselves. These processes can be reentrant. They can share the same code, but they have separate data areas to avoid conflict between program variables.

One of the simplest server architectures is shown in Figure 1-8. Processes from the NETWRKIF are named NtwrkIP<x> where <x> is the ASCII representation of a hexadecimal number between 0 and F. SERVER processes are named SERV<x>PR.



A NtwrkIP<x> process writes the address of an input message to a queue named NtwrkQI<x>. A SERV<x>PR process reads this queue while waiting for an input message. Because the queue is empty when the requester is not requesting service, the SERV<x>PR process is suspended and consumes no CPU resources.

When the NtwrkIP<x> process writes to the queue, the SERV<x>PR process is dispatched, and it begins to operate on the message. As soon as the NtwrkIP<x> process has finished sending the incoming message to NtwrkQI<x>, NtwrkIP<x> immediately tries to read a second queue, named NtwrkQO<x>. This queue is empty, and the NtwrkIP<x> process is consequently suspended until the SERV<x>PR process writes the response message to it. The NtwrkIP<x> can then transmit the message back to the requester.

Server functions can be divided into four categories:

- session control functions
- file serving functions
- print serving functions
- non-CP/NET functions

Session control functions permit a requester to log on to a server, log off, set compatibility attributes, set default passwords, and examine the server configuration table.

File serving functions make up the bulk of the server's work. These functions include opening and closing networked files, reading and writing files, and managing disk devices.

The server can operate as a print server in two different modes. If the MP/M module SPOOL.RSP is present in the system, requester outputs to a networked list device are spooled to a file for future printing. If no spooler exists in the system, the server manages the attaching and detaching of various print devices.

Finally, the NETWRKIF module can be designed to recognize a logical message that has no meaning to the SERVER module, but that can be operated on by a user-defined process. This feature allows you to use functions CP/NET does not provide.

Section 2

CP/NET User's Guide

This section describes the requester commands that enable you to access the network and use its resources. All the requester commands are actually COM files that reside on disk at the requester.

2.1 The LOGIN Command

The LOGIN command allows a requester to log in to a specified server. A requester must log in before any resources on the server can be accessed. Once a requester has logged in, it is not necessary to log in again even though the requester might power down and then power up again. A requester can only be logged off a server by an explicit LOGOFF command issued from the requester. The command takes the general form:

```
LOGIN {password}{[mstrID]}
```

where `password` is an optional 8 ASCII-character password; the default password is `PASSWORD`. `[mstrID]` is an optional two-digit server processor ID; the default is `[00]`. The simplest form is

```
A>LOGIN
```

2.2 The LOGOFF Command

The LOGOFF command allows a requester to log off from a specified server. Once a requester has logged off, the server cannot be accessed again until you issue a LOGIN command. The command takes the general form:

```
LOGOFF {[mstrID]}
```

where `[mstrID]` is an optional two-digit server processor ID; the default is `[00]`. The most simple form is

```
A>LOGOFF
```

2.3 The NETWORK Command

The NETWORK command enables a requester to assign selected I/O to the network. The NETWORK command updates the requester Configuration table. The command takes the general form:

```
NETWORK {local dev}{=}{server dev{[srvrID]}}
```

where `local dev`server dev is the specification of a server device such as A:, B: ... P: in the case of a disk device or 0, 1 ... 15 in the case of CON: or LST:. A missing `server dev` defaults to 0 in the case of CON: or LST:. `[srvrID]` is an optional two-digit hexadecimal server processor ID. The default is `[00]`. Typical assignments are

```
A>NETWORK LST:
```

```
A>NETWORK LST:=3[07]      (list dev #3 on server 07)
```

```
A>NETWORK CON:=2         (console #2 on dflt srvr)
```

```
A>NETWORK B:=D:[F]       (logical B: is D: on server 0F)
```

Note: when networking drive A: to a server, the file CCP.SPR must reside on the networked drive, or warm boot operations fail. Do not network a device to a nonexistent or off-line server because network errors could result.

2.4 The LOCAL Command

The LOCAL command enables a requester to reassign selected I/O back to local from the network. The LOCAL command updates the requester configuration table. The command takes the general form:

```
LOCAL {local dev}
```

where `local dev` is the specification of a local device such as LST:, A:,... CON:. The following are typical assignments:

```
A>LOCAL LST:  
A>LOCAL B:
```

2.5 The ENDLIST Command

The ENDLIST command sends a hexadecimal 0FF to the list device, signaling that a list output to a networked printer is finished. If a spooler is resident on the server, the spool file is closed and enqueued for printing. If no spool file is present, the networked list device is freed for use by another requester.

Note: the CCP implements an endlis every time a program terminates, provided that CTRL-P is not active at the time. Turning CTRL-P off also causes an endlis.

```
A>ENDLIST
```

2.6 The DSKRESET Command

The DSKRESET command functions exactly like the PRL that executes under MP/M II. DSKRESET resets the specified drive, so a disk can be changed. The command takes the general form:

```
DSKRESET {drive(s)}
```

where `drive` is a list of the drive names to be reset. If any of the drives specified cannot be reset, the console displays the message:

```
***Reset Failed***
```

The following are typical disk resets:

```
A>DSKRESET          (resets all drives)  
A>DSKRESET B:,F:   (reset drive B: and F:)
```

2.7 The CPNETLDR Command

The CPNETLDR command loads the requester CP/NET system. Specifically, the SNIOS.SPR file loads and relocates directly below the CP/M BDOS. The NDOS.SPR file loads and relocates directly below the SNIOS.

From that point on, the BIOS, BDOS, SNIOS, and NDOS remain resident in memory. The CPNETLDR requires no user customization. CPNETLDR displays an error message when loader errors are encountered. Listing 2-1 is a typical CPNETLDR execution.

```
A>CPNETLDR  
CP/NET 1.2 Loader  
-----
```

```

BIOS          F600H 0A00H
BDOS          E800H 0E00H
SNIOS   SPR  E500H 0300H
NDOS    SPR  DB00H 0A00H
TPA          0000H DB00H

```

```

CP/NET 1.2 loading complete.
<Warm Boot>
A>

```

Listing 2-1. A Typical CPNETLDR Execution

2.8 The CPNETSTS Command

The CPNETSTS command displays the requester configuration table. The requester configuration table indicates the status of each logical device that is either local or assigned to a specific server on the network. Listing 2-2 shows a typical CPNETSTS execution.

```

A>cpnetsts

CP/NET 1.2 Status

Requester processor ID = 34H
Network Status Byte = 10H
Disk device status:
  Drive A: = LOCAL
  Drive B: = LOCAL
  Drive C: = Drive A: on Network Server ID = 00H
  Drive D: = Drive B: on Network Server ID = 00H
  Drive E: = LOCAL
  Drive F: = LOCAL
  Drive G: = LOCAL
  Drive H: = LOCAL
  Drive I: = LOCAL
  Drive J: = LOCAL
  Drive K: = LOCAL
  Drive L: = LOCAL
  Drive M: = LOCAL
  Drive N: = LOCAL
  Drive O: = LOCAL
  Drive P: = LOCAL
Console Device = LOCAL
List Device = List #0 on Network Server ID 00H
A>

```

Listing 2-2. A Typical CPNETSTS Execution

2.9 CTRL-P

A CTRL-P causes console output to be echoed to the list device until the next CTRL-P. The messages

```
CTL-P ON
```

and

```
CTL-P OFF
```

are displayed at the console. When the requester list device has been networked, the local system uses the

server printer. The second CTRL-P causes a hexadecimal FF to be sent to the server, causing the server to close and print the spool file.

Note: when the requester uses the server printer with a CTRL-P active, the requester must issue a second CTRL-P to cause the server to close the spooled file and begin printing it. When the requester is using the server printer and has invoked it with a program such as PIP, the warm boot at program termination causes the required endlist character to be sent to the server to close and print the spooled file.

The program ENDLIST is not needed to terminate network list output in these situations.

2.10 The MAIL Utility

The MAIL utility allows you to send, receive, and manage electronic mail in a network environment. MAIL operates using file based function calls, so special processing by the server is not required. MAIL runs transparently on either server or requester, so only one program is required throughout the entire electronic mail system.

MAIL allows you to send messages to a single node, broadcast messages to all nodes currently logged in, or receive messages.

Messages are stored for your future examination on the temporary file drives of CP/NET servers. A user's mail file is named

xxMAIL . TEX

where xx corresponds to your node ID. For example, if requester #5C wants his mail, the MAIL program accesses files named 5CMAIL.TEX on the temporary file drives of all the servers that node 5C currently has logged in. Every server in the CP/NET system might have one of these files, so other nodes in the network that do not have direct access to all of node 5C's servers can still send messages indirectly to it.

Menu-driven operation allows you to run the program with a minimum of instruction. Messages are limited in size to 1.7K bytes. You can enter messages into the system directly from the keyboard or through a predated file. Options allow you to answer a message immediately while reading your mail and to delete unwanted entries.

2.10.1 Menus

Three basic menus can appear during a MAIL session:

- Main Menu
- Input Source Menu
- Receive Response Menu

The Main Menu determines the basic operation to be performed. The Input Source Menu specifies whether input comes from a file or whether you enter it directly. Finally, the Receive Response Menu determines the disposition of messages you receive.

Enter a menu selection by typing the number associated with the selection, followed by a carriage return. If you type an invalid character or no character at all, the menu system defaults to the last item on the menu. You simply press the carriage return for common operations.

Main Mail Menu

The main mail menu appears when you enter the mail program and when any of its options have completed execution. Main mail menu options are

- 1 - Broadcast
- 2 - Send Mail
- 3 - Receive Mail
- 4 - Exit Program

A simple carriage return or an invalid entry at this level return you to CP/M or MP/M II command level.

Input Source Menu

The input source menu allows you to specify how message input is entered into the system. The input source menu has only two options:

- 1 - File
- 2 - Console Input

Receive Response Menu

The receive response menu determines the disposition of messages once the user has examined them. The options are

- 1 - Stop Receiving Mail
- 2 - Answer Message
- 3 - Delete Message From Mail File
- 4 - Answer Message, Then Delete
- 5 - Re-Examine Last Message
- 6 - Get Next Message

2.10.2 Data Entry

In addition to the menus, MAIL prompts you for a variety of inputs. These inputs determine the destination of messages, input files, and subjects.

Destination ID Prompt

When using the send mail option, MAIL requires an explicit destination to deliver the message properly. The system prompts for the destination. The legal value is a 2-digit hexadecimal number, followed by a carriage return. This value corresponds to a CP/NET server or requester ID value.

If you enter a value that is not a legal hexadecimal number, the system displays an error message, and prompts you again. The system does not check, however, to determine whether a requester or server with this ID exists on the network.

Subject Prompt

With both the broadcast and send mail options, MAIL prompts for a subject header. This header is displayed as the title of the message and is also used for answering mail to the message that is sent.

When the system prompts for subject, you can enter a subject header from 0 to 80 bytes long, followed by a carriage return.

Input File Prompt

If a preedited file contains the text of a message, MAIL prompts for the filename. You can then enter a valid CP/M file specification. If the file specified does not exist, the system displays an OPEN ERROR, and the program aborts.

Console Input Prompt

If you choose to enter a message directly from the console, MAIL prompts for input. You can then simply type the message. Individual message lines can be up to 78 characters long. A message, whether input from the console or from a file, must be no longer than 1764 characters, about enough to fill a standard terminal display. Longer messages are truncated.

To terminate input, the user presses CTRL-Z, followed by a carriage return.

2.10.3 MAIL Options

This section explains how the CP/NET system gathers and receives mail and how you control the disposition of mail.

Broadcast

The broadcast option sends a message to every node that it can find logged in to the CP/NET system.

MAIL works differently when it is running on a server under MP/M II, from the way it works when it is running on a requester under CP/M or CP/NOS. If a requester is broadcasting, MAIL sends the specified message to every server on which it is logged in as well as to every other requester logged in to those servers. If a server is broadcasting, MAIL sends the message only to every requester logged in to that server. A server has no means of initiating transactions with other servers, although it can use its own local MP/M II system to file mail for its own requesters.

A message cannot be broadcast to the broadcasting node.

To send a message to a given server and its associated requesters, MAIL must reference that server's temporary file drive across the network. If a requester has not networked the temporary file drive of a server, no messages are sent to that server.

When the broadcast option is entered, MAIL prompts you for a subject and message. When the operation is completed, it returns to the main menu.

Send Mail

The send mail option sends a message to a specific node in the CP/NET system. The destination can be either a server or a requester. If the option is running on a requester, it first searches the network to see if the node specified is logged in. If the option finds the node is logged in, it sends the message. If the option does not find the node, it leaves the message on the first server located when MAIL searches the local configuration table. If a destination requester logs in later, its mail will be waiting for it. Mail files can accumulate that were erroneously sent to nonexistent requesters or to servers that the requester sending the message had not logged onto when it sent the message.

If the option is running on a server, mail is left on that server, whether the node it is being sent to is logged in or not.

Upon selecting the send mail option, MAIL prompts you for a destination ID, a subject, and for the message

itself. MAIL then attempts to send the message. If MAIL cannot find a server with a temporary file drive to accept the message, the error NO SERVER MAIL DRIVE NETWORKED is displayed, and the program aborts.

Receive Mail

The receive mail option permits you to examine messages left for you on all the servers on which you are currently logged in. After each message is displayed, you are presented with a number of message-handling options.

If you are running MAIL on the server, only the mail file on the server is accessed. However, if MAIL is being run on a requester, each server to which the requester is logged in is searched for messages.

Each message is preceded by a header that tells you what node the message came from and the subject of the message. The actual message is then displayed. As a message is being displayed, you can halt the display by pressing CTRL-S and resume display by pressing CTRL-Q. At the end of the message, bring up the receive response menu by pressing any key. You can then take one of the options listed in Table 2-1.

Option	Explanation
Stop receiving mail	MAIL stops searching for more entries or additional files and returns to the main menu.
Answer message	MAIL prompts you to type in a reply message. The reply message is sent back to the sender of the original message. The subject of the reply message is the characters "RE: ", followed by the original subject.
Delete message	MAIL flags the message in the file as deleted. At the end of each file, or if you decide to stop receiving mail, deleted messages are physically removed from the file.
Answer, then delete	This option answers the message message just displayed, then deletes the message.
Display next message	Messages continue to be displayed in this fashion, allowing the user to respond to each one, until no more can be found. The message "No More Messages" is then displayed, and the program returns to the main menu.

Table 2-1. Receive Mail Message-handling Options

Upon completion of any message-handling options, with the exception of the reexamine option, the next message is displayed.

2.10.4 Error Messages

In addition to the error messages already mentioned, CP/NET returns file system errors. These errors display

```
ERROR READING FILE
ERROR WRITING FILE
```

or
ERROR OPENING FILE

followed by a filename. After displaying such an error, MAIL aborts.

It is possible to get the ERROR OPENING FILE message by specifying a nonexistent input file for sending or broadcasting a message. Almost all other instances of the messages, however, indicate possibly serious trouble with the network, the server file system, or the mail-handling system.

Section 3

CP/NET Programmer's Guide

This section provides information for the applications programmer who wants to write programs to run under CP/NET or to evaluate the performance and correctness of programs written for CP/M or MP/M II under the CP/NET operating system.

MP/M II performs all operations on a networked device and makes file security checks that CP/M does not usually make. Because MP/M was designed to run unmodified CP/M applications, these checks seldom prevent the use of a CP/M application under CP/NET.

3.1 CP/NET Interprocessor Message Format

The simple message format that CP/NET uses for interprocessor communication includes packaging overhead and the message itself. The packaging overhead is a header consisting of a message format code, a CP/NET destination address, a CP/NET source address, a CP/M function code, and a message size. The actual CP/NET message follows the header.

3.1.1 Message Format Code

The message format code is a single byte that specifies the format of the message itself. Digital Research reserves message formats 0-127 for general interprocessor message format codes and future use. The general interprocessor format codes follow the message format shown below, but differ in length of the individual fields. (See [Appendix B](#).)

The odd-numbered format codes are for response messages sent back from servers to requesters. Thus, a CP/M disk read function sent from a requester to a server has a message format code of 0, and the return code sent back from the server to the requester has a message format code of 1.

Implement the general interprocessor message formats 0 and 1 as shown in [Appendix A](#) because these formats promote standardization among microcomputers from different vendors.

3.1.2 Message Destination Processor ID

The message destination processor ID field is one byte long. Destination IDs can be in the range 0-0FE hex. An ID of 0FF is illegal. Many CP/NET utilities use a server destination of 0 as a default. For this reason, assign the most commonly used network server a node ID of 0.

3.1.3 Message Source Processor ID

The message source processor ID field is usually one byte long. The node sending the message always fills this field with its own ID. Valid source IDs range from 0 to 0FE hex. An ID of 0FF is illegal.

3.1.4 CP/M Function Code

The CP/M function code field is one byte long. The size of the message data field depends on the CP/M function. Each CP/M function has a specific number of bytes to be sent to the server and a specific number of bytes to be returned to the requester. [Appendix C](#) provides the logical message specification for each of the CP/M functions. Some of the CP/M function codes have no equivalent network function.

3.1.5 Size

- MP/M II file security checking can cause certain CP/M applications to abort because these applications modify fields in the File Control Block that make the FCB invalid to MP/M II. Special compatibility modes have been added to CP/NET to allow these applications to run without modification.
- Temporary filenames, like \$\$\$SUB or FILENAME. \$\$\$, are modified under CP/NET. If more than one requester requires a temporary file with the same name, this modification prevents collisions between filenames that otherwise cause an application to abort. The modification is transparent to the application, but it can be confusing when trying to analyze aborted programs.
- A CP/NET requester presents a different version number to an application program when it calls Function 12 (Return Version Number). Under CP/M 2.2, this function returns a 002x value. Under CP/NET, it returns a 022x value. Application programs checking this version number might not function properly. They must be modified. Modifications to CP/NET, to present the same version number as CP/M, are now included as application notes in all releases of the CP/NET product.
- You can protect files on networked drives from unauthorized access by requiring a requester to specify a predefined password. You can also assign default passwords to all servers logged on to a particular requester.
- Certain files that exist only on user zero can be opened by any other user number if they are opened in the proper mode.
- The operating system must handle the printer differently under CP/NET from under CP/M because printer output is buffered into 128-byte packets. The operating system must have some way of deciding when an application program has finished using the printer. Also, several requesters might be competing for the same printer.
- The allocation vector for a networked drive is returned into the NDOS's default message buffer on a call to function 27 (Get Allocation Vector Address) and register pair HL is set to the address of the message buffer. Because of this, the allocation vector must be used or moved before the next network message is sent, or the vector is destroyed.

Differences between the CP/M 2.2 BDOS and MP/M II file systems are more fully described in the following sections.

3.2.1 MP/M II vs. CP/M File Systems

MP/M II is a real-time, multitasking operating system. To function properly, MP/M II requires a file system capable of sharing files among multiple processes and resolving access conflicts among those processes. In contrast, CP/M is a single-task operating system, so no such conflicts can arise.

One of MP/M II's key methods for maintaining file system integrity is the File Control Block checksum. The FCB checksum takes into account the process controlling the FCB, the physical blocks allocated to the file, whether the file is open in a mode that allows other processes to share it, and other factors. When file-related functions are submitted to MP/M II, the checksum is examined. If the checksum is found to be invalid, MP/M II returns an error to the calling process.

MP/M II also returns an error if

- a process attempts to open a file in a mode incompatible with the mode of a file already opened by another process
- a valid password is not supplied for the file
- a user tries to write to a file opened for Read-Only access
- a process exceeds certain predefined parameters for the operating system

Because a single process handles all CP/NET activity on a server all of these limitations apply to a CP/NET requester performing file operations on a remote device. These limitations, however, do not apply to a

requester accessing a local device. The systems implementer should take these factors into account when designing servers for a CP/NET system.

3.2.2 Error Handling Under CP/NET

Most CP/NET function calls result in specific values returned in the CPU registers. These values can be pointers to data objects, bit vectors specifying drive status, directory codes, or success or error conditions. Directory, success, and error codes are returned in register A. Pointers and bit vectors are returned in register HL. Register A is always equal to register L, and register B is equal to register H for all CP/NET return codes.

Error Handling for Local Devices

When a CP/NET requester performs a local file operation, the function parameters pass untouched to the CP/M BDOS. The BDOS checks those parameters for validity and calls the BIOS to perform physical I/O functions. Two types of errors can arise from these local operations.

The BDOS can detect certain logical problems with a file function and return a logical error. If it does, an error code is returned in register A, but the calling application program is allowed to continue.

A physical error is returned when the BIOS is unable to successfully perform a physical operation requested by the BDOS. When the BDOS is presented with a physical error, it prints the following message on the console:

```
BDOS Err on <x>:  
<error message>
```

where <x> is the drive referenced when the error occurred, and <error message> is one of the four following errors:

- Bad Sector
- Select
- File R/O
- R/O

After the physical error message is printed, the BDOS waits for the user to respond to the error with one of two actions. Pressing CTRL-C causes the BDOS to perform a warm boot, aborting the program. Pressing any other key causes the BDOS to ignore the physical error and continue as if it had not occurred.

For a more complete discussion of CP/M 2.x errors, see the CP/M Operating System Manual, published by Digital Research.

Error Handling for Network Devices

When an application references a networked device, the MP/M II server performs the actual file operation and returns a message defining whether the operation was successful or not. Unlike the local case, the requester has only indirect knowledge of any error status. Direct physical error indications are impossible to obtain because a requester has no contact with the MP/M II XIOS. Instead, if an error occurs, MP/M II returns a message indicating that an error occurred and the type of error it was.

When referencing a remote device, the two types of errors possible under CP/NET are logical errors and extended errors.

Like logical errors under local CP/M, logical network errors define nonfatal error conditions, such as reading past the end of a file or attempting to open a nonexistent file. Some serious error conditions are returned as logical errors for functions that expect to process their own errors. These functions are

- 20 Read Sequential
- 21 Write Sequential
- 33 Read Random
- 34 Write Random
- 40 Write Random with Zero Fill
- 42 Lock Record
- 43 Unlock Record

Errors for these functions are returned in the return code field of a CP/NET message. The NDOS formats this field into register A, so the condition code upon return to the application program looks exactly as it does under local CP/M.

Some of the following codes can be returned in register A for each of the preceding functions:

- 00 Function Successful
- 01 Reading Unwritten Data or No Directory Space Available
- 02 No Available Data Block (Disk Full)
- 03 Cannot Close Current Extent
- 04 Seek to Unwritten Extent
- 05 No Directory Space Available
- 06 Random Record Greater than 3FFFF
- 08 Record Locked by Another Process
- 09 Invalid FCB
- 0A FCB Checksum Error
- 0B File Verify Error
- 0C Record Lock Limit Exceeded
- 0D Invalid File ID
- 0E No Room in System Lock List

Extended errors indicate that a potentially fatal condition has occurred during the execution of an MP/M II function. The condition can be a physical error, similar to the physical errors that can occur under CP/M. Or the condition can be an error produced by the file system, indicating that the specified operation violates the integrity of the file system.

When an extended error occurs under MP/M II, the default mode of operation displays the extended error message on the console attached to the calling process, and the process aborts, MP/M II provides, however, for returning extended errors to the calling process without aborting that process. In this return error mode, register A is set to FF hexadecimal, and register H contains the extended error code.

The CP/NET server uses return error mode because if the server aborted, it could not communicate further with the requester it was servicing until MP/M II was restarted. When the server detects an extended error, it constructs a special CP/NET message. The message is two bytes long, with the first byte (the return code) set

to FF. The second byte is set to the extended error code.

When the requester detects one of these special messages, it checks the error mode set by the application program with Function 45 (Set BDOS Error Mode). There are three possible modes:

- Default Mode
- Return Error Mode
- Return and Display Error Mode

If the NDOS is in default mode, it prints the following error message:

```
NDOS Err <xx>, Func <yy>
```

where <xx> is the extended error code in hexadecimal, and <yy> is the function being performed when the error occurred, also in hexadecimal. The NDOS then performs a warm boot, aborting the program.

In return error mode, the NDOS does not display a message or abort the program. Instead, the NDOS sets register A to FF and register H to the extended error code; then it returns to the application program.

If an extended error is detected in return and display error mode, the NDOS displays the error message on the console. But the NDOS does not abort the program, setting the registers in the same manner as return error mode.

Function 45 (Set BDOS Error Mode) does not exist under CP/M. Because of this, most CP/M applications automatically run in default mode. If an extended error occurs, these applications abort.

The following extended error codes can be returned to the NDOS:

- 01 Bad Sector--Permanent Disk Error
- 02 Read-Only Disk
- 03 Read-Only File
- 04 Drive Select Error
- 05 File Open by Another Process in Locked Mode
- 06 Close Checksum Error
- 07 Password Error
- 08 File Already Exists
- 09 Illegal ? in an FCB
- 0A Open File Limit Exceeded
- 0B No Room in System Lock List
- 0C Requester not Logged on to Server or Function Not Implemented on Server
- FF Unspecified Physical Error

Extended error 0C hex is returned, not by MP/M II, but by the server itself. This error indicates that the server is unable to process an otherwise valid CP/NET message, either because the requester is not logged in to that server or because the function code contained in the message is invalid.

Extended error FF can result only from two special functions, Get Allocation Vector Address and Get Disk Parameter Address. Because these functions return a pointer in register pair HL, it is not possible to detect a regular extended error. Instead, these functions return an FFFF value in HL if a physical error occurs. The

NDOS ensures that the address returned for these functions (including Get Server Configuration) never return an address with FF in the low byte, so if they return with A (or L) = 0FFH then the caller should assume an error.

Not all CP/NET functions are capable of returning extended errors. However, extended error 0C can be returned on any function, even on MP/M II functions that normally have no extended error associated with them. If an extended error is returned for such a function, the NDOS ignores it. The following functions can result in the performance of a network access but cannot produce an extended error:

- 1 Console Input
- 2 Console Output
- 5 List Output
- 9 Print String
- 10 Read Console Buffer
- 24 Return Login Vector
- 28 Write Protect Disk
- 29 Get Read-Only Vector
- 37 Reset Drive
- 39 Free Drive
- 64 Login
- 66 Send Message on Network
- 67 Receive Message on Network
- 70 Set Compatibility Attributes
- 106 Set Default Password

Any other function can cause a program to abort if an MP/M II extended error occurs, if an unsupported function is passed to the server, or if the server is not logged in.

3.2.3 Temporary Filename Translation

Many common application programs use temporary files. The names of these files often have the form FILENAME.*** or ***.SUB. When multiple copies of these applications run on different requesters logged on to the same server, a number of these temporary files can have the same name, causing extended MP/M II errors that abort the application program.

To solve this problem, each requester's NDOS recognizes temporary filenames destined for networked drives and implicitly renames them, so the filename an application presents to the operating system is not the one the NDOS presents to the MP/M II file system.

Each occurrence of the string *** in the first three bytes of a filename, as well as any filetype of ***, forms a CP/NET message with a filename or filetype of \$<xx>, where <xx> is the ASCII representation of the requester ID byte. Because all requesters have a unique ID, this modification guarantees the uniqueness of temporary filenames.

This modification is transparent to the calling application program. When the NDOS modifies a filename in a CP/NET message, it converts the filename back to its original form before updating the application's FCB. The only possible change to the FCB is that interface attributes set in the high-order bits of the filename strings modified are reset. This change poses no problems if temporary files are truly temporary. Treat temporary files like Read-Write files with the DIR attribute; delete them before the application program

terminates.

Functions 17 (Search For First Directory Entry) and 18 (Search For Next Directory Entry) do not perform temporary filename translation when referencing a networked drive. If a user creates file with a temporary filename and then attempts to locate it within his directory, this can be confusing.

For example, suppose that a user working on requester 5A enters the command:

```
REN $$$.$$=$BLAH.TMP
```

Then the user enters a DIR command. The file previously renamed appears as

```
$5A.$5A
```

in the directory.

If a temporary file is referenced on a drive that is local to the CP/NET system, the filename passes unmodified to the BDOS. -No conversion is necessary, because there is no possibility of conflict.

3.2.4 Opening System Files on User 0

Under MP/M II, a requester running in a user number other than 0 can access certain networked files in user 0. If an MP/M II file has its t2' interface attribute set, the file is a system file. If a networked file is opened in locked or Read-Only mode from a nonzero user number, the following actions are taken:

- If the file exists in the same user number, MP/M II opens the file.
- If the file does not exist in the same user number, MP/M II searches user 0. If the file exists on user 0 and it is a system file, MP/M II opens it just as though the file existed under the other user number.
- If the file exists on user zero as a system file, but it is also a Read-Only file (interface attribute t1'), MP/M II automatically opens the file in Read-Only mode.

The user of a CP/NET requester can make convenient use of these options. Because the CCP.SPR always opens files in Read-Only mode, all COM files can be placed in user 0 and marked as system files, making them accessible to all user numbers.

Because this facility does not exist under CP/M 2.x, all COM files on local devices must exist within the user numbers from which they are to be executed.

3.2.5 Compatibility Attributes

Because of MP/M II's added file security, applications written under CP/M might not work properly under MP/M II. Two basic factors contribute to the incompatibility. The first is the FCB checksum computation that MP/M II performs on open FCBs. Certain CP/M applications modify their FCBs in a way that makes their checksums invalid. Second, MP/M II defaults to opening all files in locked mode, allowing only one process to have a file open at a time. Although files can be opened in an unlocked or shared mode, an application must explicitly specify that the file is to be opened unlocked. CP/M applications have no knowledge of this procedure.

To enable CP/M applications to run unmodified under MP/M II, a system of compatibility attributes has been added. This feature is supported under CP/NET. Using compatibility attributes, a user can selectively disable parts of the MP/M II file security mechanism.

When a requester's CCP opens a COM file for loading and subsequent execution, it examines the high-order bits of the first, second, third, and fourth bytes of the filename. These bits are referred to as interface attributes F1', F2', F3', and F4'. The CCP constructs a byte based on the interface attribute set. It then uses this byte as a parameter for Function 70 (Set Compatibility Attributes). Function 70 causes the NDOS to send a logical compatibility attribute message to every server of which it has knowledge.

Table 3-1 defines the interface attributes.

Attribute	Meaning
F1'	causes MP/M II to behave as though all files were opened in Read-Only mode, although write accesses are still permitted. F1' is functionally equivalent to opening a file in unlocked mode, except that record locking is not possible. Using this attribute, two programs can update the same record simultaneously, leaving the file in an indeterminate state.
F2'	causes all file close operations to convert to partial close operations. A partial close uses the current FCB to update the directory but permits the application program to continue using the file without reopening it.
F3'	disables FCB checksum verification during close operations. Files are closed successfully as long as MP/M II can tell the file was initially opened and still has an item on the system lock list. If the file was not opened, an error is still returned.
F4'	disables all FCB checksum verification. F4' implicitly sets attributes F2' and F3' as well. Use this attribute with extreme caution because it is possible to perform valid file operations using corrupt FCBs. Doing this could result in serious damage to the files on the disk drive being referenced.

Table 3-1. Interface Attributes

The CCP uses the interface attributes to construct a one-byte parameter for the set compatibility attributes call by setting the following bits:

- F1' bit 7
- F2' bit 6
- F3' bit 5
- F4' bits 4, 5, and 6

All other bits are set to zero.

The set compatibility attributes logical message causes the server to change its process descriptor if the user has enabled compatibility attributes during the MP/M II GENSYS operation. Otherwise, the message is ignored.

When an application program terminates, the CCP resets all compatibility attributes. This prevents a subsequent program from operating in an environment with insufficient file security.

It is advisable to enable the minimum number of compatibility attributes necessary to allow a program to run properly. Use the following guidelines for setting the attributes:

- If the program aborts with NDOS Error 05, FILE OPEN BY ANOTHER PROCESS, set F1'.
- If the program aborts with NDOS Error 06, CLOSE CHECKSUM ERROR, set F3'.
- if an error code is returned in register A on I/O operations under CP/NET, but no error is returned under CP/M, try setting F2' If the problem persists, try setting both F2' and F3'. if the problem still

persists, set user attribute F4'. Make sure there is no possibility of corrupting the file system before using attribute F4'.

You can use the SET utility under MP/M II to enter compatibility interface attributes into a .COM file's directory entry from an MP/M II console. For example,

```
SET <filespec> [F1=0N,F3=0N]
```

If you cannot use MP/M II, you can set the interface attributes under program control using Function 30 (Set File Attributes).

3.2.6 Password Protection Under CP/NET

The MP/M II file system limits file access by unprivileged users through password protection for individual files. There are three levels of password protection for files:

- All access is denied without the password.
- The file can be read without the password, but it cannot be written to.
- The file can be read and written to without the password, but not deleted.

Use the SET utility to assign passwords under MP/M II. The procedure for assigning passwords is described in the MP/M II Operating System User's Guide. CP/NET does not support the assignment of passwords across the network.

CP/NET does, however, allow an application program to send a Password across the network when a file is opened. This allows a user on a CP/NET requester the most basic form of password support: operation on networked files that have been previously password protected.

If a read-protected file is opened and no password is specified, an extended error is returned across the network, and the Calling application aborts. The same error is also returned when an application attempts to write to a write-protected file for which no password was provided when the file was opened. Finally, any attempt to delete, rename, or change the attributes of a delete protected file without providing a password results in an extended error.

CP/NET also supports Function 106 (Set Default Password). Function 106 provides a password against which all protected files are checked if no password is provided or if the password is incorrect. This function can relieve an application of the responsibility to parse passwords constantly into the first eight bytes of the current DMA buffer.

CCP.SPR does not support MP/M II's facility of supplying passwords when the user enters a command line. Because of this, do not password-protect COM files unless a default password utility is provided to the user.

Because CP/M 2.x does not support any kind of file protection, passwords are ignored when referencing files on drives local to a CP/NET requester.

3.2.7 Networked List and Console Devices Under CP/NET

In addition to the 16 disk devices, CP/NET allows the user to map the list and console devices across the network. A number of requesters can share a printer, or a console can be logically attached to a completely independent system running CP/NET or CP/NOS. Such a system needs only a network interface to support full CP/M capability.

Unlike most requester BDOS calls, whether a console or list device is local or networked is determined, not at the BDOS intercept level, but at the BIOS-intercept level. This feature enables application programs to make direct BIOS calls for console and printer I/O and to continue to run transparently across the network.

List device I/O is handled in the following manner: when the BIOS call is made to LISTOUT, the NDOS traps it. The NDOS examines the configuration table to determine whether the list device is local to the CP/NET system or networked. If the list device is local, the call is passed through to the BIOS unchanged.

If the list device is networked, however, the NDOS stores the character to be listed in a special buffer, located directly below the requester configuration table. When 128 characters are stored, the NDOS sends a List Output logical message to the server upon which the list device is mapped. This buffering process improves system performance because one-character messages that would congest the network communication interfaces need not be sent between each requester and server.

Under CP/M, there is no need to tell the list device when a listing is complete because only one application can list at a time, and that application has complete control of the device during that time. Under CP/NET, however, more than one requester can share a printer. So a mechanism must be included to notify the server that a listing is done and that the list device is available to other requesters.

A special provision must be included so a partially filled list buffer can be flushed to the server when a listing is finished, and so the server can release the list device. Endlist, a special character equal to FF hex, is intercepted by the NDOS as the signal to terminate a listing.

The endlist character can come from one of four sources:

1. The CCP.SPR sends an endlist character every time it is entered and detects that a list is in progress. This causes an endlist every time a program terminates.
2. An application can issue an endlist to terminate its own listing.
3. Every time a CTRL-P is toggled to off, the NDOS console input routine detects this and issues its own endlist.
4. You can use the ENDLIST utility to terminate the listing.

The server can handle listing in two different modes. If the module SPOOL.RSP is present in MP/M II, the server takes all list output messages and writes them to a dedicated spooler file. When the server detects an endlist, it inserts a CTRL-Z end-of-file character into the message, closes the spooler file, and directs the SPOOL process to begin printing the file on the appropriate list device

If a SPOOL process is not resident under MP/M II, the server, upon receiving an initial list out message, performs an explicit attach list function on the specified list device. This prevents other requesters from using the list device until the requester being serviced is finished listing. All other requesters are suspended or receive network errors if they try to use the same list device. When the server finally receives the endlist character, it issues a detach list function, freeing the list device for another process.

Both server modes have potential disadvantages. A printer that uses a CTRL-Z as an escape sequence for special printing functions cannot be used with the SPOOL.RSP. Using CTRL-Z causes the spooler to terminate a print job prematurely, assuming that an end-of-file was encountered. On the other hand, explicit attaching and detaching of list devices can cause a network error if a requester attempts to attach a list device that is already in use, has its server become suspended, and eventually times out.

Console I/O cannot be buffered and sent across the network in large blocks because it is not possible to determine when input critical to the operation of an application is needed. The NDOS must therefore send such I/O across the network one character at a time.

As with list output, the NDOS traps console-related BIOS calls. The NDOS determines whether the console is local or networked. If the console is local, no action is taken, and the local BIOS is entered. If the console is networked, a raw or unfiltered console I/O message is sent to the server. The server performs the I/O function and sends a response back to the requester.

If a networked console is used with CP/NET, the system behaves unreliably when the console is also being used as a regular MP/M II terminal because MP/M II allocates a Terminal Message Process (TMP) to each known user console. Both a server process and a TMP can be waiting for input from the same console. Because of this, typed characters can be echoed normally, doubly echoed, or not echoed at all. The actual processes might or might not receive every character.

A networked console user should also be aware that, because each character must be sent over the network, networked consoles drastically degrade the performance of the entire CP/NET system. Networked consoles are not recommended unless there is no way to support a local console, as in certain industrial process-control applications.

The CTRL-P facility of CP/M is partially handled by the NDOS. The NDOS must know when CTRL-P is active because it must send an endlist character when the facility terminates. If the CCP detects that CTRL-P is active, it will not send an endlist, even if a program terminates.

3.3 CP/NET Function Extensions to CP/M

Applications accessing networked drives use the MP/M II file system to perform file operations. Many of those operations have slightly different meanings than they do under CP/M. For example, by setting the high-order bits of an FCB filename, a file can be opened or made in locked mode, unlocked mode, or Read-Only mode. CP/NET also allows an application to place a password in the current DMA buffer for opening password-protected files. Similarly, a close operation can perform either a permanent close or a partial close.

The return codes and side-effects of MP/M II functions also differ. Error-handling differences are discussed in [Section 3.2.2](#). The open and make functions also differ. These functions return a two-byte value, called the file ID, in the random record field of the opened FCB. The file ID is necessary for performing record locking functions.

For a complete description of how individual CP/M functions work under MP/M II, see the MP/M II Operating System Programmer's Guide.

This section describes CP/NET functions that have no counterpart under CP/M. These include MP/M II functions that do not exist under CP/M, as well as a set of dedicated CP/NET functions. All of these functions adhere to exactly the same calling conventions as the rest of CP/M and all follow the same conventions regarding return codes.

FUNCTION 38: ACCESS DRIVE		
Prevents Drives from Being Reset		
	Register	Value
Entry Parameters	C	26H
	DE	Drive Vector
Return Values	A	Return Code
	H	Extended Error

The Access Drive function inserts a dummy open file item in the system lock list for each drive specified in the drive vector. The drive vector is a 16-bit vector in which each possible drive is presented. Bit 0 represents drive A:, bit 1, drive B:, continuing through 15 for drive P:.

The NDOS separates the drive vector into a number of drive vectors, one per server that the NDOS can find in the requester's configuration table. The NDOS then sends a logical message to each of these servers. If any of these messages result in an extended error, the function aborts.

If a server's system lock list does not have enough room to fit all the dummy items for all the drives specified, or if the open file limit for the server process is exceeded, none of the items is inserted and Function 38 returns an extended error.

Because the NDOS sends messages to each server in sequence, an extended error on one server does not indicate that servers accessed previously failed to insert open file items. This differs from MP/M II, where only one file system controls the entire lock list. Note that drives might have to be freed after a failure resulting from an access drive call.

If the NDOS is in return error mode, an error condition on function 38 causes register A to be set to 0FFH, and register H contains one of the following codes:

- 0A Open File Limit Exceeded
- 0B No Room in the System Lock List
- 0C Server Not Logged In

Because Function 38 is meaningless to local drives under CP/NET, no call to the local BDOS is made.

FUNCTION 39: FREE DRIVE		
Free Specified Disk Drives		
	Register	Value
Entry Parameters	C	27H
	DE	Drive Vector

The Free Drive function purges servers' lock lists of all items pertaining to the drives specified. The drive vector is a 16-bit vector in which each possible drive is represented. Bit 0 represents drive A:, bit 1, drive B:, continuing through 15 for drive P:.

Because dummy drive accesses, locked records, and open files are all purged, close all important files before issuing the free drive call. Otherwise, a checksum error is returned on the next file access, and data might be lost.

The CP/NET CCP issues a free drive every time a program terminates. This prevents the server process associated with the requester from becoming clogged with useless files.

Because Free Drive is meaningless under CP/M, the operating system ignores entries in the drive vector that specify drives local to the requester.

Free Drive has no error return.

FUNCTION 42: LOCK RECORD		
Lock Records in a File		
	Register	Value
Entry Parameters	C	2AH
	DE	FCB Address
Return Values	A	Return Code
	H	Extended Error

The Lock Record function grants a requester exclusive write access to a specific record of a file opened in unlocked mode. Using this function, any number of requester processes can simultaneously update a common file.

To lock a record, a requester application must place the logical record number to be locked in the random record field of the file's FCB. The file ID number, a two-byte value that is returned in the random record field when a file is opened in unlocked mode, must be placed in the first two bytes of the current DMA buffer. When the lock function is called, a pointer to the FCB must exist in register pair DE.

The record to be locked must reside within a block currently allocated for the file. The lock fails if the record is locked by another process or requester. This prevents two processes from simultaneously updating the same record and leaving it in an indeterminate state.

If a file was opened in locked mode, the Lock Record function always returns successfully, but no explicit action is taken because the whole file is locked in the first place.

To use the Lock Record function, follow these steps:

1. Open the file in unlocked mode. Save the file ID returned in the random record field of the open FCB.
2. When the application needs to update the record, lock the record, even before attempting to read it. Reading a record that is locked by another process can result in leaving the record in an indeterminate state. If an error results because the record is locked by another process, repeat this step until the record is locked successfully. Place a timeout value on retrying the lock in case another requester has locked the record and then gone off line.
3. Read the record.
4. Update the record.
5. Write the record back.
6. Unlock the record.

The Lock Record function returns a 0 in register A if successful. Otherwise, the Lock Record function returns one of the following error codes in register A:

- 01 Reading Unwritten Data
- 03 Cannot Close Current Extent to Access Extent Specified
- 04 Seek to an Unwritten Extent
- 06 Random Record Number Greater than 3FFFF
- 08 Record Locked by Another Process
- 0A FCB Checksum Error
- 0B Unlock File Verification Error

0C Process Record Lock Limit Exceeded
 0D Invalid File ID in the DMA Buffer
 0E No Room on the System Lock List
 FF Extended Error

These extended errors can occur:

01 Permanent Error
 04 Select Error
 0C Requester Not Logged In to Server

The Lock Record function has no meaning when a drive local to the requester is referenced. The function returns with register A set to 0.

FUNCTION 43: UNLOCK RECORD		
Unlock Records in a File		
	Register	Value
Entry Parameters	C	2BH
	DE	FCB Address
Return Values	A	Return Code
	H	Extended Error

The Unlock Record function releases a previously locked record, allowing it to be locked and written to by another requester. The record to be unlocked must be placed in the random record field of the file's FCB. The file ID is a two-byte value that is returned in the random field when a file is opened in unlocked mode. The file ID must be placed in the first two bytes of the current DMA buffer. Register pair DE must contain a pointer to the FCB.

The Unlock Record function returns successfully if

- the file was opened in locked mode.
- the record specified is already unlocked.
- the record is locked by another process.

In all these cases, no action is performed.

Do not unlock a record until the requester's application program has finished updating the locked record and has written it back out to the file. Otherwise, another process might inadvertently destroy the updated information.

The Unlock Record function returns a 0 in register A if Successful. Otherwise, the function returns one of the following error codes in register A:

01 Reading Unwritten Data
 03 Cannot Close Current Extent to Access Extent Specified

- 04 Seek to an Unwritten Extent
- 06 Random Record Number Greater than 3FFFF
- 0A FCB Checksum Error
- 0B Unlock File Verification Error
- 0D Invalid File ID in the DMA Buffer
- FF Extended Error

These extended errors can occur:

- 01 Permanent Error
- 04 Select Error
- 0C Server Not Logged In

The Unlock Record function is meaningless when it references a requester's local drive; it returns a 0 in register A.

FUNCTION 45: SET BDOS ERROR MODE		
Defines CP/NET Error Handling		
	Register	Value
Entry Parameters	C	2DH
	E	Error Mode

The Set BDOS Error Mode function provides the NDOS with these options:

- aborting on extended errors
- returning the extended error to the calling application for handling
- returning the error to the application and displaying it on the console

All requester application programs are initially loaded in a default environment that causes the NDOS to abort on extended errors and to display the extended error code. Use Function 45 to change this default mode, according to the contents of register E.

Register	Explanation
0FFH	Return Error Mode. BDOS returns extended errors coming from the network to the application program. Register A is set to 0FFH, and register H contains the extended error code. No error message is displayed on the console.
0FEH	Return and Display Mode. BDOS returns the extended error in the same manner as in Return Error Mode, but also displays an extended error message.
Any Other Value	Default Mode.

Table 3-2. BDOS Error Modes

Function 45 is not implemented across the network. The NDOS maintains its own internal error mode flag and acts upon returning network messages according to that flag.

The Set BDOS Error Mode function has no effect on physical errors returned by the requester's local BIOS. These errors always display an error message, then they give the user the option of aborting the application program or continuing.

FUNCTION 64: LOGIN		
Initiate Session Between a Requester and a Server		
	Register	Value
Entry Parameters	C	40H
	DE	Ptr to Login Msg
Return Values	A	Return Code

The Login function identifies a requester to a server and initiates a session with that server. The Login function must always be successfully called before a requester can access a server's resources. Register pair DE must contain a pointer to a data structure that contains the following two fields:

- 00-00 Server ID byte
- 01-08 Password

The NDOS uses this structure to construct a logical LOGIN message to the server specified. Only the LOGIN message can be passed to the SERVER module without generating an extended error 0C, requester not logged in.

The server checks to see whether the password matches the password defined in the server configuration table. The server then scans the configuration table to find out whether logging in another requester exceeds the number of servers present in the system. If a server exists for the requester, and the password matches, the NDOS returns a 0 in register A. Otherwise, an error is flagged by returning an 0FFH in register A. The NDOS also returns a 0 in register A if the requester is already logged in.

FUNCTION 65: LOGOFF		
Terminate a Session Between a Requester and a Server		
	Register	Value
Entry Parameters	C	41H
	E	Server ID
Return Values	A	Return Code
	H	Extended Error

The Logoff function completes a session and breaks the logical binding between the server specified in register E and the calling requester. Once a Logoff has been performed, the server process is free to begin a session with another requester, if the the server's NETWRKIF can support the dynamic binding of requester

nodes to server processes.

Function 65 returns a 0 if successful. It returns an extended error 0C, requester not logged on to server, if unsuccessful.

FUNCTION 66: SEND MESSAGE ON NETWORK		
Send a Message to Another Network Node		
	Register	Value
Entry Parameters	C	42H
	DE	Pointer to Message
Return Values	A	Return Code

The Send Message on Network function sends messages across the network that might have no defined function on the MP/M II server. This allows applications to be written under CP/NET that use non CP/NET messages. Point-to-point communications packages, special electronic mail systems, implementation of requester synchronization functions, and special print spooling systems are examples of such applications.

To use Function 66, the address of the message to be sent must be passed in register pair DE. The message pointed to might have the standard CP/NET structure of FMT, DID, SID, FNC, SIZ, and MSG, or it might take some nonstandard format. In the latter case, the SNIOS must be able to recognize the nonstandard message and send it properly.

Unlike the usual CP/NET session protocol, the Send Message on Network function does not automatically attempt to receive a response to the message that was sent. So an application can send throw-away messages that do not require a logical acknowledgment or response. You can also define message types that can be broadcast to every node in the network.

If an application requires a logical response to a message sent using Function 66, make an explicit call to Function 67 (Receive Message on Network).

As a rule, set the FMT field of the message header of any nonstandard message sent through a CP/NET system to a value other than those reserved for use by Digital Research. Future releases can then run applications using Function 66, with minimal modification.

Function 66 returns an FF in registers A, H, and L if a network error occurred and the message was not sent.

FUNCTION 67: RECEIVE MESSAGE ON NETWORK		
Receive Message from Another Network Node		
	Register	Value
Entry Parameters	C	43H
	DE	Receive Buffer Address
Return Values	A	Return Code

The Receive Message on Network function is the counterpart of Function 66, Send Message on Network. Invoke it immediately after performing a send message if a logical response is expected. Function 67 can also

be used to wait for an unsolicited message from another node.

To use Function 67, an application must pass a pointer to a buffer area into which the message can be received in register DE. Upon return, registers A, H, and L are set to 0FFH if the function failed to receive the message properly.

Like Function 66, Function 67 can handle nonstandard messages across a CP/NET network, provided that the requester's SNIOS is equipped to handle them. For a more detailed discussion on how to use Functions 66 and 67, see section 3.4.

FUNCTION 68: GET NETWORK STATUS		
Get Network Status Byte from the Configuration Table		
	Register	Value
Entry Parameters	C	44H
Return Values	A	Network Status Byte

The Get Network Status function returns the configuration table's network status byte in register A. It also resets any error conditions in the status byte.

For a description of the fields contained in the network status byte, see [Section 4.2.1](#).

FUNCTION 69: GET CONFIGURATION TABLE ADDRESS		
Get Configuration Table Address		
	Register	Value
Entry Parameters	C	45H
Return Values	HL	Table Address

The Get Configuration Table Address function returns the address of the requester configuration table maintained in the SNIOS. Using this function, an application can dynamically modify the mappings of devices across the network. The utilities NETWORK and LOCAL use Function 69 to accomplish this kind of modification

For a description of the fields in the configuration table, see [Section 4.2.2](#).

FUNCTION 70: SET COMPATIBILITY ATTRIBUTES		
Configure Server File Systems for an Application		
	Register	Value
Entry Parameters	C	46H
	E	Compatibility Attribute Byte

The Set Compatibility Attributes function selectively disables the file security mechanism on all MP/M II

servers to which the calling requester has networked drives. This allows certain applications that run under CP/M but not under the MP/M II file system to run under CP/NET and access networked devices.

The CCP.SPR checks the compatibility interface attributes of all COM files that it loads for execution and performs a Set Compatibility Attributes function based on the pattern it finds. This is the only time to use this function. Applications should not modify their compatibility mode in mid-execution. Doing so might produce unpredictable results.

The compatibility attribute byte is set according to the interface attributes found in the COM file's name. The following attributes cause the corresponding bits to be set in register E prior to the call to Function 70:

- F1' bit 7
- F2' bit 6
- F3' bit 5
- F4' bits 4, 5, and 6

For a complete description of how to use compatibility attributes, see [Section 3.2.5](#).

Function 70 has no error return. Extended error messages from servers to which the requester is not logged in are ignored.

FUNCTION 71: GET SERVER CONFIGURATION TABLE ADDRESS		
Get Information About a Server		
	Register	Value
Entry Parameters	C	47H
	E	Server ID
Return Value	HL	Server Configuration Table Address

The Get Server Configuration Table Address function returns a pointer to parts of the specified server's configuration table. The ID of the server to be examined is passed in register E prior to calling Function 71, and a pointer to the received information is returned in register pair HL.

The data structure addressed by HL has the following format:

- 00-00 Server Temporary File Drive
- 01-01 Server Network Status Byte
- 02-02 Server ID
- 03-03 Maximum Number of Requesters Permitted on the Server
- 04-04 Number of Requesters Currently Logged In Bit Vector of Requesters Logged In in the Requester
- 05-06 ID Table
- 07-16 Requester ID Table

The information is identical with that contained in the server configuration table, except that the login password has been removed, and a byte containing the server's temporary file drive has added to the front of

the table.

Function 71 can determine whether other requesters are logged into a server. The temporary file drive can be used when an application wants to leave a file on a server but does not know the capacity or type of the server's disk drives. The MAIL utility makes frequent use of Function 71.

The server configuration table is returned across the network in a Special buffer in the NDOS. If more than one call is to be made to Function 71, and the calls reference a different server each time, the buffer is overwritten by each successive call. If an application must examine more than one server configuration table at once the table must be copied down into a buffer defined by the application.

If Function 71 passes a server ID to which the calling user is not logged on, an extended error 0C, requester not logged in, is returned.

FUNCTION 106: SET DEFAULT PASSWORD		
Establish a Default Password for File Access		
	Register	Value
Entry Parameters	C	46H
	DE	Password Address

The Set Default Password function allows an application to specify a password that is checked if an incorrect password is presented during an Open File function. If a file is password protected, MP/M II first checks for a password in the current DMA buffer. If no match is found, MP/M II then checks the default password set by Function 106. If MP/M II finds a match, it allows the requested operation to succeed. Otherwise, MP/M II returns an error.

When Function 106 is performed on a requester, the requester's NDOS attempts to set the default password on every server to which a drive is networked by that requester. Since Function 106 has no error return, extended requester not logged in errors are ignored

Each server process uses an MP/M II default password slot, starting with console 0 and using as many slots as there are requesters supported.

The default password set by Function 106 persists until another default password is set.

3.4 CP/NET Applications

In addition to running standard CP/M applications packages on a CP/NET requester, you can implement special applications using the network functions available in CP/NET. The applications can handle message processing in a distributed environment. Examples include high-performance print spoolers, node-to-node transfer utilities, and network management tools.

Using Functions 66 (Send Message on Network) and 67 (Receive Message on Network) , you can define an entire set of specialized messages to provide network functions. These messages must be recognized and processed by the SNIOS and NETWRKIF, but once implemented, they can be used by application programs as though they were functions themselves.

Suppose a specific network application requires a print spooler that provides special formatting features. You can write an application program that creates messages with a special code in the format byte of the CP/NET

message header. When the application wants to spool data to the special spooler on the server, it uses Function 66 to send the data.

On the server side, the NETWRKIF must be capable of recognizing the specially defined format code. When the NETWRKIF sees this format, instead of routing the message to a server process, it writes the message to a special queue. The actual spooler can reside as a process under MP/M II. The spooler reads the queue and spools the data.

Notice that Functions 66 and 67 are independent of the logical protocol of CP/NET, where every message sent by a requester implies that the requester waits to receive the message. This independence permits an application using a feature like a special spooler to return immediately after sending its message. The application need not wait for a logical acknowledgment.

Another convenient application is a file copy program that works without server intervention. Under the regular CP/NET protocol, the only way to copy a file on a local requester drive to the local drive of another requester is first to copy the file to a common networked drive, then copy it back to the other requester's drive. This is inefficient.

Instead, suppose that the users of the two requesters agree to cooperate in the copying of the file. They can do this by sending each other mail. One user invokes an application program called RECEIVE, while the other brings up an application program called SEND.

The SEND program merely reads the file into memory, then sequentially sends it to the other requester, using Function 66. The SEND program might or might not request verification from the receiving requester via Function 67. In the meantime, the RECEIVE program reads the messages from the network. No server intervention is required; only the two SNIOS modules of the requester are involved in the transmission. Even though the two requesters are only capable of sequential processing, they are still able to send and receive messages synchronously. This application does not require modifications to the SNIOS and NETWRKIF; the standard CP/NET protocol is sufficient, because such applications never reference the server.

Finally, a complex network might require automatic system monitoring and maintenance utilities. Using special message formats, you can design a set of messages that check which drives are usable on various servers, compute the best path from a requester to a given server and back, and notify the system's users of servers and requesters going on or off line. These messages can be handled automatically by the SNIOS or NETWRKIF software, or they can be implemented under the control of special application programs.

Section 4 CP/NET System Guide

The requester's NDOS and the server's SERVER module are key components in the logical structure of the CP/NET operating system. These modules, however, do not deal with the physical problems of moving a logical message from the source requester to the destination server and back again. Implementing this task varies depending on network topology, hardware, and the characteristics of the host computer systems. These modules are therefore not portable from machine to machine. You must customize them.

This section provides the network systems implementer with the information necessary to design and implement a CP/NET system efficiently. Section 4 is divided into four parts. [Section 4.1](#) discusses general network design issues that affect CP/NET implementation. [Section 4.2](#) details how to implement the requester network software, the SNIOS.SPR. [Section 4.3](#) discusses the design and implementation of the server communications software, the NETWRKIF.RSP. [Section 4.4](#) describes the design of a CP/NET server that runs under an operating system other than MP/M II. Appendixes to this manual contain several example network communications packages.

4.1 General Network Considerations

This section explains some of the basic functions of network communications software and describes, in the most general way, how communications software fits into the overall architecture. If any of the material in this section is unfamiliar to you, consult one of the many excellent textbooks available on modern networking technology. Theoretical knowledge can help you enormously in the design and implementation of your network system.

4.1.1 Functions of the CP/NET Physical Modules

The SNIOS and NETWRKIF modules function on four levels. At the lowest level, they must handle the physical transfer of a bit stream from one network node to another. This physical layer must take into account the I/O port numbers being used for communication, the physical characteristics of the network medium, network contention schemes, and other factors.

The next layer of functions must address the problem of getting complete messages from one node to another with no errors or redundant data. This data-link layer takes the bit stream from the physical layer and processes it according to its own protocol.

If any routing from node to node is required, you must include, a network-level protocol. The network layer can be as simple as identifying when a message is destined for a particular node, or it can perform complex store-and-forward operations, compute the best route from node to node, and maintain open circuits for nodes that want to communicate.

The last layer the SNIOS and NETWRKIF must address provides an interface between the low-level communications software and the logical level operating system software. In the SNIOS, this layer must transport messages to and from the NDOS. In the NETWRKIF, the transport layer reads and writes message from and to the appropriate server queues.

The layered architecture presented here can be indistinct in implementations, with single subroutines sometimes handling all four layers at once. [Figure 4-1](#) shows the relationship of the various layers to the network interface. Notice that the physical, data link, and network layers might have to participate in the interface to recover information to perform their functions.

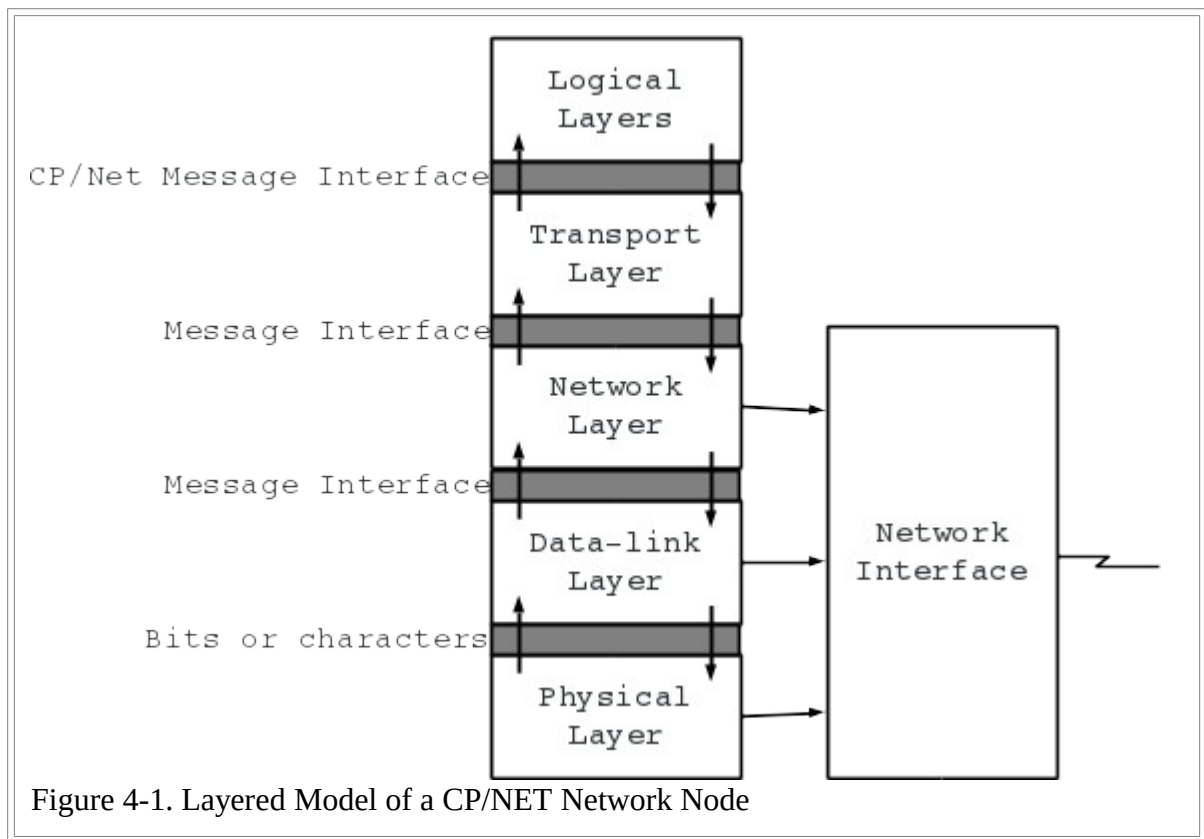


Figure 4-1. Layered Model of a CP/NET Network Node

Notice also the interfaces between the various levels. As a message migrates through the layers, the data in the message can change. The interface between the physical layer and the data-link layer yields bit or character data; the message itself is incomplete. The interface between the data-link and network layers produces messages, but the messages might contain routing information irrelevant to the transport layer. When a message reaches the transport layer, it might be in a format unusable by the higher logical layers of the operating system. Only when the message is passed to those logical layers must it be complete and in the standard format of a CP/NET message.

The architecture described above corresponds to the four lowest layers of the network model described by the International Standards Organization (ISO). However, there are some slight differences. For example, the ISO definition of the transport layer concerns itself mostly with migrating messages from a centralized network controller to one of many possible hosts. In the model described above, the transport layer deals with moving messages that have already reached a host into the correct portion of the operating system. The model in [Figure 4-1](#) is the basis for the following, more detailed discussion.

4.1.2 Interfacing a Computer to a Network

All network nodes need some method of controlling the communication functions that take place on the communications medium of the network. The simplest method is to have the node's CPU directly control all network communications protocols.

In this case, the network interface is a direct line into the host computer. When the communications software is called upon to send a message, the CPU must initiate the message, possibly waiting for an appropriate handshake response from the destination node. The CPU must then transmit the message, receive and process any acknowledgments, and determine whether the message should be retransmitted. If the node is receiving a message, it must, under program control, detect when the sender is trying to initiate a message, perform any handshake with the sender, receive the message, verify its correctness, and provide acknowledgment. All these tasks must be performed using programmed I/O operations or possibly some form of DMA for parts of the transmission or reception.

These tasks can take up a significant amount of the CPU's processing power. For an SNIOS, this is not a problem, because the NDOS is idle in the time interval after a message is sent and before the response is received. For a NETWRKIF, however, the multitasking nature of the server can result in serious performance degradation.

Another drawback to this method is that it places the burden of engineering communications software on the host systems implementer. This software can be extremely costly to develop for a high performance network.

The principal advantage of this method is its simplicity. If two computers have spare RS-232 ports, you can network them together with no special hardware. Many simple protocols can be readily modified to provide low-performance networks at low cost. Such a protocol is provided in [Appendix E](#).

For higher-performance networks, it might be necessary to relieve the host CPU of the burden of physical, data-link, and network processing. In this case, an intelligent network communications controller can be useful. Many such controllers are available, and there is a variety of methods of interfacing them to a host computer.

An intelligent communications controller can perform all physical and data-link processing, as well as many network layer functions, with no host CPU intervention. The SNIOS and NETWRKIF modules must be concerned only with a nominal amount of network routing, if necessary, and with the problem of transporting the message from the controller. Because the communications controller can transfer data to the host at high speed with high reliability, the host's transport layer can be very simple and requires little CPU time. [Appendix G](#) provides a CP/NET implementation utilizing an intelligent network controller.

Intelligent controllers require special hardware that must be added to the host computer. Interfacing this hardware is not always possible. In addition, each network node needs a controller. This can be expensive.

CP/NET also works in multiprocessor environments, both loosely coupled and tightly coupled. A loosely coupled system can send messages via a high-speed, reliable bus. This reduces the data-link problem, so simply transferring data is often sufficient to ensure the message's integrity. Tightly coupled processors can share memory, so messages can be sent between nodes by mapping memory from one processor to another.

4.1.3 Developing a Network Layer

Because CP/NET is independent of the network used, the communication modules must be modified to support various network topologies. The NETWRKIF that supports a multidrop, contention network is different from the one that supports an active hub-star configuration.

Some CP/NET configurations require extremely complex interconnections. Messages destined for one server might have to pass unmodified through several servers or requesters before they reach their final destination. The network implementer must define the software necessary to accomplish this routing. For simple networks, a network layer is barely necessary. For example, a simple work station cluster, where several requesters share a single server, requires only that the destination ID field of the message match the server's ID on a request, and that the destination match the requester's ID when the server's response is sent back to the requester.

In complex networks, each node might need to keep track of other nodes on-line in the network. Some algorithms require the exchange of routing messages to maintain an accurate picture of the topology of the overall network. To do this, the communications software must recognize these routing messages as nonstandard CP/NET messages and not pass them to a server process or to the NDOS for processing.

Even requesters might need a network layer. For example, consider a daisy-chain network of several requesters with a server at one end. All the traffic for requesters farther down the chain passes through the

requester adjacent to the server.

Because a CP/M requester can only operate a single task, the communications software for receiving and forwarding a message must be written as a series of interrupt routines. Because the NDOS might call on the SNIOS to transmit or receive a message of its own, these routines must be reentrant to the extent that NDOS requests can be held up until an intermediate message has been processed.

4.1.4 Error Recovery

Network transmission media are often unreliable. Messages are occasionally garbled or lost. In addition to data-link errors, networks can route messages incorrectly, or messages can be lost due to congestion in a section of the network. Because of these problems, a node must be able to recover from transmission errors

The most common form of error is garbled data. Bits that should have been zeros are received as ones, and ones are received as zeros. The easiest way to detect this type of error is to transmit a check along with the message. The check is computed by performing an arithmetic operation on the actual message before it is transmitted. If the check does not match the result of performing the same operation when the message is received, then a transmission error has probably occurred.

Most data-link protocols provide a mechanism for acknowledging that a message was received correctly. This mechanism requires a special message as an acknowledgment. The node that received the original message sends the special message back to the node that sent the original message. If an error occurs, the receiver either sends no acknowledgment or sends a negative acknowledgment, telling the sender to retransmit the message immediately.

The sender must be able to detect a transmission error and take steps to retransmit the message. This can be a problem because the sender does not know what the receiver is doing. If an error message comes back, the sender knows something has gone wrong. But if a message is lost completely, the receiver might not know it was sent and never send an error condition.

To solve this problem, the sender can send a message, then wait a predetermined interval for acknowledgment. If no acknowledgment arrives, the interval expires, and the sender times out. A timeout condition can cause the sender to retransmit the message or take other steps to recover from the error. When the message is finally sent successfully, the sender can free up the buffer that held it and continue with other processing.

For a CP/NET requester, two different levels of timeouts might be necessary. At the data-link level, a timeout can be set on the amount of time that elapses between sending a message and receiving the acknowledgment that it was received correctly. This timeout interval can be fairly short, since the transmission path is not likely to be very long.

The second timeout addresses the logical structure of CP/NET. Every message sent to the server implies a response to be sent back to the requester. A timeout can be set upon entering the requester's receive message routine. If the requester waits too long for a response, it can be assumed that the communication link or the server itself has crashed. With this kind of timeout, the error recovery involves much more than just retransmitting the initial message. A logical initialization must take place, probably including a CP/M warm boot.

A timeout scheme can successfully retransmit lost or garbled messages. Another problem arises, however, when the receiver's acknowledgment signal is lost. The sender, not receiving the acknowledgment, eventually times out and retransmits the message. In the meantime, the message has actually been successfully received. When the message arrives from the sender a second time, the receiver must have some way of knowing that the message is a duplicate. The receiver should ignore the message, but send an acknowledgment to stop the

sender from sending the duplicate yet again.

The easiest way to detect duplicates is to assign a sequence number to each message. If the receiver does not receive the sequence number it was expecting, it ignores the message, even if the message was received correctly. Every time a message is received, the expected sequence number is incremented. Every time the sender receives an acknowledgment, the sequence number to be sent is incremented. If a message times out, however, the sequence number is not incremented.

All error recovery schemes should be free from deadlocks. A deadlock occurs when the sender is waiting for an action from the receiver, but the receiver is not performing that action because it is waiting for the sender to perform another action. Carefully analyze networks that store and forward messages from node to node for deadlocks because two nodes can try to transmit to one another simultaneously.

The means of avoiding deadlocks varies according to the network topology. A multidrop network can use collision detection. If two nodes attempt to use the network at the same time, they immediately detect that their messages are garbled and stop transmitting. To avoid continuous collisions and a consequent deadlock condition, the two nodes attempt to transmit again based on a random time interval, so that one node can start transmitting before the other.

In a point-to-point network, a properly designed message handshake can often avoid data-link deadlocks. At a higher level, enforcing a buffer allocation protocol can often prevent deadlocks. Waiting to transmit messages until the receiver has space for them minimizes the possibility of two messages continuously timing out.

4.2 Customizing the Requester's SNIOS

The communication interface between the logical NDOS and the actual network is contained in the Slave Network I/O System module, SNIOS.SPR. Because this interface varies depending on the computer system and network hardware, you must customize the SNIOS.

For most applications, the SNIOS need only be a sequential system. The SNIOS never needs to respond asynchronously to unsolicited messages. Only the NDOS must direct the SNIOS to receive messages. However, some networks require real-time response from their SNIOS modules to pass a message between two network nodes that have no direct means of communicating with one another.

This section details the design and preparation of an SNIOS for inclusion with a CP/NET requester and describes the installation of the utilities necessary to run the requester.

4.2.1 Slave Network I/O System Entry Points

The SNIOS must begin with a jump vector containing the network I/O system entry points, as shown below:

```
SNIOS:  JMP NETWORKINIT      ; Network initialize
        JMP NETWORKSTS     ; Rtn network status
        JMP CONFIGBLADR    ; Rtn Config. Tbl Adr
        JMP SENDMSG        ; Send msg on network
        JMP RECEIVMSG      ; Receive msg from ntwk
        JMP NTRKERROR      ; Network error
        JMP NTRKWBOOT      ; Network warm boot
```

Listing 4-1. SNIOS Jump Vector

Each jump address corresponds to a subroutine that performs the specific function. The exact responsibilities

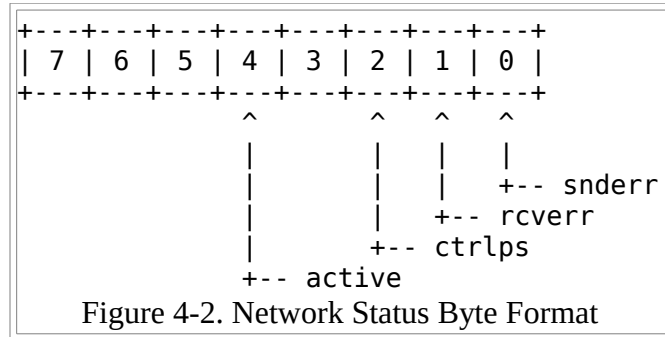
of each entry point subroutine are given below.

NETWORKINIT

This SNIOS entry point is called when control is transferred to the NDOS initialization entry point after being loaded by the CPNETLDR. This subroutine performs any required network interface initialization. Initialization includes reading back-panel switches, or some other suitable source, to obtain the requester processor ID for the configuration table. If initializing messages must be sent out over the network, send them from this routine.

NETWORKSTS

This subroutine returns a single byte in register A and determines the status of the network interface. The error bits snderr and rcvrr are reset when the call is made. The format of the network status byte is shown in [Figure 4-2](#).



- active = 1 if requester logged in
- ctrlps = 1 if control P is active
- rcvrr = 1 if error in received message
- snderr = 1 if error in sending a message

CONFIGTBLADR

This subroutine returns the requester configuration table address in the HL register pair. The requester configuration table is described in [section 4.2.2](#).

SENDMSG

This subroutine enables messages to be sent from one processor to another via the network. The passed parameter, in registers BC, is a pointer to the message. Control is not returned from this procedure until the message has been sent. Thus, the message pointed to by the BC register pair can be modified immediately upon return. The return code, in register A, has a value of 0 indicating success or 0FFH indicating failure to access the network.

RECEIVMSG

Messages are received from another processor through the network with this subroutine. The passed parameter, in registers BC, is a pointer to a message buffer. Control is not returned from this procedure until the message has been received and placed into the message buffer. Thus, the message in the buffer is valid immediately upon return. The return code, in register A, has a value of 0 indicating success or 0FFH indicating failure to access the network.

NTWRKERROR

When network errors are encountered, this procedure is called. Any required network interface device reinitialization should be performed. In typical SNIOS implementations, executing a return from the NTWRKERROR procedure results in a retry. If a retry is not wanted, an appropriate message is displayed on the console, and a warm boot is performed.

NTWRKBOOT

This SNIOS procedure is called each time the NDOS reloads the CCP. The sample SNIOS in Appendix E displays a

<Warm Boot>

message on the console only as a demonstration of NTWRKWBOOT. More practical applications of this procedure include interrogating the CP/NET server for messages. In this way, each time a warm boot is performed, the user is notified of messages posted for him.

4.2.2 Requester Configuration Table

The configuration table that resides in the CP/NET requester's SNIOS allows reassignment of logical devices to networked servers. The configuration table creates a mapping of logical to physical devices that can be altered during CP/NET processing. The configuration table specifies the system I/O to be accessed through the network.

The requester configuration table is defined in [Table 4-1](#).

Offset	Explanation
000-000	Requester status byte
001-001	CP/NET requester processor ID
002-033	Disk Devices; 16 two-byte pairs, first byte high-order bit on = drive on network, with the server drive code in the least significant 4 bits; the second byte contains the server processor ID.
034-035	Console Device; first byte high-order bit on console I/O on network, with the server console number in the least significant 4 bits; the second byte contains the server processor ID.
036-037	List Device; first byte high-order bit on = list to network, with the server list device number in the least significant 4 bits; the second byte contains the server processor ID.
038-038	List Device buffer index.
039-043	List Device logical message header: FMT, DID, SID, FNC and SIZ.
044-044	List Device server list device number.
045-172	List Device buffer.

Table 4-1. Requester Configuration Table

4.2.3 Preconfiguring the Configuration Table

In many network systems, there is never any need to modify the device mappings specified through the NETWORK utility. In such systems, you can preconfigure the device mappings in the configuration table. To do this, select the devices to be networked and set the high-order bit of the first byte in the entries corresponding to those devices. Set the remote device to which the local device is to be mapped in the low-order four bits of the same byte. Finally, set the server ID of the remote device in the second byte of the entry.

Be careful when preconfiguring devices to servers that might be off line. Some CP/NET functions send messages to all servers referenced in the configuration table. If one of these servers is not capable of receiving messages, functions that might subsequently send messages to servers on line can prematurely abort.

For example, the CCP might issue a free drive function to initialize the server environment for a subsequent application program. If the previous application had left files open on two on line servers, but a third server was off line, those files are left open if the free drive message was sent to the off-line server before the on-line servers. The next application program might damage the files that were inadvertently left open.

You can solve this problem by having the error recovery in the SNIOS remove any networked device that experiences continuous timeouts, converting it back into a local device. This prevents the NDOS from making continuous references to the off-line server. A major drawback of this scheme, however, is that an application might suddenly begin referencing a local device, possibly destroying files on a local disk drive. A more secure, but less friendly protocol for dealing with off-line servers is to force a warm boot whenever a network error is encountered.

It is wise to enforce a protocol that prohibits devices from being networked until the server to which they are assigned is on line. Special utilities can be written to accomplish this by sending a dummy message to every server to which drives are mapped.

4.2.4 Sending and Receiving Messages Asynchronously

In some networks, a requester might have to receive and retransmit asynchronously a message destined for another node. For example, consider a loop network, where every node has two network ports. The network protocol specifies that all messages are sent via port #1, and all messages are received via port #2. If there is only one server in the network, but more than one requester, all messages must pass through every other requester, either as they are sent to the server or as the response returns from the server.

If a requester must asynchronously handle a communication channel, it must do so outside of the facilities provided by the single-tasking CP/M operating system. The communication protocol must be interrupt driven. An interrupt service routine must at least detect the start of a message; after that, the rest of the message can be handled sequentially or under control of additional interrupt routines. If a requester cannot support interrupts, asynchronous handling of messages might be impossible. Neither the application program nor the NDOS can periodically check for incoming messages.

A mechanism must be provided so that the NDOS, sequentially calling the SNIOS to send a message, does not collide with the asynchronous transmission of another message. Receiving messages cannot collide because only one message can come over the network at a time. To accomplish this, consider implementing the loop network described above.

As a requester's application is running, another node suddenly starts sending a message to it. The requester must now receive the message, verify its correctness, and retransmit it to another node. All of these operations must be performed without damaging the local application program. If the data-link routines do not make CP/M system calls and do not modify the message buffers used by the NDOS, the entire message can be received and transmitted transparently. When this operation is finished, the interrupt service routine returns to the application program, and processing continues. When the NDOS needs to use the network, the same data-link routines that handled the asynchronous message can be used to handle the sequential one.

It is even possible to transmit a message from the NDOS while receiving a message from some other node. To do this, the message must be able to be received a piece at a time, giving both the send and receive routines enough processor time to avoid timing out. Such a system requires a mechanism for preventing both the NDOS and the interrupt service routine from attempting simultaneous transmission. A semaphore variable can be used to control the system.

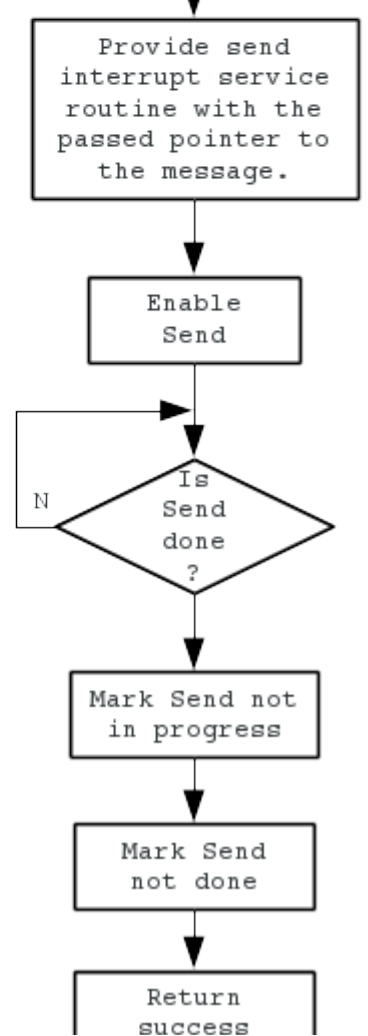
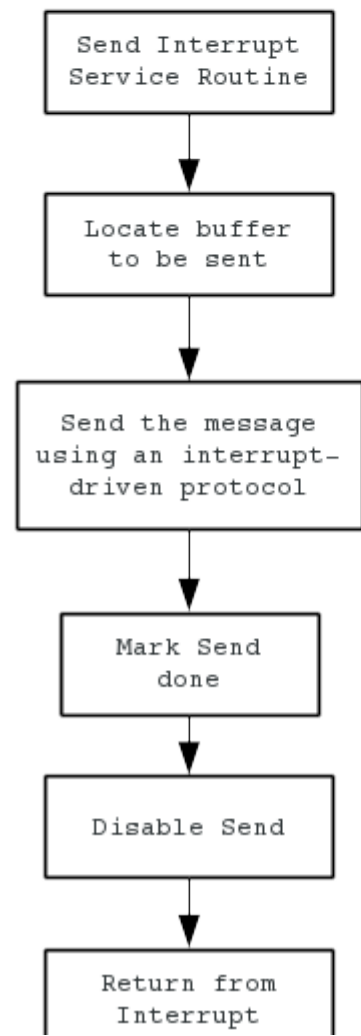
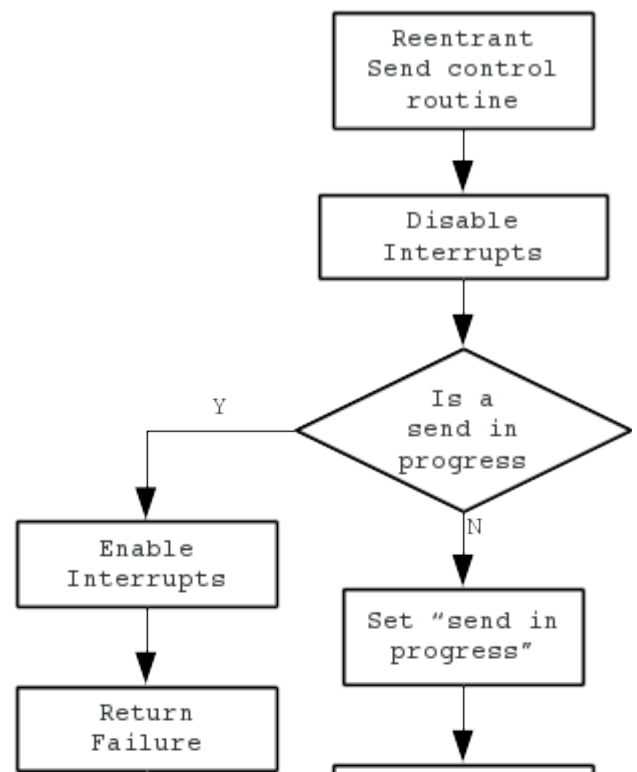
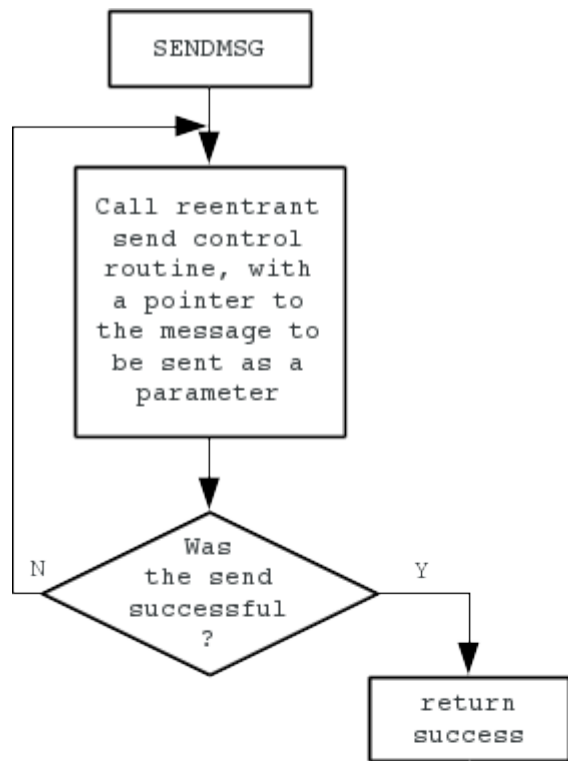
[Figure 4-3](#) outlines a possible protocol for such a system. Both the SNIOS SENDMSG routine and the asynchronous receive interrupt service routine access a piece of reentrant code to control access to the message transmission system.

Three external events drive the system:

- The NDOS can request to send a message.

- The NDOS can request to receive a message.
- A message, unbidden, can cause an interrupt so that it can be received.

In this implementation, the message sending software is interrupt driven, started by enabling a transmitter interrupt. The message sending software can also operate sequentially, called by the reentrant routine that controls its use.



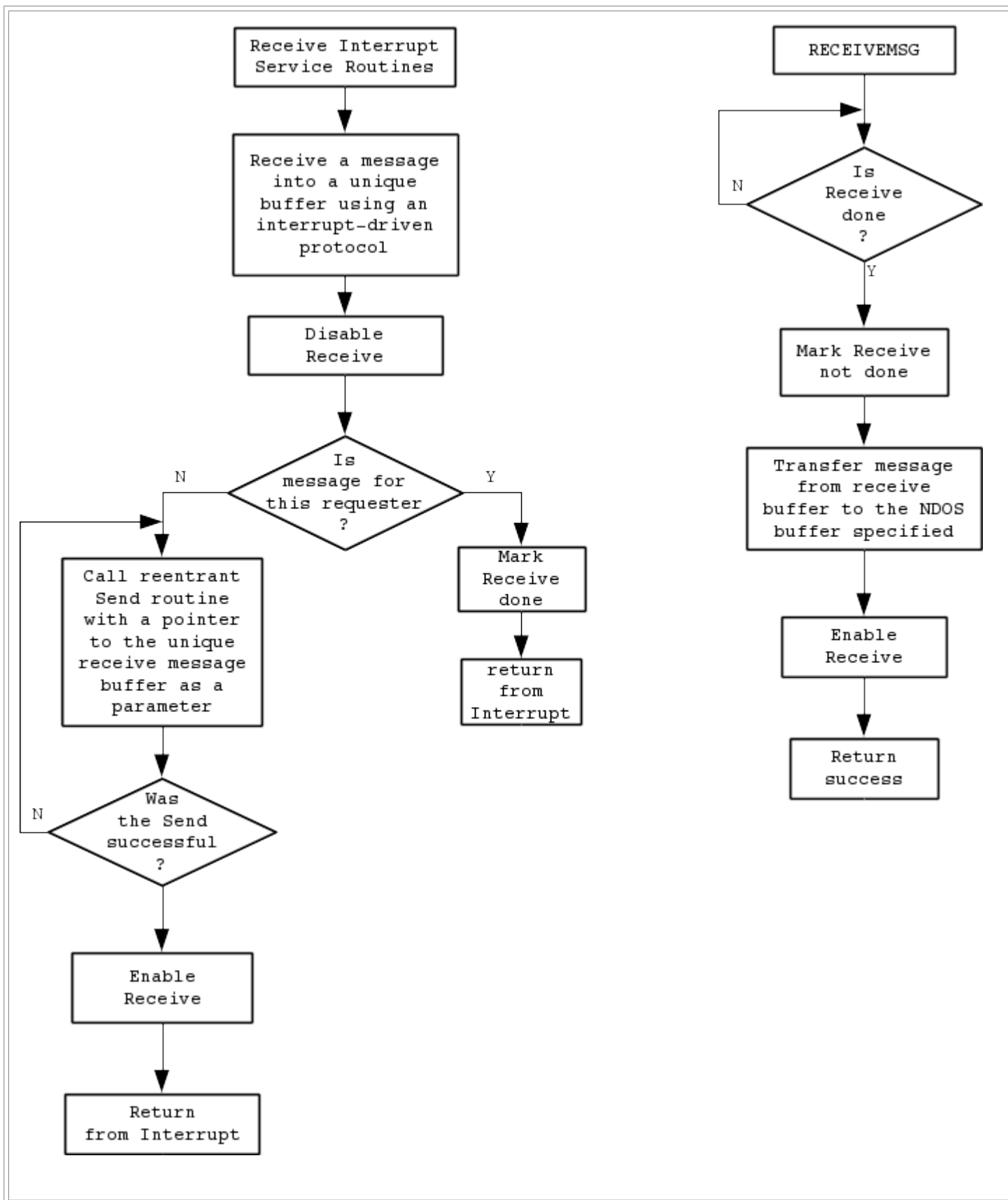


Figure 4-3. Algorithm for Interrupt-driven Requester Node that Stores and Forwards Messages

4.2.5 Generating and Debugging a Custom SNIOS

Follow these steps to generate and debug a custom SNIOS.

1. Prepare the SNIOS.SPR file, as shown below:

```
A>RMAC SNIOS
A>LINK SNIOS[0S]
```

The output of the linker is the SNIOS.SPR file.

If you do not use RMAC and LINK-80 use ASM, PIP, and GENMOD, as shown below:

Assemble with ORG 0000H.

```
A>ASM SNIOS
A>REN SNIOS0.HEX=SNIOS.HEX
```

Edit the SNIOS.ASM ORG statement. Assemble with ORG 0100H.

```
A>ASM SNIOS
A>REN SNIOS1.HEX=SNIOS.HEX
```

Concatenate the HEX files.

```
A>PIP SNIOS.HEX=SNIOS0.HEX,SNIOS1.HEX
```

Generate the SNIOS.SPR file.

```
A>GENMOD SNIOS.HEX SNIOS.SPR
```

The GENMOD program uses the difference in code origins to produce a bit map of addresses to be relocated. GENMOD then places this bit map at the end of a copy of the origin 0 code and constructs a 256-byte header to create an SPR file.

2. Copy the following files to the requester:

- CPNETLDR.COM = Loads CP/NET (NDOS.SPR and SNIOS.SPR)
- CPNETSTS.COM = Displays status of the system I/O
- NETWORK.COM = Redirects I/O from local to network
- LOCAL.COM = Redirects I/O from network to local
- DSKRESET.COM = Resets specified logical drives
- LOGIN.COM = Logs on to server
- LOGOFF.COM = Logs off from server
- MAIL.COM = Electronic mail utility
- NDOS.SPR = Network Disk Operating System
- SNIOS.SPR = Previously Customized Slave Network I/O System
- CCP.SPR = Console Command Processor

you can use DDT to debug the SNIOS as follows:

```
A>DDT CPNETLDR.COM
*IB
```

```
*s103
0103 07 xx
*g
```

where xx is the restart the debugger uses, usually 7.

At this point, CP/NET loads, displaying the memory map, and then breaks at the specified restart. You can place breakpoints at desired locations, and then issue a G command specifying the address following the restart instruction where the CPNETLDR broke.

Communications software is difficult to debug. Because of its real-time nature, when the program is interrupted to find out what is going on, the other side of the network overruns or times out. These pointers might help you:

- Before debugging, disable any timeout logic in both the SNIOS and the NETWRKIF. This allows one node to be examined without causing errors on the other node. The SNIOS example in [Appendix E](#) accomplishes this with a conditional assembly switch called ALWAYS\$RETRY.
- Never set a breakpoint in the SNIOS without setting a corresponding breakpoint in the NETWRKIF.
- Write a simulation module that mimics how you think the NETWRKIF should behave in response to the actions the SNIOS takes to send a message. Disable the actual network transmission until the SNIOS can successfully send messages to and from the simulation. Gather copious statistics because when you finally transmit over a real network link the simulation and the real NETWRKIF probably will not correspond. The statistics can help point up what was wrong with the simulation, the NETWRKIF, or both.
- Carefully verify any communications handshakes between the two nodes. You can do this by stepping through the code of both nodes simultaneously, using debuggers. Discover which data link operations can be performed while the other node is halted or disabled. Quite often, making a mistake in your debugging session points up holes in your protocol design. Once you have the protocol working with this method, have someone step one node while you step the other. Do not coordinate the actions of the two debuggers. If your protocol works without conscious synchronizing, try running it full speed.
- If possible, write one data-link module for both the SNIOS NETWRKIF, then interface them to the appropriate module. This enhances the uniformity of the protocol, making it easier to debug.

4.3 Customizing the Server

This section addresses the problems of designing and implementing an efficient CP/NET server under the MP/M II operating system. Because a CP/NET server must be capable of handling several simultaneous requests in real-time, the Network Interface module (NETWRKIF) must take full advantage of the real-time primitives of MP/M II.

The server's logical module, SERVER.RSP, consists of a set of processes, one for each requester supported. This section also discusses how the NETWRKIF sends and receives messages to and from those processes.

Finally, this section explains the system generation options available to the server implementer once the NETWRKIF has been implemented.

4.3.1 Detecting and Receiving Incoming Messages

The server is a passive, asynchronous system; it does not initiate CP/NET transactions. The server performs two distinct functions:

1. The server must detect an incoming message and initialize the communications software to receive.

2. The server must actually receive the message.

The server detects incoming messages in two ways. The first is polling, where the server periodically checks the status of the network interface. If the status changes from an idle to a ready state, the server receives a message. The second method relies on the network interface's interrupting the server. The server then transfers control to a service routine that receives the message. Either of these methods can accomplish the two functions listed above. Both methods have advantages and drawbacks.

Polling the Server

Polling is a more active method, requiring more processing overhead. If the server has a fairly heavy, continuous load of network traffic, then the status of the poll operation often indicates that a message is to be received. In this kind of system, polling has a marked advantage: the server can immediately begin receiving the message without switching contexts. But if the network traffic is subject to bursts of data mixed with periods of traffic, then the extra overhead of interrogating the network interface is inefficient.

Interrupting the Server

Interrupt driven operation is excellent for communication that occurs in bursts because no overhead is required when no communication is taking place. But very high network loads cause the server to waste a great deal of time saving the state of the process currently executing when the interrupt occurred.

Once a message has been initiated, it can be received under interrupt control, where data is processed on demand as it comes in, or under direct program control, where a process is dedicated to monitoring the incoming message. The most efficient choice depends on the type of network being used and the amount of traffic the network must handle.

In an interrupt driven communication scheme, the server responds to network events asynchronously. The network interface determines when data is processed by the host CPU. For example, when the network interface presents characters to the host, each character causes an interrupt. When the network interface performs direct memory access to transfer blocks of data, only each complete DMA transfer causes an interrupt. Depending on the protocol, each interrupt causes a specific action to be performed. The CPU is free, however, to process other tasks in between processing each piece of data. Like interrupt-driven message detection, saving the state of an interrupted process requires CPU overhead. The greater the number of interrupts required to process a message, the more system performance is degraded.

Overruns

One of the greatest problems of an interrupt-driven communications scheme develops when the interrupts occur faster than the CPU can service them. This condition is known as an overrun, and it can cause data to be lost. When an overrun occurs, the message appears to be garbled, and the sender must retransmit it. If overruns occur only when the host is extremely busy, it might be more efficient to accept the occasional garbled message in exchange for better overall response. If the number of overruns is too high, however, serious system degradation sets in. Many protocols prevent overruns by allowing the receiver to signal the sender that data is coming in too fast.

Disabling Interrupts

The other approach to message processing uses MP/M II's facility to control processes. Unlike an interrupt service routine, which is largely transparent to MP/M II, a process is a logically complete task. Using a process-oriented protocol, you can eliminate the overrun problem by disabling interrupts while the message is being received. Disabling interrupts gives the communication program exclusive control of the CPU, so all other processing comes to a halt. If messages are fairly short, however, this method might be preferable to an interrupt-driven scheme, because no overhead is incurred by switching back and forth between a process and an interrupt service routine continually.

Selecting a Protocol

The actual data-link protocol used to process messages has not been discussed. Consider the selection

of a protocol when designing how the server is going to respond to incoming messages. For example, in a CP/NET system where loosely coupled processors are communicating over a high-speed bus with little or no error checking, DMA transfer of data can be efficiently interrupt driven. But complex cyclic redundancy checks that involve extensive arithmetic operations require careful design in an interrupt-driven system, or overruns might result. Such a protocol might be better implemented using a process-oriented system.

4.3.2 NETWRKIF Module Architecture

[Section 4.3.1](#) discusses general strategies for implementing a data-link layer protocol under MP/M II. This section deals with integrating the data-link layer into a network and transport layer. This integration allows the entire communications package to send logical requester messages to the SERVER.RSP module, and then receive the SERVER's response message for transmission back to the requester.

A dedicated server process is associated with each requester logged on to a server node. These processes are named `SEVR<x>PR` where `<x>` is an ASCII character between 0 and 9 or A and F. This character is a sequence number that serves as a unique identifier for the server process. Each server opens two queues that it expects the NETWRKIF module to have created. They are named `NtwrkQI<x>` and `NtwrkQO<x>` where `<x>` is the same character as the server's sequence number. The server process always reads the address of incoming messages from `NtwrkQI<x>`, and it always writes the address of the response message to `NtwrkQO<x>`.

This is the basic interface between the SERVER.RSP module supplied by Digital Research and the user-customized communications software. However, there are a variety of ways to implement the processes driving the interface.

[Appendix E](#) includes an example of the simplest NETWRKIF architecture. In this architecture, one network interface process is associated with each server. All processes execute the same reentrant code, but each process maintains local data that identifies the communications port it is using and the sets of queues through which it interfaces to the server process. This implementation handles its data-link software at the process level. It uses polled console I/O functions in the XIOS to detect incoming messages. This architecture is illustrated in [Figure 4-4](#).

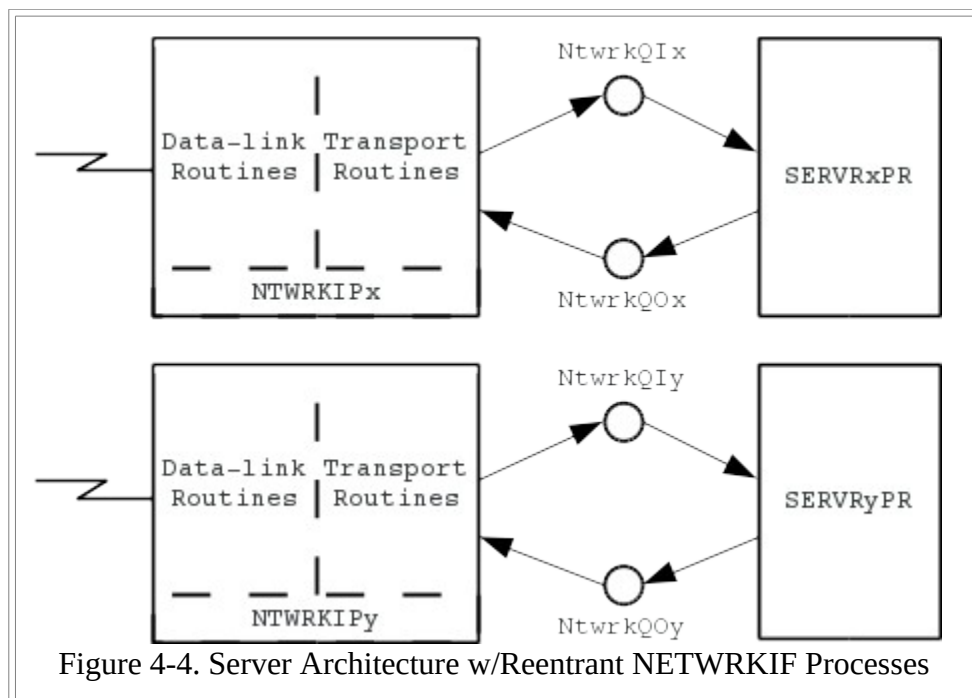


Figure 4-4. Server Architecture w/Reentrant NETWRKIF Processes

Another possible NETWRKIF architecture has only two network interface processes. An input process receives data from the network, identifies the requester that sent the message, and writes the message to the appropriate queue. An output process conditionally reads all the output queues and sends any messages it finds back out over the network.

It is also possible to force all the server processes to write their messages to a single queue by patching SERVER.RSP. In this case, the output network interface process reads the single output queue. When a message is written to it, the output process sends the message out across the network and goes back to read the queue again. An application note details how to patch SERVER.RSP. [Figure 4-5](#) illustrates both strategies. Note that a small patch to the SERV<x>PR processes can consolidate the output queues.

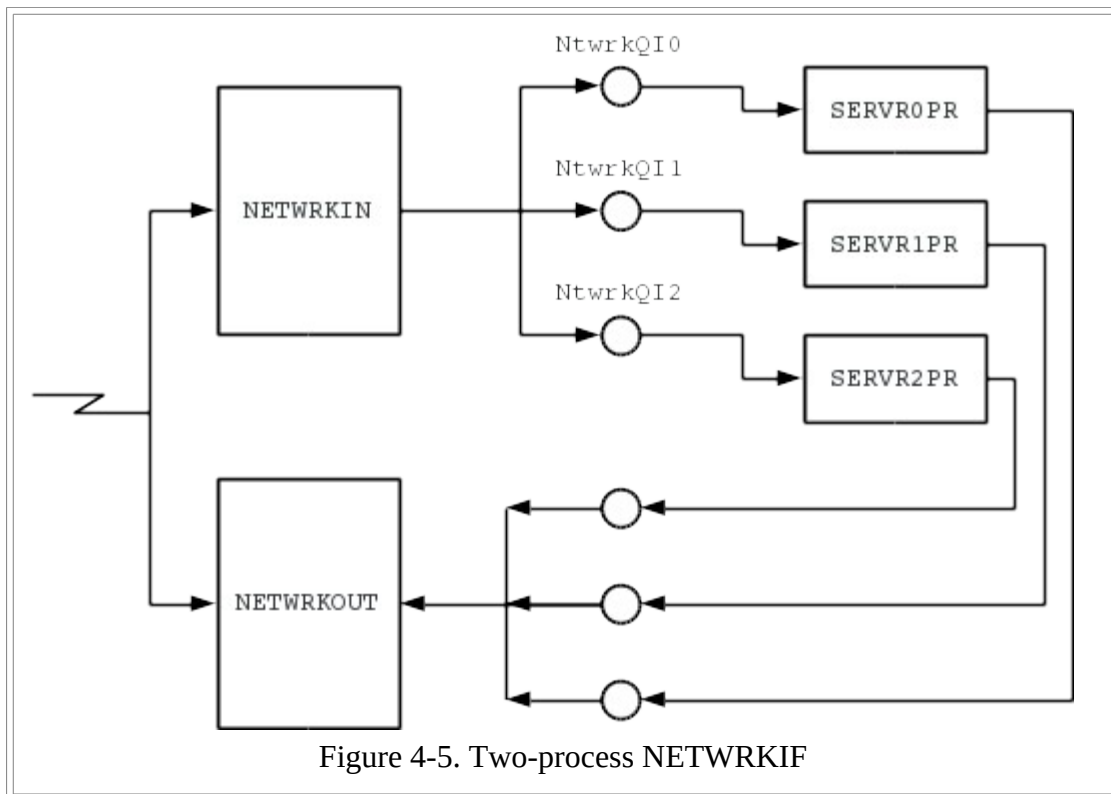


Figure 4-5. Two-process NETWRKIF

You can design a single NETWRKIF process that receives a message, writes it to the appropriate queue, then checks for any output activity. If NETWRKIF finds a message to send, it sends it, then it returns to checking for input. This kind of process has the disadvantage of being constantly busy; there is no point at which it can allow itself to become blocked. To do so might result in a deadlock or serious performance degradation.

Consider the network topology when designing the NETWRKIF architecture. For example, a NETWRKIF that uses one process per requester is suitable in an active hub-star configuration, where a unique network line is dedicated to each requester. This allows several messages to arrive at the server simultaneously.

For a multidrop topology, however, a single output and single input process NETWRKIF might be more suitable, because the network hardware guarantees that only one message is active on the network at any one time. The same type of architecture could be applied to a loop topology.

For an active hub-star network that services several multidrop lines, it might be necessary to combine the two architectures, so that several reentrant processes are routing input to the server processes, while a set of output processes are collecting data from output queues and sending it back out of the appropriate multidrop line.

Also consider what the NETWRKIF does when it has no traffic to process. If the NETWRKIF loops madly while waiting, it will gobble up precious CPU resources, degrading the overall performance of the server system. On the other hand, the NETWRKIF must be able to respond to traffic quickly.

A number of MP/M II system calls cause a process to become blocked, so that the operating system dispatcher does not pass control back to the process until a critical condition is fulfilled. Reading an empty queue, waiting on a flag, and performing a poll call are three of the most common ways to suspend the execution of a process conditionally. Such quiescent points should be built into all NETWRKIF systems to minimize the overhead of maintaining the process when it is idle.

The processes driving the input and output queues constitute one half of a message transport layer. The NETWRKIF must also deal with how the raw message is received from the data-link and network layers that are performing the actual communication control. This interface is governed by how the data-link and network layer software is implemented.

Consider an architecture that has little or no network layer, so that the data-link software interfaces directly with the transport processes. If the data-link is included in the processes that are also performing the queuing functions, then no special interface is needed. The process can pass control from one function to another, first performing input data-link and network activities to receive a message; then computing the routing to the appropriate server input queue; then reading the response from an output queue; and finally returning to the data-link level to send the response back to the requester. The sequence can be repeated indefinitely.

Some implementations require the data-link and network layers to be under process control, with a separate set of processes controlling the transport layer. In these cases, the transport processes can use queuing for both the low-level interface to the data-link layer and the upward interface to the server processes.

This kind of architecture has the drawback of slowing down the MP/M II dispatcher with extra queuing overhead. For a small number of processes, however, the impact is slight. The architecture has the advantage of being highly modular, facilitating the future upgrade of the data-link and network layers or the transport layers. [Figure 4-6](#) details the architecture.

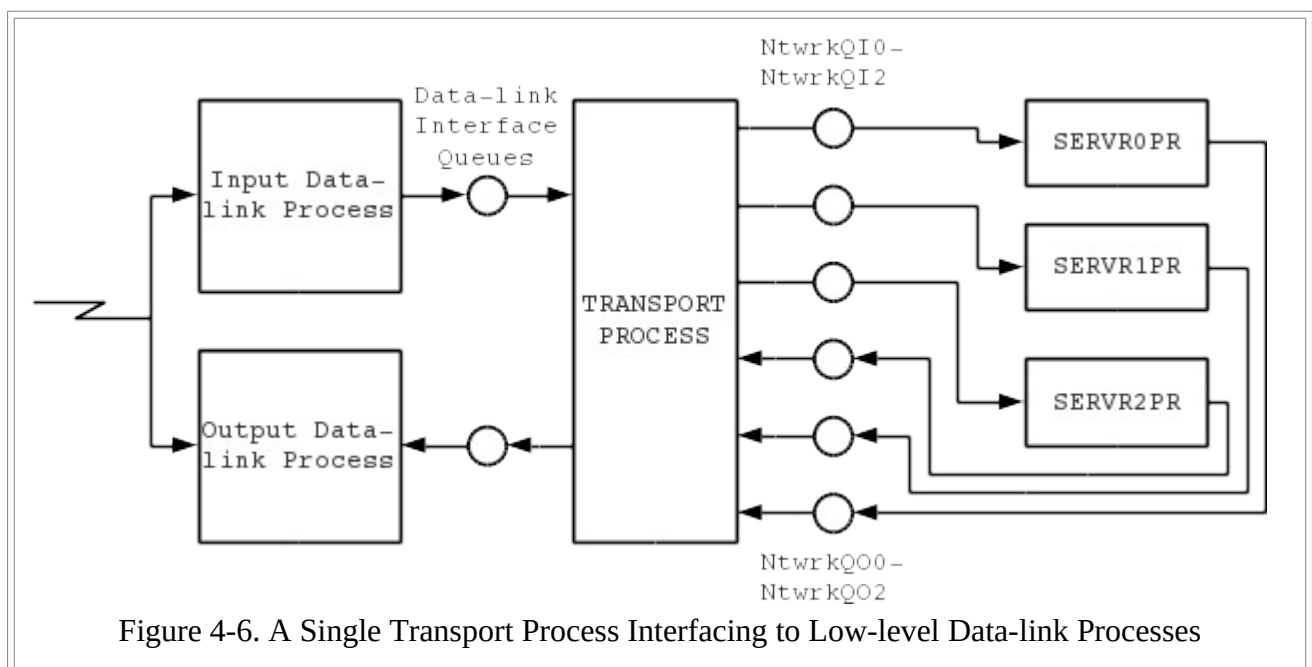


Figure 4-6. A Single Transport Process Interfacing to Low-level Data-link Processes

To implement some network interfaces, it is necessary to modify the MP/M II XIOS. Interrupt service routines must access the system interrupt vector, which is usually maintained by the XIOS. If an interface

routine requires polling, the routine to accomplish the polling must be placed on the list maintained by the XIOS POLLDEVICE routine.

Interfacing to data-link and network routines that reside in the XIOS is slightly more complex than interfacing to routines contained in the NETWRKIF. These routines are often not processes, but shared code fragments or interrupt service routines. They cannot use queues as an interface mechanism. Routines that are not process-oriented must communicate through a direct function linkage, through polling, or through the Flag Set/Flag Wait functions supported by MP/M II.

Because the NETWRKIF might not be able to resolve references to such routines directly, it is often necessary to enter the XIOS through its jump vector. The XIOS jump vector table is always page aligned; a pointer to that page is located in byte 7 of the MP/M II system data page - From this point, data-link routines can be called by specifying dummy console I/O or dummy list device I/O.

If dummy console or printer I/O is used, the NETWRKIF loads a non-existent device number in register D and, if necessary, a pointer to a message buffer. The I/O routine specified checks for the non-existent device number and dispatches the call to the appropriate network routine.

[Figure 4-7](#) illustrates how the NETWRKIF module can perform calls to subroutines resident in the XIOS.

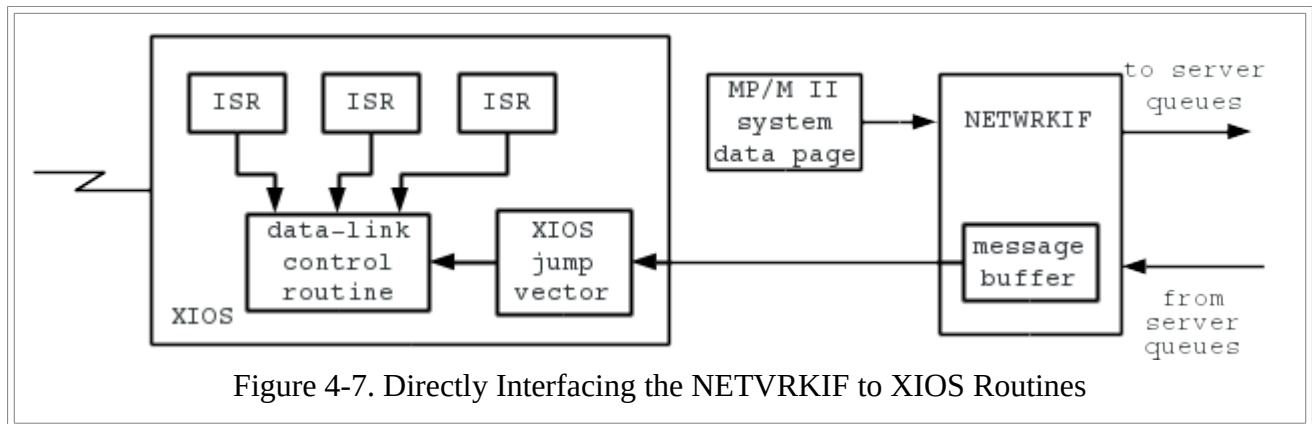


Figure 4-7. Directly Interfacing the NETWRKIF to XIOS Routines

Another method of interfacing data-link and network layer routines to a transport NETWRKIF is to have the low-level routines set a flag when a message has been processed. For example, consider a data-link routine that reads in an incoming message and checks it for validity. This routine might be a set of vectored interrupt service routines.

At this point, the NETWRKIF is not synchronized with the data link routine. When the NETWRKIF requires a message, it issues a flag-wait call to MP/M II. When the data-link routine has a complete message, it issues a flag set call. The NETWRKIF does not proceed until the flag has been set. The NETWRKIF can then transfer the message from a predefined buffer and transport it to the appropriate server process.

This type of architecture is ideal for allowing intelligent network controllers to drive the NETWRKIF transport processes. A simple interrupt service routine locates the message, builds a control block, and sets a flag to inform the NETWRKIF of the status and location of the message. [Figure 4-8](#) shows a similar interface.

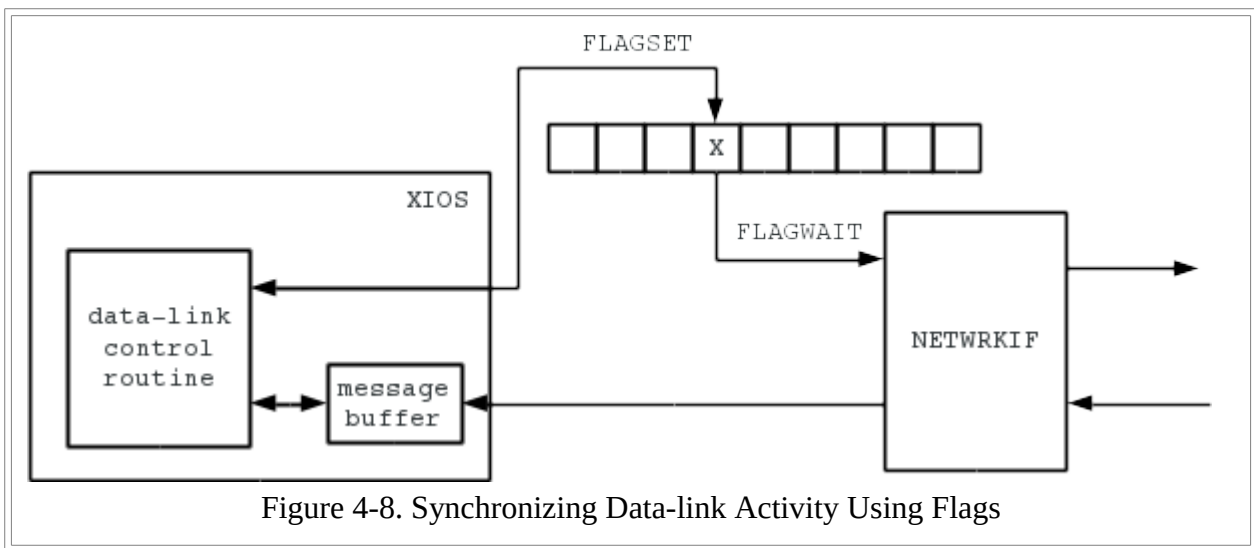


Figure 4-8. Synchronizing Data-link Activity Using Flags

To send a response message back to a requester using flags, the transport process must first identify the message to be sent and instruct the data-link layer to send it. A predefined control block can accomplish both operations. The transport process then waits on a flag until the message is sent and the flag set by the data-link.

Another possible synchronization mechanism is through the MP/M II Poll function. With this function, MP/M II suspends the calling NETWRKIF process but periodically interrogates the status of the data-link and network software through a small code fragment defined in the XIOS POLLDEVICE routine. When the status becomes true, MP/M II allows the NETWRKIF process to proceed.

If the server system supports vectored interrupts, and the location of the system's interrupt vector is known, you can write interrupt service routines that reside inside the NETWRKIF module. When the NETWRKIF performs its initialization, it simply writes the addresses of various interrupt service routines into the vector. From then on, any reference to those vector locations results in the execution of the NETWRKIF's ISRs.

This approach preserves system modularity and allows the network implementer to implement low-level routines when the XIOS itself is not available for modification. This approach still requires a synchronization mechanism between code fragments that are not part of any process and the more well-defined transport processes of the NETWRKIF.

In addition to synchronizing with low-level communications software, NETWRKIF processes might have to compete for data-link resources. For example, a transport process that wants to send a message might have to be suspended while another process is busy receiving a message. Or two reentrant processes might try to send a message out across the same network line simultaneously. These conflicts can be resolved through use of mutual exclusion (MX) queues.

An MX queue contains only one dummy message, called a token. In order to control a resource, a process must first acquire the token, leaving the MX queue empty. If another process already has the token, the first process is suspended until the second completes its resource-critical operation and replaces the token.

In this way, two low-level data-link routines--one for sending and one for receiving--can be driven without collisions by their higher-level transport processes, even if the low level routines have no explicit mechanism for sharing a network resource.

Just as the design of the network topology and error recovery schemes for CP/NET must be examined for potential deadlocks, so must the server architecture itself. A simple example of a deadlock is a process that competes for a resource using an MX queue but never restores the token to the queue when it is finished with

the resource. All the other processes waiting for the resource come to a grinding halt, the network becomes congested, and eventually everything stops.

Finally, you can design an architecture that distinctly divides the data-link, network, and transport layers. The preceding synchronization strategies can be generalized to work across several layers just as easily as they can work when the server architecture divides the communications software into low-level and high-level segments. Remember that as the architecture grows more and more complex, performance of the MP/M II dispatcher and nucleus software degrade further and further. It is always wise to keep the architecture as simple as possible.

4.3.3 Elements of the NETWRKIF

This section defines the data objects and processing required to allow the server to be initialized and to operate smoothly and continuously. Through these objects, you define how many requesters a server can handle at once and how many messages can be simultaneously processed.

The following objects must be present to create the NETWRKIF.RSP module:

- XDOS entry point
- Transport Process Process Descriptors
- Transport Process Stacks
- Queue control blocks (QCBs) for the interface between the NETWRKIF and the server processes
- User queue control blocks (UQCBs) to allow the NETWRKIF to access the queues
- Message buffers
- The server configuration table
- Stack space for additional server processes, if more than one requester is to be serviced at a time
- Areas allocated to contain more server Process Descriptors, if more than one requester is to be serviced at a time
- Network initialization code
- Data-link interface code
- Message validity checking and reformatting
- Server process interface code

XDOS Entry Point

All resident system processes (RSPs) require a linkage to MP/M II's XDOS entry point because the Command Line Interpreter does not prepare an execution environment for them. This linkage is always the first two bytes of the module. When the implementer runs the MP/M II GENSYS utility to include the server modules into the operating system, GENSYS automatically fills in these two bytes with a pointer to the XDOS entry point. This allows the execution of MP/M II system functions within the body of the RSP by setting up the function parameters, loading this pointer, and dispatching.

NETWRKIF Process Descriptors

Immediately following this pointer, MP/M II expects to see a Process Descriptor. It automatically creates and executes the process to which the Descriptor refers. In the case of the NETWRKIF, this Process Descriptor controls the execution of one of the server transport processes. These processes perform the queue read and write operations to move messages into and out of the server processes. The first process must also be responsible for server and network initialization and for creating any additional transport processes.

Process Descriptors for additional transport processes must also be included, if the processes are necessary. These processes can be automatically created by linking them to the first Process Descriptor. Linking is accomplished by placing a pointer to the second Process Descriptor in the PL field of the first Process Descriptor, a pointer to the third in the PL field of the second, and so on. The chain of

links terminates with a zero in the PL field of the last Process Descriptor to be created.

If you choose to have processes automatically created, remember that once processes are created, they are completely independent unless they are explicitly synchronized. The processes should not be dependent upon the first process to perform initialization for them.

Run transport processes at a very high priority, so that messages tie up the communications software for as little time as possible. The example in [Appendix E](#) runs at priority level 64, exactly the same priority as the server processes. For compute bound NETWRKIF processes, it is advisable to give the server a slightly higher priority than the NETWRKIF. The implementation in [Appendix F](#), for example, runs at a priority of 66. This forces MP/M II always to process logical messages first if both the server and transport processes are ready at the same time.

Each transport process must have its own local stack area. Because RSPs do not have access to the extra user stack space on system calls, each stack must be capable of supporting the local storage required by the MP/M II XDOS and XIOS in addition to its own local storage.

When a process is created, its Process Descriptor's STKPTR field should point to the top of its associated stack. The top of the stack must contain the starting execution address for the process.

Queue Control Blocks

The NETWRKIF module must contain all of the queue control blocks for the entire server system. The number of QCBs varies depending on how many requesters the server system supports at one time. For each requester, there must be one input queue, named NtwrkQIO, NtwrkQI1, and so on. There must also be one output queue per requester, named NtwrkQOO, NtwrkQOI, and so on. These queues must also be created by the NETWRKIF module.

You can patch the server process code so that all processes open the same output queue, NtwrkQOO. If this patch is applied, the NETWRKIF need only include the one output QCB. The NETWRKIF examples in [Appendixes E](#) and [G](#) use this method.

The input and output queues communicate the address of the message buffer containing the message to be processed by the server or the response to be sent back to the requester. Because the message passed through the queue is only two bytes long, circular queues can be used. Both input and output queues need only buffer one message at a time because a requester must have always received a response before sending another request. Consequently, there is never more than one message from a given requester at the server at a time.

A queue capable of buffering more than one message is required only when the server processes have been patched to write all of their responses to a single queue. In this case, the queue must be capable of buffering the output from all of the servers simultaneously.

User Queue Control Blocks

Transport processes must read and write queues using user queue control blocks. These data structures contain a pointer to the appropriate QCB and a pointer to the message to be written. The queue passes only the addresses of message buffers rather than the message buffers themselves. The address of the message buffer to be accessed must be written to a location in memory, and a pointer to that location must be loaded into the appropriate UQCB.

If the UQCB can resolve the address of its associated QCB, there is no need for the NETWRKIF to open the queue using MP/M II Function 135 once the queue has been created. A pointer to the QCB can be placed in the UQCB at link time, instead. If, however, the QCB address cannot be resolved, an open queue operation must be performed. This might be the case if the system implementer breaks the NETWRKIF module into an RSP and a Banked Resident System Process (BRS).

Message Buffers

The message buffers must each be at least 262 bytes long, 5 bytes for the CP/NET header information, and 257 bytes for the actual CP/NET message. Even though the longest CP/NET message is only 256 bytes long, the extra byte is required because the server processes use the message buffer they are passed as a temporary scratch area.

If the data-link and network layers require additional header information, the message buffers must be even longer. If the message format used by the network is different from that used by CP/NET, the message must be converted into the standard CP/NET format before it is passed to the server process. The server process expects a one-byte format code of 0, a one-byte destination code equal to the server ID, a one-byte source code, a one-byte function code, a one-byte size code, and a contiguous message in binary format. The server returns an error for any deviation from this format.

A server process always returns its response to a requester in the same message buffer that it is passed. Consequently, no transport process should modify a message in between the time that its address is written to NtwrkQI<x> and the time that its address is read back from NtwrkQO<x>. To do so can cause the server to crash.

It is not always necessary to have one buffer for every server process in the server system. Fewer buffers can be provided if the network implementer limits the number of transactions that can occur simultaneously. It is important to recognize the distinction between the number of requesters supported (the number of sessions that can be ongoing at any one time) and the number of simultaneous transactions supported (the number of messages the server can process at any one time).

Because many server processes can be idle, the number of transactions can be much lower than the number of requesters. Limiting the number of transactions can sometimes drastically improve the performance of a CP/NET server because it reduces the amount of time the operating system switches from process to process trying to service a number of file-oriented requests simultaneously.

The Server Configuration Table

The server process must interface directly with a set of objects within the NETWRKIF to perform its own initialization, maintain its own reentrant processes, and perform validity checking on its incoming messages. These three sets of objects are the server configuration table, server Process Descriptor areas, and server process stacks.

The server configuration table is defined in Table 4-2.

Offset	Explanation
00-00	Server status byte. The communications software can use this byte to signal the current state of the network. This byte has no fixed function, however.
01-01	Server processor ID. The server processes compare this field against the destination ID field of all incoming messages. An error is returned if they do not match. A server ID of FF hex is illegal. Requester utility programs use a default server ID of 0, so a CP/NET network containing only one server identifies it as node 0, for convenience.
02-02	Maximum number of requesters supported at once. Up to 16 requesters can be supported.
03-03	Number of requesters currently logged in. This field is incremented by a server process when a login takes place and decremented when a logoff takes place. Logins return an error if the maximum equals the number currently logged in.
04-05	Log-in vector. Each bit of this field indicates whether the corresponding requester ID table entry is valid and refers to a logged-in requester. When a successful login takes place, a bit is set in this vector and the corresponding table entry is updated. When a logoff occurs, the table is searched and the corresponding bit is reset.

Offset	Explanation
06-21	Requester ID table. When a requester is successfully logged in, a server process locates an empty slot by checking the log-in vector, marks the slot as used, and then writes the source ID of the log-in message into this table, using the bit vector position as an index.
22-29	Log-in password. The password sent in the log in message must match this password, or the login fails, and an error is returned.

Table 4-2. Server Configuration Table

Just as the requester configuration table can be preconfigured to map certain devices as networked, the server configuration table can be preconfigured to define certain requesters as logged in without performing a login operation.

To do this, set the current number of logged-in requesters to the number of predefined logins desired. Make sure the number is less than the maximum number of requesters permitted. Otherwise, the server's behavior becomes unpredictable.

The log-in vector should have a bit set for every requester to be prelogged in, and the requester ID table should contain the logged-in requesters. For example, for a five-requester server where requesters 1, 2, and 5 are defined as already logged in, the server configuration table might look like this:

```

configtbl:    db 0          ; server status
              db 0          ; server ID
              db 5          ; max number of requesters
              db 3          ; currently logged in
              dw 8009h      ; log-in vector
              db 1          ; requester ID table
              ds 2
              db 2
              ds 11
              db 5
              db 'WUGGA'    ; password

```

The requester ID table is position independent. When a server process checks to see if a requester is logged in, it searches the entire requester table, using the entire log-in vector to check the entries for validity. Consequently, the configuration table is not sufficient to specify the process to which an incoming message should be routed.

The transport software must maintain its own routing mechanism. For example, the NETWRKIF in [Appendix E](#) maintains its routing implicitly as local data in its reentrant processes. The example in [Appendix F](#), on the other hand, relies on a requester control block that associates a source ID number with a UQCB.

Descriptors and Stacks

The module SERVER.RSP contains only one Process Descriptor and stack area. It is consequently initialized as only one process. SERVER.RSP must have some way of creating additional copies of itself. To do this, SERVER.RSP must know how many copies to create, and where to put the additional Process Descriptors and stacks.

By convention, the NETWRKIF process writes the address of the server configuration table into location offset 0009 in the system data page. The SERVER module uses this address to locate the maximum number of requesters from the configuration table. It then creates the maximum number, less

one, of processes. To locate storage to create the additional processes, the SERVER module expects to find stack areas for the extra processes directly following the configuration table.

Server process stacks must be exactly 150 bytes long, and there Must be one stack for each additional server. For example, to support a total of five servers, $4*150 = 600$ bytes of storage must he allocated after the configuration table.

The server expects the top of each additional server stack to contain a pointer to a 52-byte data area in which to create the new Process Descriptor. All of the Process Descriptor data areas must be contiguous.

Here is an example of the structure required for a four requester server:

```
server$pds:    ds (4-1)*52    ;server Process Descriptors
; (other data or code can be defined here)
configtbl:    ds 30          ;configuration table allocation
srvr$stkl:    ds 148         ;second server stack area
              dw server$pds
              ds 148         ;third server stack area
              dw server$pds+52
              ds 148         ;fourth server stack area
              dw server$pds+104
```

Listing 4-2. Stack and Process Descriptor Allocation for a Four-requester Server

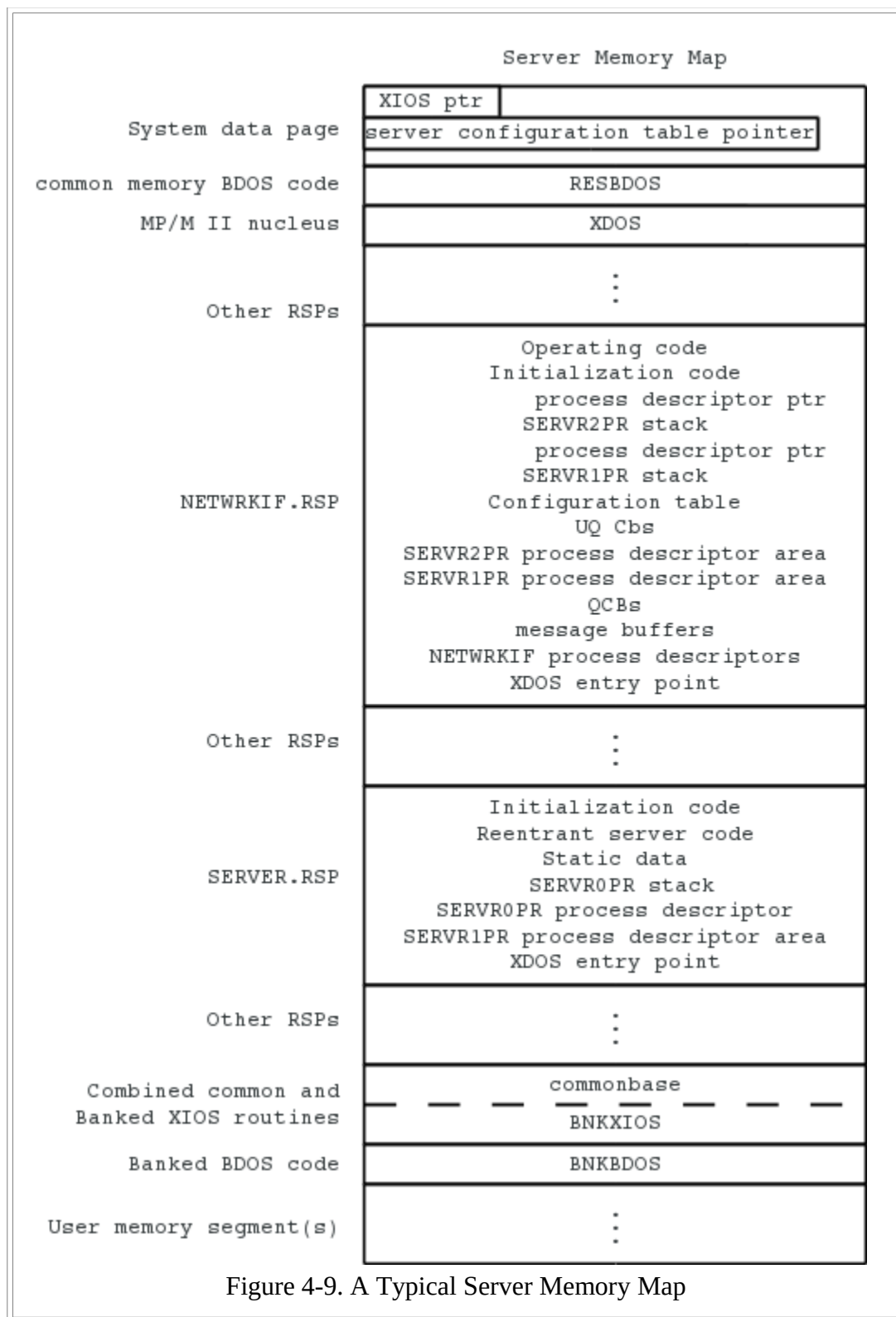
NETWRKIF Execution Requirements

The initialization code must perform the following actions:

- Initialize the network hardware, or cause lower-level routines to initialize it.
- Via MP/M II Function 134, make all input and output queues required to run the server.
- Write the address of the configuration table into the system data page.

These initialization functions need not be performed by a single process; they can be distributed among a variety of processes and interrupt service routines. The address of the configuration table should be written to the system data page with interrupts disabled. This prevents the server from loading an incorrect partial address and making its process-creation decisions on invalid data.

[Figure 4-9](#) shows a memory map, detailing how the SERVER.RSP and NETWRKIF.RSP modules fit into the rest of MP/M II, and how they communicate with one another during initialization.



Most of the other NETWRKIF run-time functions are discussed in previous sections. The general form of the NETWRKIF is the following:

1. Allocate a message buffer and receive a message. Check the message for data-link or network errors.
2. Reformat the message, if necessary, into the standard CP/NET format.
3. Compute the server process to which the message should be routed.
4. Write the message to the server's input queue.
5. Read the response from the server's output queue.

6. Send the response back to the requester, and free the buffer.
7. Repeat this process indefinitely.

4.3.4 Enhancements and Additions to the NETWRKIF

This section deals with extensions to the basic elements required to allow a CP/NET server to run under MP/M II. These extensions can increase the capabilities and improve the performance of the basic system.

Network Initialization and Maintenance

The network interface initialization can do much more than get the server processes ready to run. In addition to passing information about the network environment to the server and physical device initialization, the NETWRKIF can interrogate the network environment to identify other nodes in the system, their status, and their resources.

For example, the NETWRKIF network layer software might send out special packets to discover on-line nodes. When other NETWRKIFs and SNIOSs detect these packets, they respond with special routing packets of their own. If these routing messages are carefully designed, each node can build a table of routes to various nodes and mark other nodes as inaccessible.

Once the network has been initialized, a special network communications process intermittently circulates the routing packets. This circulation keeps the network routing information current as nodes go on and off line.

Nodes can be interrogated to identify their system resources for networking. For example, when a process similar to the routing process just described detects the existence of a node, it logs in to the node and sends out a series of dummy select disk messages. According to the error conditions returned, the process can identify the disk drives the node has available. This can also be accomplished by having a network-layer process issue its own select disk calls in response to receiving a special message.

In implementing these schemes, make sure these special messages do not interfere with regular CP/NET traffic. Some provisions are required to ensure that requests are not made to requesters that ignore the requests or mistake them for legitimate responses to previous requests. You might have to modify the SNIOS to allow it to deal with these strange messages.

Error Handling with Timeouts

Although the transport layer software of a CP/NET system is probably extremely reliable, and the possibility of garbled messages can be ignored, network data-link errors are likely in the long run. [Section 3.2.2](#) includes a general discussion of error handling. This section details a specific error-handling implementation, using timeouts.

Once the data-link software sends a message, it waits for an acknowledgment that the message was received. If no acknowledgment arrives, a timeout is triggered and the message is retransmitted.

You can implement a watchdog timeout mechanism as an interrupt service routine or as a process. When the transport process requests transmission from the data-link software, the process initializes a timeout variable and then waits on a flag. If the watchdog routine is implemented as an interrupt service routine (ISR), it decrements the timeout variable as a multiple of the clock interrupt frequency. If the watchdog routine is implemented as an extremely high priority process, it simply decrements the variable and then executes the MP/M II delay function for a fixed number of cycles.

With either method, a timeout status and the flag on which the transport process is waiting are set if the timeout variable is decremented to zero. At the same time, the data-link software sets the same flag and a transmission success status if it receives an acknowledgment.

When the transport process resumes processing after the flag wait operation, it checks the status variable to see which event occurred first. If the transmission timed out, the process attempts to retransmit. If the transmission succeeded, the transport process continues.

There are many variations to this method. The preceding one assumes that the message is transmitted with no handshake or initial signal to the receiver that a message is about to follow. If a handshake is implemented, it might require a timeout of its own. Several timeout points might have to be set throughout a single message, depending on how the receiver intends to acknowledge that message.

Other error conditions can occur; they can be integrated into the error-handling structure described above. For example, the receiver can transmit a negative acknowledgment, indicating that the message was received but that it was garbled. In this case, the data-link software need only set the same event flag, but instead of setting a message received status, it sets a transmit error variable. The transport process must now differentiate between three statuses rather than two when it resumes execution, but the overall structure is the same. The architecture required to implement timeouts is shown in [Figure 4-10](#).

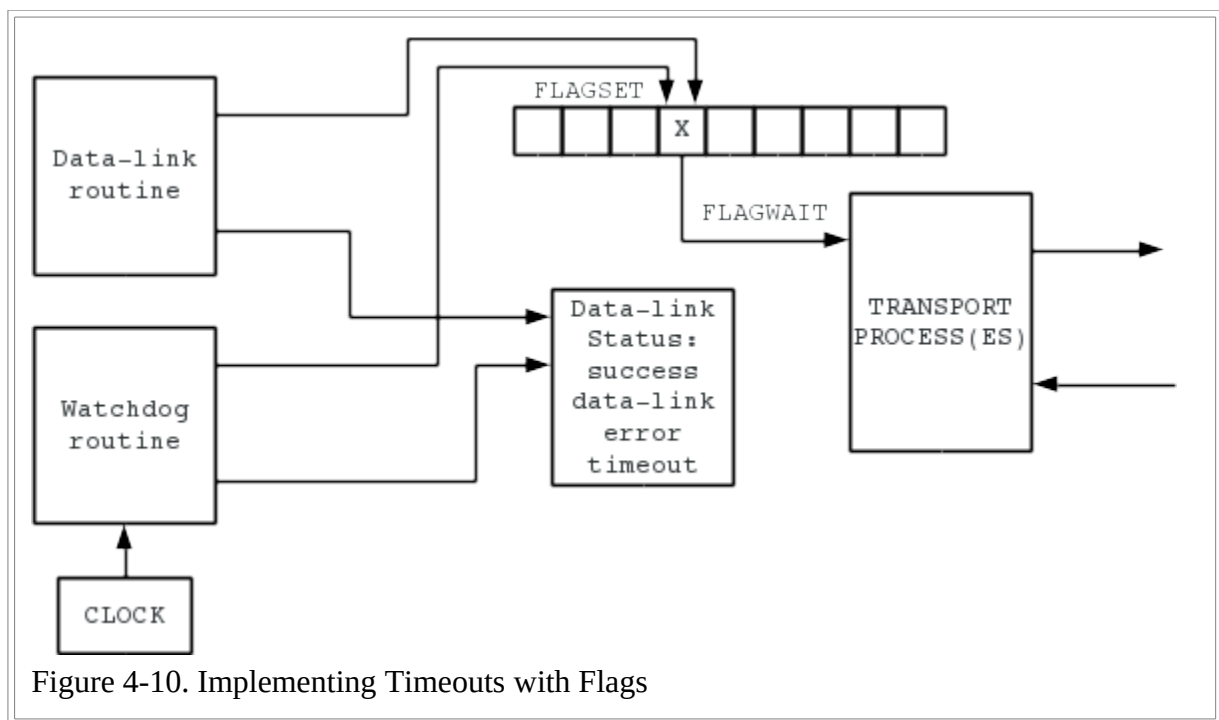


Figure 4-10. Implementing Timeouts with Flags

Store-and-Forward Networks

In some networks, the NETWRKIF can receive a message destined for another node that the sender could not reach directly. For these networks, implement network layer software to check the ultimate destination and send the message out along some other network line. These NETWRKIFs might need some of the following features.

The NETWRKIF might need more message buffers than there are supported requesters. Some messages are actually destined for the server processes resident on the current node, but a potentially high volume of the messages might be headed elsewhere.

The NETWRKIF must have a mechanism for receiving a message and then immediately sending it elsewhere without an intervening Queue Write-Queue Read operation. You can facilitate this type of operation by making the NETWRKIF software highly modular. It is advisable to have both network layer processes and transport layer processes, in addition to the data-link implementation you use. This gives the network layer process exclusive control of the data link layer, simplifying interprocess

competition for the data-link resource.

Finally, the network software must have a method of knowing which nodes can be reached through which network lines. This method can be a static, predefined table or a dynamic message-passing scheme like the one described in the preceding "Network Initialization and Maintenance" section.

Dynamic Login Handling

A CP/NET server under MP/M II can handle 16 requesters at a time. Many more physical requester nodes might want to access the server. The source ID byte in the standard CP/NET message allows up to 255 nodes. Theoretically, 254 requesters can be waiting to access one server.

Obviously, it would be useful to have a method whereby a server process can be reused by another requester after its previous owner has logged off. Unfortunately, the information contained in the server configuration table is not sufficient to identify which specific server processes are free and which are in use.

To solve this problem, define one requester control block (RCB) for each requester to be simultaneously supported by the server. The RCB is defined in [Table 4-3](#).

Offset	Explanation
00-00	Requester ID, If the control block is not in use, this field is set to FF hex.
01-03	Pointer to a particular server's input QCB.
04-05	A predefined pointer to byte 6 of this RCB.
06-07	A buffer that contains the address of the received message to be handled by this server process.

Table 4-3. Requester Control Block

Notice that this control block is a requester ID that can be matched with an incoming source ID, followed by a user queue control block. With this simple data structure, servers can be dynamically allocated to requesters with the following algorithm:

1. Receive a message.
2. Scan the RCBs for a match between the source ID of the message and the requester ID field of the RCB.
3. If a match is found, write the message buffer address into the RCB's message buffer address field in bytes 6 and 7. Then write to the queue, using the RCB's internal UQCB.
4. If a match is not found, but the scan reveals a free RCB (indicated by a requester ID field of FF), and the incoming message is a login, then flag the RCB in use by writing the message's source ID into the RCB; update the message buffer address field; and write to the queue.
5. If a match is not found and the message is not a login, send a message back to the requester with extended error 12, requester not logged in.
6. If a match is not found, and there are no free RCBs, and the message is a login, send a message back indicating login failed.
7. When a response message is read from the queue and the message is a logoff, then free the appropriate RCB before sending the message back to the requester.

This algorithm still does not allow more than 16 requesters to be logged in at the same time. But the algorithm does permit more than 16 requester nodes to compete for access to the server node. When more than 16 requester nodes log in, they receive login failed messages. These requester nodes cannot access the server until another requester logs off. In this kind of network it is advisable to implement an

automatic logoff feature for requesters that have not used the network for a fixed period of time.

Handling Special Messages

Special messages exchange network maintenance information between nodes. These messages have almost unlimited uses. For example, you can define a special message format for a special feature, high-performance print spooler. Once the format has been implemented, custom application packages can access it using Function 66 (Send Message on Network).

There are two basic steps to processing special message formats. First, the transport processes must be able to recognize special message formats and prevent them from entering the server processes. Second, the transport processes must have an interface to pass the messages off for special processing.

The first step can be accomplished by defining additional codes in the format field of the standard CP/NET message. When the transport software recognizes a strange format, it takes the appropriate action. If the message does not contain the standard CP/NET header, the data-link software can recognize this fact and notify the transport layer.

The problem of what to do with the message once it has been recognized can be solved using the same methods that are used for transporting messages throughout the more normal portions of the NETWRKIF. For example, the special print spooler and the transport layer can communicate via a predefined queue.

Some special formats require a logical response message. Functions 66 and 67 are intentionally exempt from the standard logical protocol of CP/NET. If a logical acknowledgment is required, then the transport layer must know how to accept it from the defined interface. Otherwise, the transport layer can forget that the special message occurred.

Bank-switched NETWRKIF Modules

Because of the size of the SERVER.RSP and NETWRKIF.RSP modules in a CP/NET server, MP/M II servers usually need more common memory than is available on the server system. Because of this, CP/NET users can use only one bank of their systems, completely wasting additional banks that might be used to run auxiliary processes or as additional disk buffer.

However, you can reduce the common memory requirements of an RSP by breaking it into two modules. One, still named a resident System process, contains only the code and data that must reside in Common memory to allow MP/M II to work. The rest of the module is reformatted and placed in a banked resident system process (BRS) that can be banked out when it is not executing, allowing its address space to be used by another process.

Process Descriptors and queue control blocks are the only sections of the server code that must reside in common memory. Prepare source module containing the XDOS entry point, all transport Process Descriptors, area for server Process Descriptors, all the NETWRKQIx QCBs, and all NTWRKQOx QCBs.

The first NETWRKIF Process Descriptor still must be allocated immediately after the XDOS entry point for the module, at relative address 0002H. However, this Descriptor's memory segment value should be 0FFH identifying that a BRS module is associated with it.

If any other processes exist in the NETWRKIF--for example, watchdog timeout processes--their Process Descriptors must also be included in this module. Assemble this source module and link it into RSP format. Name the object module <netprocess>.RSP where <netprocess> is the name of the first Process Descriptor in the module.

Then use the main body of the NETWRKIF source module to form a second source module. Remove all Process Descriptors and QCBs and place the following header at relative location 0:

```
rsp$adr:      ds 2          ;address of associated RSP
stk$adr:      dw stk$top    ;top of stack containing entry point
brs$name:     db '<netprocess>'
```

where `stk$top` is the address of the top of the stack for the first process, and `<netprocess>` matches the name of the associated RSP. This is the standard format for a BRS module; it is described in more detail in the *MP/M II Operating System System Guide*.

Because the Process Descriptors and queue control blocks are in a completely separate RSP, they cannot be resolved as simple externals. They must be defined in terms of known offsets from the beginning of `<netprocess>.RSP`. At run-time, the variable `rsp$adr` contains a pointer to the beginning of this RSP, placed there by MP/M II's GENSYS utility. Using this pointer and the predefined offsets, required references to these data objects can be resolved.

On startup, the NETWRKIF processes perform the following initialization:

1. Initialize the stack pointer fields in all NETWRKIF Process Descriptors with a pointer to the top of the stack associated with each process. This is not necessary for the first process because GENSYS provides the stack pointer linkage via the header data in the BRS.
2. The make queue operations the NETWRKIF requires can be complicated because the QCB addresses must be resolved. Once they are, however, update the UQCBs associated with them with those addresses, avoiding the necessity of performing open queue functions.

The NETWRKIF.BRS module requires a different way of referencing the operating system because it does not contain a pointer to the XDOS entry point. The RSP associated with the BRS module, however, does contain such a pointer as its first two bytes. The following subroutine performs operating system calls transparently:

```
do$os:  lhld rsp$adr
        mov a,m
        inx h
        mov h,m
        mov l,a
        pchl
```

you must also assemble this module and link it into RSP format; but name it `<netprocess>.BRS`.

Banking out the NETWRKIF module alone might raise the BNKXIOS COMMONBASE entry point above the hardware bank-select point, allowing banked operation of MP/M II. If banking out the module does not accommodate this, you can use a patch to convert SERVER.RSP into a banked module in a similar way. The patch is detailed in *CP/NET V1.2 Application Note #2, 11/11/82*.

Perform GENSYS with a specified banked system. You can add memory segments to occupy the new banks. The address ranges of the new memory segments are prompted for at the end of GENSYS.

If the number of requesters to be supported still requires more common memory than is available, there is no purpose in implementing a banked version of the server.

A banked-out server has a marginally slower response time because the dispatcher must select the system bank and because of the added level of indirection in calling the operating system. This degradation, however, is insignificant.

Although banking out the server provides additional user Segments under MP/M II, resist the temptation to add additional consoles to the system. Because of the extremely high priority at which the server runs, performance on additional consoles is very poor. However, these extra banks do provide the user with a means of performing occasional jobs directly from the MP/M II level. more importantly, extra segments can enhance the server itself by using special CP/NET messages.

4.3.5 MP/M II Performance Factors Affecting the NETWRKIF

The characteristics of the network for which a server is being implemented influence the architecture of the NETWRKIF and the rest of the server software. Another important factor in designing efficient servers under MP/M II is the nature of MP/M II itself. This section points out the overhead MP/M II incurs in implementing multitasking programming environment.

The heart of the MP/M II operating system is its dispatcher. This routine is entered every time a system call is made. The dispatcher protects system resources, tests for events that could influence the execution of any process in the system, and finally chooses the processes to execute and their order. The dispatcher takes roughly 900 microseconds to execute, but interrupts are disabled for no longer than 90 microseconds. This overhead is incurred on every system call.

The limitations of the dispatcher alone place some basic constraints on communications speed. If the network is using a serial I/O device capable of buffering three characters at 10 bits per character, then the NETWRKIF had better not rely on a system call like console input to receive network messages if the transmission rate is faster than 33K bits per second and the sender sends characters as fast as possible. Even below this speed, overruns are likely if there are any other processes in the system. This assumes an extremely simple protocol. If the network has extra signal lines, most serial I/O devices permit the receiver to signal a clear to send condition back to the sender. But networks often must work without these extra signals.

Because interrupts are disabled for no longer than 90 microseconds, a network that works at the character-interrupt level functions properly at transmission speeds up to 333K bits per second. Beyond that speed, overruns are likely to occur too often for adequate performance.

At speeds higher than 333K bits per second, the network interface software can use one of three approaches:

- A process can disable interrupts and perform no system calls, preventing the dispatcher from being entered, and perform its own direct network I/O, character by character.
- The network interface can use DMA to transfer large blocks of message data and perform validity checking after the message has been transferred.
- The network interface can use an intelligent protocol controller that also does DMA or it can map completed messages from its own memory space into MP/M II's memory space.

Serial I/O is not the only possible network transmission medium. The example is provided to acquaint you with the performance of MP/M II.

The amount of time spent in the dispatcher varies depending on a number of factors. Because the dispatcher must check suspended processes against system events, keep the number of processes, queues, flags, and poll calls to a minimum. Poll calls are especially degrading. Every time the dispatcher is entered, it executes every code fragment associated with every outstanding poll call. If all 16 reentrant NETWRKIF processes polled output ports at once, the dispatcher would be very busy. In fact, enough poll calls can lengthen the dispatcher's execution time so much that it exceeds the clock interval. When this happens, the dispatcher is reentered before it has even been exited.

The design of interrupt service routines Must take the structure of the dispatcher into account. ISRs must first

of all save the register image of the process they interrupted--the service routine then executes. When the ISR terminates itself, it should restore the interrupted process's registers and take one of two actions:

- If the service routine winds up setting a flag, the flag set call to MP/M II should be made, followed by a jump into the dispatcher. This allows the dispatcher to ready the process waiting on the flag as quickly as possible.
- If no flag is to be set, the ISR can simply return to the interrupted process.

ISRs should perform no MP/M II system calls except for the Flag Set function. There are two reasons for this. First, ISRs are not processes, so the dispatcher has no way of saving the status of the ISR in a Process Descriptor before allowing the function to be performed. Second, the dispatcher reenables interrupts and possibly dispatches another process, leaving the ISR and the interrupted process in an indeterminate state. The Flag Set function is specifically recognized by the dispatcher to avoid dispatching a different process.

Several factors determine how often the NETWRKIF and server processes are dispatched. The most obvious is, once again, the number of processes. If MP/M II must share the CPU with more tasks, there is less CPU available. Consider the priority of the various network server processes carefully. All processes in the SERVER module run at a high priority level of 100. processes in the NETWRKIF might require other priorities. In general, assign compute-bound processes lower priorities than I/O-bound processes, to prevent processes that perform few system calls from hogging the CPU.

The dispatcher always schedules processes according to priority. Improperly setting priorities can cause processes to be permanently suspended. For example, consider a NETWRKIF module that performs all direct I/O and busy-waits for network input. Suppose this process has a priority of 60, slightly higher than the server processes. Although the dispatcher is entered every time the system clock ticks, the NETWRKIF is ready. Because the NETWRKIF has a higher priority than the server processes, the server processes never execute.

Note that because of the extremely high priority of the server process, normal user processes running under MP/M II perform very poorly. In addition, the extra process load degrades the server performance. It is recommended that normal work station terminals not be provided on an MP/M II system that is functioning as a server, although a system console can be convenient for monitoring system performance and giving the operator a means of maintaining the server's data base.

The last factor affecting the dispatch rate is the system clock frequency. Every time a clock tick occurs, the dispatcher is entered and recomputes the process to be executed next. Processes of equal priority are dispatched on a first come, first served basis. The system clock can be tuned for optimal network performance. There are no general rules on tuning because each network and the applications run on the network determine the optimal clock period. Experiment with the clock frequency to determine the best performance for the server.

In addition to designing the NETWRKIF for the server system, you might want to reexamine the XIOS used in the system. Many CP/NET users discover that once their communications system has been optimized, server performance has improved only slightly because several requesters are forcing the disk system to thrash.

Thrashing can be minimized if the XIOS is provided with efficient blocking/deblocking algorithms like those discussed in the *MP/M II Operating System System Guide*. These algorithms buffer disk accesses, deferring physical Read-Write operations until they are absolutely necessary. As a result, many file record Read-Write operations occur at memory speed, instead of having to wait for physical I/O from a disk drive.

Extra blocking/deblocking buffers can also improve overall server performance enormously. Because a

dedicated server only requires a single tiny user program segment, or, in some cases, no user segment at all, almost all additional memory remaining after the server has been implemented can be used for disk buffers. In a bank-switched or memory-managed system, potentially huge amounts of memory can be made available for disk buffers. Providing one or more disk buffers per supported requester potentially eliminates competition between two requesters for buffer resources.

Another way to improve disk performance with limited memory for disk buffers is to limit the number of transactions that can be present in the server at one time. Even if a server is supporting 16 requester sessions, it is possible, for example, to permit only four or five messages to be active at a time. This limit reduces the amount of competition between actual processes, although competition continues between individual transactions. Quite often, however, the overhead incurred by refusing network messages and forcing requesters to retransmit them is considerably less than the overhead incurred by repeatedly having to flush disk buffers for use and reuse by individual processes.

You can estimate the average number of disk accesses an application program is likely to perform in a short time. The NETWRKIF processes can then selectively transport messages from only one requester for a short amount of time, then service another requester for an equal amount of time. The scheme allows a single process to take maximum advantage of the blocking and deblocking algorithms implemented in the server's XIOS. The major disadvantage of such a scheme is that it is extremely complex and difficult to implement efficiently. Carefully tuned, however, it can greatly improve performance.

4.3.6 Generating the NETWRKIF

To create the MP/M II server, perform the following steps:

1. If the XIOS has been modified, generate a new version of RESXIOS.SPR or BNKXIOS.SPR or BNKXIOS.SPR, according to the instructions provided in the *MP/M II Operating System System Guide*.
2. Assemble and link the NETWRKIF module:

```
A>RMAC NETWRKIF
A>LINK NETWRKIF[NR,OR]
```

The linker generates the NETWRKIF.RSP file.

If RMAC and LINK are not available, you must use ASM, PIP, and GENMOD, as shown below:

Assemble with ORG 0000H.

```
A>ASM NETWORKIF
A>REN NTWRK0.HEX=NETWORKIF.HEX
```

Now edit the NETWRKIF.ASM ORG statement to locate the module at 100 hex. Assemble with ORG 0100H.

```
A>ASM NETWRKIF
A>REN NTWRK1.HEX=NETWRKIF.HEX
```

Concatenate the HEX files.

```
A>PIP NETWRKIF.HEX=NTWRK0.HEX,NTWRK1.HEX
```

Generate the NETWRKIF RSP file.

```
A>GENMOD NETWRKIF.HEX NETWRKIF.RSP
```

3. Copy the following files to the server boot disk.

- SERVER.RSP = Server process Module
- NETWRKIF.RSP = Custom Network Interface Process
- MAIL.COM = Mail Utility

4. Perform a GENSYS on the MP/M II system. The GENSYS must include the SERVER.RSP file and the customized NETWRKIF.RSP; it can also include the SPOOL.RSP.

When GENSYS asks for the number of consoles, do not include the consoles (character I/O drivers) that support the requesters. Usually, the response is 1.

You must also configure the file system for the types of applications CP/NET runs, enable compatibility attributes, if necessary, and so on. These issues are discussed in the *MP/M II Operating System System Guide*.

4.3.7 Debugging the NETWRKIF

The MP/M II server is now ready to be debugged. There are three general strategies for debugging the server.

Debugging MP/M II Under CP/M

To debug MP/M II under CP/M, follow these steps:

1. GENSYS the MP/M II with the top of memory set below where a CP/M system running on the same hardware would reside when it is running DDT, SID, or ZSID.
2. Boot up CP/M on the server target computer system.
3. Run MPMLDR under the debugger. You can halt the loader just before passing control to MP/M II through the following sequence:

```
A>DDT MPMLDR.COM
*I$B
*G
```

When the loader breaks, you can insert breakpoints and restart the loader.

When using this method, remember that, because CP/M is a single-tasking operating system, the entire CP/M operating system becomes part of the process in which a breakpoint is inserted every time the system encounters a breakpoint. Furthermore, DDT and SID reenables interrupts on breakpoints. If a clock tick goes off, the MP/M II dispatcher is likely to suspend CP/M and continue with other processing. This might not inconvenience you because the process that was breakpointed is also suspended. If it does affect the operation of the system, you might have to disable the system clock.

Debugging the NETWRKIF as a COM file

The example in [Appendix E](#) is set up to debug the NETWRKIF as a COM file. Debugging instructions are also included in [Appendix E](#).

Inserting Trace Code Into the NETWRKIF

Gather run-time statistics by inserting trace code into the NETWRKIF. Although this is not very helpful for debugging real-time problems, it is the least destructive method of gathering real-time

statistics. This method can also be useful when tuning the network for increased performance.

4.4 Implementing Non-MP/M II Servers

It is possible to implement a CP/NET server on any computer system, under any operating system. There are several reasons why you might choose another operating system:

- MP/M II servers limit the number of requesters to 16. You might want more than 16 work stations to have access to a common database.
- You might require higher performance levels. The high speed of a mainframe CPU can substantially increase CP/NET performance.
- You might want your system to take advantage of the large base of CP/M applications programs, but maintain its files under another operating system. Or you might want to create a gateway to one of the other commercially available network systems. A special server could translate CP/NET messages into an appropriate format for the other network.

The module SERVER.RSP cannot be used on a different processor or under a different operating system. So you must not only create the equivalent of the NETWRKIF for the target computer system; you must also write the logical portion of the server.

The server processes under MP/M II act essentially as a proxy for the requester assigned to them. For example, the requester wants to open a file on a networked drive but it does not have access to the operating system controlling that drive. Instead, the requester sends a message to a server process that does have direct access to the controlling operating system and asks that process to open the file for the requester. The server obligingly performs the operation for the requester and tells it what happened. This is often referred to as a ghosted process model of a server because the operating system thinks it is running the entire application program as a process, while in fact the application is running somewhere else, but has a friend to help out.

Using the logical messages included in this manual, you can write a ghosted process server for CP/NET under almost any multitasking operating system. You can even write a CP/NET server under a single-tasking operating system. (CP/NET servers have actually been implemented under CP/M.)

The basic elements of such a server are

- A communications interface.
- A function interpreter. This module must interpret the logical messages sent by the CP/NET requester and take the appropriate action.
- A file system translator. This module must convert CP/M BDOS File Control Blocks passed by the requester into native operating system File Control Blocks.
- An operating system interface. This module must translate a CP/NET function that corresponds exactly to a function supported by MP/M II into a function or set of functions supported by the native operating system.

Each of these functional modules varies depending on the environment under which it is forced to execute. The communications interface is governed by the types of process architectures the target operating system can support. The remaining modules can be a set of reentrant processes, as they are under MP/M II, or they can be a single process that keeps track of the requester it is currently servicing. If the latter method is used, the server must keep track of such context sensitive information as directory search first/search next information and shared files.

It might not be possible to support all CP/M functions under a non-MP/M II server. If this is the case, choose applications that do not require the use of the unsupported functions.

Finally, it might be necessary to have several different computer systems and operating systems acting as servers in the same network. It is best to make the server implementation as portable as possible. Implementing the server in a high-level language is a first step to portability.

Making the system highly modular can improve its portability. For example, break the communications interface into a hardware interface module, a data link module, a network module, and a transport module. All of these modules, with the exception of the hardware interface, can port to different systems with minimal modification.

The server's function interpreter should be completely portable, but you will probably have to rewrite the file system interpreter and the operating system interface modules.

Appendix A

CP/NOS Overview

A.1 overview

CP/NOS is a version of the CP/M operating system that performs all file handling across a CP/NET network system. CP/NOS supports one local console and one local printer, but it supports only remote mass storage media. Because of this, the BDOS and BIOS modules in a CP/NOS system are considerably smaller than their counterparts in a standard CP/M system. This allows CP/NOS to fit in a fairly small (usually 4K bytes) Read-Only memory, so you do not need a bootstrap loader. CP/NOS can also be downloaded from a server. Using a small loader, you can also download a CP/NOS system from a centralized server.

Programs written under any CP/M 2.x system are fully compatible with a comparable CP/NOS system, provided that mass storage devices referenced by the application are available across the network. When BDOS calls that service, these devices are automatically translated into network functions.

Unlike CP/NET, CP/NOS cannot be loaded under an existing CP/M system. The network modules and CP/M modules must be linked together and executed in a stand-alone environment. The special problems this creates in debugging CP/NOS are discussed in this appendix.

A.2 System Requirements

CP/NOS can run on an 8080, 8085, or Z80 microprocessor, with a maximum of 64K of memory. A usual CP/NOS system can be placed in a 4K ROM.

The CP/NOS requester must be networked to an MP/M II server. The server is the same as the one used by CP/NET. CP/NOS and CP/NET requesters can even be networked to the same server.

A.3 Customizing CP/NOS

Three of the modules incorporated in CP/NOS are system dependent and must be modified to work on a particular hardware configuration. They are the CPBIOS, CPNIOS, and NETWRKIF modules

The CPBIOS can be exactly the same as the BIOS used in a CP/M system that runs on the same hardware, except that only a small portion of the BIOS is required. The only routines required are:

BOOT	cold start
CONST	read console status
CONIN	read console character
CONOUT	write console character
LIST	write character to the list device
LISTST	read list device status

The CPBIOS jump vector must be the same as that of a regular BIOS, but all other entry points can be null.

The CPNIOS module takes the place of the SNIOS module in CP/NET and requires only minimal modification. The only difference is that all variables must be initialized upon cold start, including the requester configuration table. The utilities NETWORK and LOGIN are not sufficient to define the configuration table after cold start because CP/NOS has no local disk drives from which to load these utilities. The CPNIOS must also prompt the user for login information upon cold start, or a warm boot results

in continuous requester not logged in extended errors as the CP/NOS requester tries to load the file CCP.SPR from a server that has no knowledge of the requester.

The SNIOS example in Appendix E contains a sample CPNIOS, conditionally assembled out. To obtain the CPNIOS version, equate the literal CPNOS to true.

Note: if the two preceding routines are to reside eventually in ROM, all variable data must be contained in data segments and cannot be initialized at run-time. Initializing values must reside in a code segment, and they must be copied down to their corresponding data segment locations at cold start. The assembly of these modules requires an assembler capable of supporting separate code and data segments; the segments must be assembled into REL file format. Use RMAC with 8080 source files.

The NETWRKIF module resides on the server and is identical to the NETWRKIF required to support CP/NET. See [Section 4.3](#) for a discussion of NETWRKIF preparation.

A.4 Building the CP/NOS System

To generate a CP/NOS system ready for insertion into ROM, follow these steps:

1. Assemble the modules CPBIOS and CPNIOS.
2. Link the following modules together in the order shown, using LINK-80:

CPNOS, CPNDOS, CPNIOS, CPBDOS, CPBIOS

Locate the code segment where the ROM sits in the address space of the finished system. At least 1K (400 hexadecimal bytes) of RAM must be allocated for data segments. If the code segments are to be loaded into high memory (at F000H for a 4K system), data must be explicitly linked, using the D option, at least 1K in front of the code segments. For example,

```
A>LINK CPNOS,CPNDOS,CPNIOS,CPBDOS,CPBIOS[LF000,DEC00]
```

These two steps produce an executable CP/NOS, capable of being programmed into ROM. At this stage, however, the system cannot be debugged from CP/M.

A.5 Debugging the System

You can create a version of CP/NOS that can be cold started from CP/M if a CP/M system with 64K RAM is available. First, type the following commands:

```
A>RMAC CPNIOS
A>RMAC CPBIOS
A>LINK CPNOS,CPNDOS,CPNIOS,CPBDOS,CPBIOS[LF000,DEC00]
A>GENHEX MVCPNOS 0100
A>GENHEX CPNOS 0200
A>PIP LDCPNOS.HEX=MVCPNOS.HEX[I],CPNOS.HEX[H]
A>LOAD LDCPNOS
```

This procedure produces a file LDCPNOS.COM that is directly executable from CP/M. LDCPNOS relocates the CPNOS module to location F000H and passes control to it, destroying CP/M and replacing it with CP/NOS.

Because CP/M is destroyed by this procedure, it is not advisable to run LDCPNOS under software debugger like DDT or SID, although you can run LDCPNOS under an in-circuit emulator. To run CP/NOS under DDT

or SID, use the following procedure:

1. Link CPNOS so that all code and data reside below the address specified as END when the debugger is brought up:

```
A>LINK CPNOS,CPNDOS,CPNIOS,CPBDOS,CPBIOS[L<org>,D<org-400H>]
```

where <org> is the link origin.

2. A>DDT CPNOS.COM
3. Relocate CPNOS from location 100, where DDT loads it, to its link origin:

```
-M100,<100+next-1>,<org>
```

where next is the field specified by NEXT when the debugger loads CPNOS.COM, and <org> is the link origin.

4. Begin execution with appropriate diagnostics:

```
-G<org>
```

where <org> is the link origin.

Appendix B
CP/NET 1.2 Standard Message Formats

FMT	DID	SID	FNC	SIZ	MSG
<ul style="list-style-type: none"> • FMT = Message format code • DID = Message destination processor ID • SID = Message source processor ID • FNC = MP/M function code • SIZ = Data field length - 1 • MSG = Actual message, SIZ + 1 bytes long 					

Figure B-1. CP/NET 1.2 Logical Message Format

FMT CODE	FMT	DID	SID	FNC	siz	MSG	Comment
00	1	1	1	1	1	1-256	Preferred format
01	1	1	1	1	1	1-256	Returned result
02	1	1	1	1	2	1-65536	
03	1	1	1	1	2	1-65536	Returned result
04	1	2	2	1	1	1-256	
05	1	2	2	1	1	1-256	Returned result
06	1	2	2	1	2	1-65536	
07	1	2	2	1	2	1-65536	Returned result

Table B-1. Message Field Length Table

Appendix C

CP/NET 1.2 Logical Message Specifications

Messages for all CP/NET functions are defined in this appendix. These messages are logical messages. Any implementation of the SNIOS or NETWRKIF modules must always present messages to the NDOS or SERVER modules in the form presented here.

You must adhere to these formats when implementing a server that runs under an operating system other than MP/M II.

Notes:

- ss = Server ID
- rr = Requester ID
- xx = Don't care byte
- nn = Value specified
- All numeric values are in hexadecimal.
- All functions capable of returning extended errors are marked *EE*. Extended errors are assumed w bytes in length, with the following message format:

```
SIZ = 01
MSG(0) = FF
MSG(1) = Extended Error Code
```

Any response with SIZ = 01 is interpreted as an error, regardless of the value in MSG(0).

- Any message can return the *server not logged in* or *function not implemented on server* extended error

For functions that return with the user's FCB updated (messages that have an FCB in their response), the first byte of the FCB (drive designator) is never copied back from the response message. In some cases, the random record bytes are also not copied back.

For search functions, the entire directory entry (which is NOT an FCB) is copied back to the current DMA buffer, into the position indicated by the Directory Code result byte. This means that the DMA buffer is not the actual directory sector from the disk, but merely an accumulation of directory entries in an order determined by how they were found.

Functions that return the address of a system resource (Get Allocation Vector, Get DBP, Get Server Config), the data is kept in an NDOS buffer which is overwritten on subsequent calls. The user must copy data out as needed. The NDOS guarantees that the low byte of the address is never 0FFh, so that a valid address can be distinguished from an Extended Error Code.

For functions that return a drive vector (Get Login, Get R/O), there is no way to distinguish between valid vectors and errors. The message format does not allow for errors, and the NDOS ignores the possibility of an error in the response. The actual vector returned to the user is a composite of data retrieved from all known servers and the local BDOS.

FMT	DID	SID	FNC	SIZ	MSG
0 SYSTEM RESET: NOT IMPLEMENTED AT SERVER					
00	ss	rr	00	00	• 00-00 = xx
01	rr	ss	00	01	• 00-00 = 0FFh • 01-01 = 00Ch

FMT	DID	SID	FNC	SIZ	MSG
1 CONSOLE INPUT: NOT IMPLEMENTED AT SERVER					
00	ss	rr	01	00	• 00-00 = xx
01	rr	ss	01	01	• 00-00 = 0FFh • 01-01 = 00Ch
2 CONSOLE OUTPUT: NOT IMPLEMENTED AT SERVER					
00	ss	rr	02	00	• 00-00 = xx
01	rr	ss	02	01	• 00-00 = 0FFh • 01-01 = 00Ch
3 RAW CONSOLE INPUT:					
00	ss	rr	03	00	• 00-00 = Server Console #
01	rr	ss	03	00	• 00-00 = Character Input
4 RAW CONSOLE OUTPUT:					
00	ss	rr	04	01	• 00-00 = Server Console # • 01-01 = Character to Output
01	rr	ss	04	00	• 00-00 = 00
5 LIST OUTPUT:					
00	ss	rr	05	nn	• 00-00 = Server List # • 01-nn = Characters to List Device (nn = 01 to 80)
01	rr	ss	05	00	• 00-00 = 00
6 DIRECT CONSOLE I/O: NOT IMPLEMENTED AT SERVER					
00	ss	rr	06	00	• 00-00 = xx
01	rr	ss	06	01	• 00-00 = 0FFh • 01-01 = 00Ch
7 GET I/O BYTE: NOT IMPLEMENTED AT SERVER					
00	ss	rr	07	00	• 00-00 = xx
01	rr	ss	07	01	• 00-00 = 0FFh • 01-01 = 00Ch
8 SET I/O BYTE: NOT IMPLEMENTED AT SERVER					
00	ss	rr	08	00	• 00-00 = xx
01	rr	ss	08	01	• 00-00 = 0FFh • 01-01 = 00Ch
9 PRINT STRING: NOT IMPLEMENTED AT SERVER					

FMT	DID	SID	FNC	SIZ	MSG
00	ss	rr	09	00	• 00-00 = xx
01	rr	ss	09	01	• 00-00 = 0FFh • 01-01 = 00Ch
10 READ CONSOLE BUFFER: NOT IMPLEMENTED AT SERVER					
00	ss	rr	0A	00	• 00-00 = xx
01	rr	ss	0A	01	• 00-00 = 0FFh • 01-01 = 00Ch
11 GET CONSOLE STATUS:					
00	ss	rr	0B	00	• 00-00 = Server Console #
01	rr	ss	0B	00	• 00-00 = Console Status Byte
12 RETURN VERSION NUMBER: NOT IMPLEMENTED AT SERVER					
00	ss	rr	0C	00	• 00-00 = xx
01	rr	ss	0C	01	• 00-00 = 0FFh • 01-01 = 00Ch
13 RESET DISK SYSTEM: NOT IMPLEMENTED AT SERVER					
00	ss	rr	0D	00	• 00-00 = xx
01	rr	ss	0D	01	• 00-00 = 0FFh • 01-01 = 00Ch
14 SELECT DISK: *EE*					
00	ss	rr	0E	00	• 00-00 = Selected Disk
01	rr	ss	0E	00	• 00-00 = Return Code
15 OPEN FILE: *EE*					
00	ss	rr	0F	2C	• 00-00 = User Number • 01-24 = FCB • 25-2C = Password
01	rr	ss	0F	24	• 00-00 = Directory Code • 01-24 = FCB 01 not copied to user; 22-24 (file ID) not copied to user unless F5' set and not F6'.
16 CLOSE FILE: *EE*					
00	ss	rr	10	2C	• 00-00 = User Number • 01-24 = FCB • 25-2C = Password (ignored)
01	rr	ss	10	24	• 00-00 = Directory Code • 01-24 = FCB 01 not copied to user; 22-24 (file ID) not copied to user unless F5' set and not F6'.

FMT	DID	SID	FNC	SIZ	MSG
17 SEARCH FOR FIRST: *EE*					
00	ss	rr	11	25	<ul style="list-style-type: none"> • 00-00 = Current Disk if FCB(0)='? • 01-01 = User Number • 02-25 = FCB
01	rr	ss	11	20	<ul style="list-style-type: none"> • 00-00 = Directory Code • 01-20 = Directory Entry Directory Entry copied to current DMA buffer based on Directory Code.
				80/0	<ul style="list-style-type: none"> • 00-00 = Directory Code • 01-80 = Directory Sector • 00-00 = Directory Code Extensions for "full search" mode, incl. CP/M 3.
18 SEARCH FOR NEXT: *EE*					
00	ss	rr	12	01	<ul style="list-style-type: none"> • 00-00 = xx • 01-01 = User Number
01	rr	ss	12	20	<ul style="list-style-type: none"> • 00-00 = Directory Code • 01-20 = Directory Entry Directory Entry copied to current DMA buffer based on Directory Code.
				80/0	<ul style="list-style-type: none"> • 00-00 = Directory Code • 01-80 = Directory Sector • 00-00 = Directory Code Extensions for "full search" mode, incl. CP/M 3.
19 DELETE FILE: *EE*					
00	ss	rr	13	24	<ul style="list-style-type: none"> • 00-00 = User Number • 01-24 = FCB
01	rr	ss	13	00	<ul style="list-style-type: none"> • 00-00 = Directory Code
20 READ SEQUENTIAL: *EE*					
00	ss	rr	14	24	<ul style="list-style-type: none"> • 00-00 = User Number • 01-24 = FCB
01	rr	ss	14	A4	<ul style="list-style-type: none"> • 00-00 = Return Code • 01-24 = FCB • 25-A4 = Sector of Data Read 01 and 22-24 not copied to user.
21 WRITE SEQUENTIAL: *EE*					
00	ss	rr	15	A4	<ul style="list-style-type: none"> • 00-00 = User Number • 01-24 = FCB • 25-A4 = Sector of Data to Write
01	rr	ss	15	24	<ul style="list-style-type: none"> • 00-00 = Return Code • 01-24 = FCB 01 and 22-24 not copied to user.
22 MAKE FILE: *EE*					

FMT	DID	SID	FNC	SIZ	MSG
00	ss	rr	16	24	<ul style="list-style-type: none"> • 00-00 = User Number • 01-24 = FCB
01	rr	ss	16	24	<ul style="list-style-type: none"> • 00-00 = Directory Code • 01-24 = FCB 01 not copied to user; 22-24 (file ID) not copied to user unless F5' set and not F6'.
23 RENAME FILE: *EE*					
00	ss	rr	17	24	<ul style="list-style-type: none"> • 00-00 = User Number • 01-24 = FCB in RENAME format
01	rr	ss	17	00	<ul style="list-style-type: none"> • 00-00 = Directory Code
24 RETURN LOGIN VECTOR: Message Sent for Each Remote Drive, Results combined with local BDOS					
00	ss	rr	18	00	<ul style="list-style-type: none"> • 00-00 = xx
01	rr	ss	18	01	<ul style="list-style-type: none"> • 00-01 = Login Vector
25 RETURN CURRENT DISK: NOT IMPLEMENTED AT SERVER					
00	ss	rr	19	00	<ul style="list-style-type: none"> • 00-00 = xx
01	rr	ss	19	01	<ul style="list-style-type: none"> • 00-00 = 0FFh • 01-01 = 00Ch
26 SET DMA ADDRESS: NOT IMPLEMENTED AT SERVER					
00	ss	rr	1A	00	<ul style="list-style-type: none"> • 00-00 = xx
01	rr	ss	1A	01	<ul style="list-style-type: none"> • 00-00 = 0FFh • 01-01 = 00Ch
27 GET ALLOCATION VECTOR ADDRESS: *EE*					
00	ss	rr	1B	00	<ul style="list-style-type: none"> • 00-00 = Current Disk
01	rr	ss	1B	02-FF	<ul style="list-style-type: none"> • 00-FF = Allocation Vector NDOS guarantees low byte of returned address cannot be FF.
28 WRITE PROTECT DISK: *EE*					
00	ss	rr	1C	00	<ul style="list-style-type: none"> • 00-00 = Current Disk
01	rr	ss	1C	00	<ul style="list-style-type: none"> • 00-00 = 00
29 GET R/O VECTOR: Message Sent for Each Remote Drive, Results combined with local BDOS					
00	ss	rr	1D	00	<ul style="list-style-type: none"> • 00-00 = xx
01	rr	ss	1D	01	<ul style="list-style-type: none"> • 00-01 = R/O Vector
30 SET FILE ATTRIBUTES: *EE*					
00	ss	rr	1E	24	<ul style="list-style-type: none"> • 00-00 = User Number • 01-24 = FCB with File Attributes Set

FMT	DID	SID	FNC	SIZ	MSG
01	rr	ss	1E	00	• 00-00 = Directory Code
31 GET DISK PARAMETER ADDRESS: *EE*					
00	ss	rr	1F	00	• 00-00 = Current Disk
01	rr	ss	1F	0F	• 00-0F = Disk Parameter Block
32 SET/GET USER CODE: NOT IMPLEMENTED AT SERVER					
00	ss	rr	20	00	• 00-00 = xx
01	rr	ss	20	01	• 00-00 = 0FFh • 01-01 = 00Ch
33 READ RANDOM: *EE*					
00	ss	rr	21	24	• 00-00 = User Number • 01-24 = FCB
01	rr	ss	21	A4	• 00-00 = Return Code • 01-24 = FCB • 25-A4 = Sector of Data Read 01 not copied to user.
34 WRITE RANDOM: *EE*					
00	ss	rr	22	A4	• 00-00 = User Number • 01-24 = FCB • 25-A4 = Sector of Data to Write
01	rr	ss	22	24	• 00-00 = Return Code • 01-24 = FCB 01 not copied to user.
35 COMPUTE FILE SIZE: *EE*					
00	ss	rr	23	24	• 00-00 = User Number • 01-24 = FCB
01	rr	ss	23	24	• 00-00 = Return Code • 01-24 = FCB
36 SET RANDOM RECORD: *EE*					
00	ss	rr	24	24	• 00-00 = User Number • 01-24 = FCB
01	rr	ss	24	24	• 00-00 = Return Code • 01-24 = FCB 01 not copied to user.
37 RESET DRIVE: *EE*					
Message Sent to All Affected Servers					
00	ss	rr	25	01	• 00-01 = Drive Vector
01	rr	ss	25	00	• 00-00 = Return Code
38 ACCESS DRIVE: *EE*					
Message Sent to All Affected Servers					

FMT	DID	SID	FNC	SIZ	MSG
00	ss	rr	26	01	• 00-01 = Drive Vector
01	rr	ss	26	00	• 00-00 = Return Code
39 FREE DRIVE: *EE*					
Message Sent to All Affected Servers					
00	ss	rr	27	01	• 00-01 = Drive Vector
01	rr	ss	27	00	• 00-00 = Return Code
40 WRITE RANDOM WITH ZERO FILL: *EE*					
00	ss	rr	28	A4	<ul style="list-style-type: none"> • 00-00 = User Number • 01-24 = FCB • 25-A4 = Sector of Data to Write
01	rr	ss	28	24	<ul style="list-style-type: none"> • 00-00 = Return Code • 01-24 = FCB 01 not copied to user.
42 LOCK RECORD: *EE*					
00	ss	rr	2A	26	<ul style="list-style-type: none"> • 00-00 = User Number • 01-24 = FCB • 25-26 = File ID
01	rr	ss	2A	24	<ul style="list-style-type: none"> • 00-00 = Return Code • 01-24 = FCB 01 not copied to user.
43 UNLOCK RECORD: *EE*					
00	ss	rr	2B	26	<ul style="list-style-type: none"> • 00-00 = User Number • 01-24 = FCB • 25-26 = File ID
01	rr	ss	2B	24	<ul style="list-style-type: none"> • 00-00 = Return Code • 01-24 = FCB 01 not copied to user.
45 SET BDOS ERROR MODE: NOT IMPLEMENTED AT SERVER					
00	ss	rr	2D	00	• 00-00 = xx
01	rr	ss	2D	01	<ul style="list-style-type: none"> • 00-00 = 0FFh • 01-01 = 00Ch
46 GET DISK FREE SPACE: Extension for CP/M 3.					
00	ss	rr	2E	00	• 00-00 = Drive
01	rr	ss	2E	02	<ul style="list-style-type: none"> • 00-02 = Free Space, little-endian Copied to current DMA buffer.
48 FLUSH BUFFERS: Extension for CP/M 3.					
00	ss	rr	30	00	• 00-00 = Purge Code

FMT	DID	SID	FNC	SIZ	MSG
01	rr	ss	30	00	• 00-00 = Return Code
64 LOGIN: *EE*					
00	ss	rr	40	07	• 00-07 = Password, 8 ASCII Chars
01	rr	ss	40	00	• 00-00 = Return Code
65 LOGOFF: *EE*					
00	ss	rr	41	00	• 00-00 = xx
01	rr	ss	41	00	• 00-00 = Return Code
66 SEND MESSAGE ON NETWORK: NOT IMPLEMENTED AT SERVER					
00	ss	rr	42	xx	• 00-FF = xx
01	rr	ss	42	01	• 00-00 = 0FFh • 01-01 = 00Ch
67 RECEIVE MESSAGE ON NETWORK: NOT IMPLEMENTED AT SERVER					
00	ss	rr	43	00	• 00-00 = xx
01	rr	ss	43	01	• 00-00 = 0FFh • 01-01 = 00Ch
68 GET NETWORK STATUS: NOT IMPLEMENTED AT SERVER					
00	ss	rr	44	00	• 00-00 = xx
01	rr	ss	44	01	• 00-00 = 0FFh • 01-01 = 00Ch
69 GET CONFIGURATION TABLE ADDRESS: NOT IMPLEMENTED AT SERVER					
00	ss	rr	45	00	• 00-00 = xx
01	rr	ss	45	01	• 00-00 = 0FFh • 01-01 = 00Ch
70 SET COMPATIBILITY ATTRIBUTES: Message Sent to All Known Servers					
00	ss	rr	46	00	• 00-00 = Compatibility Attributes
01	rr	ss	46	00	• 00-00 = xx
71 RETURN SERVER CONFIGURATION: *EE*					
00	ss	rr	47	00	• 00-00 = xx
01	rr	ss	47	16	• 00-00 = Server Temporary File Drive • 01-01 = Server Status Byte • 02-02 = Server ID • 03-03 = Maximum Number of Requesters • 04-04 = Number Logged In • 05-06 = Login Vector

FMT	DID	SID	FNC	SIZ	MSG
					• 07-16 = Requester ID's
98 FREE BLOCKS: Extension for CP/M 3.					
00	ss	rr	62	00	• 00-00 = xx
01	rr	ss	62	00	• 00-00 = Return Code
99 TRUNCATE FILE: *EE* Extension for CP/M 3.					
00	ss	rr	63	24	• 00-00 = User Number • 01-24 = FCB
01	rr	ss	63	24	• 00-00 = Directory Code • 01-24 = FCB
101 GET DIRECTORY LABEL BYTE: Extension for CP/M 3.					
00	ss	rr	65	00	• 00-00 = Drive
01	rr	ss	65	00	• 00-00 = Dir Mode Byte
102 GET FILE TIME STAMPS: *EE* Extension for CP/M 3.					
00	ss	rr	66	24	• 00-00 = User Number • 01-24 = FCB
01	rr	ss	66	24	• 00-00 = Directory code • 01-24 = FCB with timestamps in d8-d15
105 GET DATE AND TIME: Extension for CP/M 3. Not issued by NDOS.					
00	ss	rr	69	00	• 00-00 = xx
01	rr	ss	69	04	• 00-01 = Date, days since 12/31/1977 • 02 = Hours (BCD) • 03 = Minutes (BCD) • 04 = Seconds (BCD)
106 SET DEFAULT PASSWORD: Message Sent to All Known Servers					
00	ss	rr	6A	07	• 00-07 = Default Password to be Set
01	rr	ss	6A	00	• 00-00 = Return Code

Table C-1. Conventional CP/NET Messages

Appendix D
NDOS Function Summary

Code	Function Name	Input Parameters	Output Results
38	Access Drive	DE = Drive Vector	none
39	Free Drive	DE = Drive Vector	none
42	Lock Record	DE = FCB Address	A = Err Code
43	Unlock Record	DE = FCB Address	A = Err Code
45	Set BDOS Error Mode	E = Error Mode	none
64	Login	see definition	A = Err Code
65	Logoff	E = Server ID	none
66	Send Message on Ntwrk	DE = Message Adr	A = Err Code
67	Receive Msg from Ntwk	DE = Message Adr	A = Err Code
68	Get Network Status	none	A = Status byte
69	Get Config Table Adr	none	HL = Table Adr
70	Set Compat. Attrs.	E = attributes	none
71	Get Server Config.	E = Server ID	HL= Table Adr
106	Set Default Password	see definition	none

Table D-1. NDOS Functions

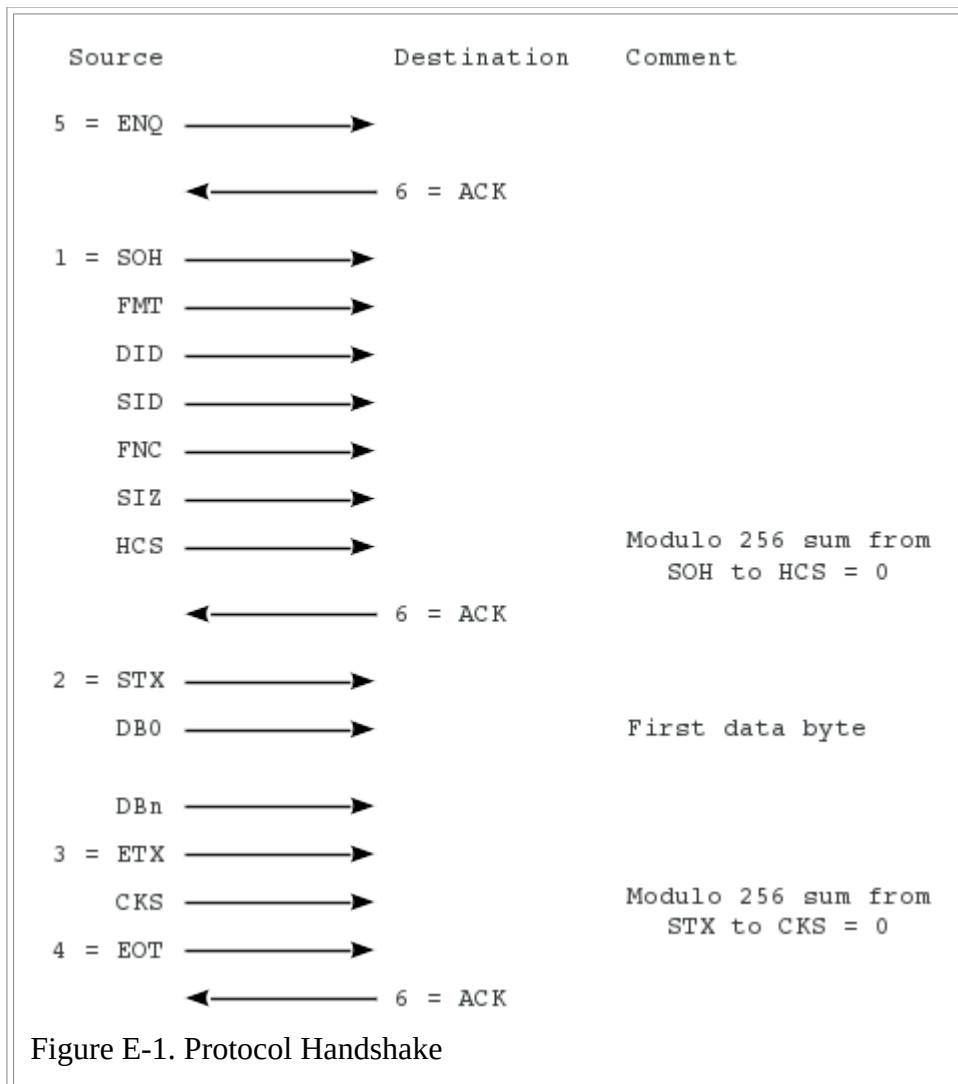
Appendix E A Simple RS-232C CP/NET System

Digital Research developed a relatively simple RS-232C point to-point protocol to provide a demonstration vehicle for CP/NET and to encourage compatibility among hardware vendors. The protocol, as implemented in the sample SNIOS and NETWRKIF in this appendix, breaks the logical message into a fixed header and a variable length data portion the size of which is obtained from the fixed header. This simplifies operation with DMA channels that need terminal counts and also provides a checksum for the header that contains the SIZ field.

This protocol can be implemented between any requester and server that support an extra RS-232 console port.

E.1 Protocol Handshake

The protocol handshake is detailed in [Figure E-1](#).



E.2 Binary Protocol Message Format

Data integrity for this protocol is maintained by a simple checksum, shown in [Figure E-2](#), on both the header and the actual message.

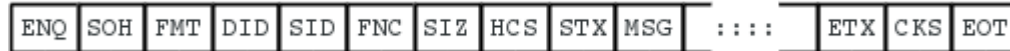


Figure E-2. Binary Protocol Message Format

Message format codes 00 & 01 are recommended.

Field Description:

- ENQ = Enquire, one byte, 05H.
- SOH = Start of Header, one byte, 01H.
- FMT,DID,SID,FNC,SIZ = as defined in [Appendix A](#), one byte per field.
- HCS = Header Checksum, one byte. This is a simple horizontal checksum, computed by adding together all the bytes of the message, starting with the SOH, to the SIZ byte of the header field modulo 256, complementing the result, and adding one. The entire message, from the SOH to and including the HCS, should add up to zero.
- STX = Start of Data, one byte, 02H.
- MSG = SIZ + 1 byte long.
- ETX = End of Data, one byte, 03H.
- CKS = Checksum, one byte. This is a simple horizontal checksum, computed by adding together all the bytes of the message, starting with the STX, to the last byte of the MSG field modulo 256, complementing the result, and adding one. The entire message, from the STX to and including the CKS, should add up to zero.
- EOT = End of Transmission, one byte, 04H.

E.3 ASCII Protocol Message Format

If the RS-232 link is not capable of transmitting 8-bit binary data, you might have to transmit each nibble of the message as a 7 bit ASCII character.

Note: the 7-bit ASCII network protocol is identical to the 8-bit protocol except that it requires twice as many bytes because each byte is transmitted in hexadecimal ASCII format.

The ASCII network protocol message format is detailed in [Figure E-3](#).

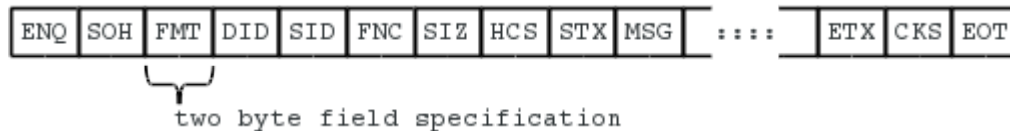


Figure E-3. ASCII Protocol Message Format

Message format codes 00 & 01 are recommended.

Field Description:

- ENQ = Enquire, one byte, 05H.
- SOH = Start of Header, one byte, 01H.
- FMT,DID,SID,FNC,SIZ = as defined in [Appendix A](#), two bytes per field.
- HCS = Header Checksum, 2 bytes (Hex-ASCII) . This is a simple horizontal checksum. It is computed by adding together all the bytes of the message, starting with the SOH, to the SIZ of the header field modulo 256, complementing the result, and adding one. The entire message, from the

SOH to the including the HCS, should add up to zero.

- STX = Start of Data, one byte, 02H.
- MSG = 2 * (SIZ + 1) bytes long.
- ETX = End of Data, one byte, 03H.
- CKS =Checksum, two bytes (Hex-ASCII) . This is a simple horizontal checksum. It is computed by adding together all the bytes of the message, starting with the STX, to the last byte of the MSG field modulo 256, complementing the result and adding one. The entire message, from the FMT to and including the CKS, should add up to zero.
- EOT = End of Transmission, one byte, 04H.

E.4 Modifying the SNIOS

The sample SNIOS can be modified for almost any requester that has a spare console port. To do so, follow these steps:

1. Obtain assembled listings of the SNIOS.ASM source file that require modification. You can use MAC, RMAC, or ASM. if you use ASM, the title, name, if, and else statements must be removed from the source files to assemble correctly. Using RMAC is highly recommended because it simplifies the task of generating the SPR files when used in conjunction with LINK. Otherwise, the SPR files must be generated in the same manner as for MP/M II XIOS.SPR generation.

```
A>RMAC SNIOS
```

2. Study the SNIOS.PRN listing. Notice the ASCII equate. If true, it specifies that the message format is 7-bit ASCII. If false, it specifies a binary 8-bit message format. The ASCII mode is sometimes useful in debugging, but in practice do not use it where it is possible to transmit 8 bit serial data.

The only code that requires modification in the SNIOS.ASM file is contained in the CHAROUT, CHARIN, and DELAY procedures. The CHAROUT and CHARIN procedures can be conditionally assembled for a Dynabyte DB8/2, now called DB8/5200, a Digital Microsystems DSC-2 or an ALTOS 8000-2. The NOPs in the CHAROUT procedure are simply padding, so the length of the DB8/2 SNIOS and DSC-2 SNIOS is the same, which helps in the debugging of these two versions.

Perhaps the most critical area in the SNIOS that requires adjustment for a specific network configuration is in the timeout code of the CHARIN procedure. If too little time is allowed, the server might not be able to complete the function because of a heavy request load from the requesters. If too much time is specified, communication breaks on the network can go undetected for a period of time, making both error recovery and precise detection difficult. Note that this is a logical timeout, not a data-link timeout. The logical timeout determines how long the requester expects the server to take between the time it receives the message and the time it returns a response message.

Another critical parameter that requires adjustment for different environments is ALWAYS\$RETRY. This equate, when true, controls conditional assembly that always produces retries on network failures. In this mode of operation, it is possible to recover from broken communication between the requester and a server. However, ALWAYS\$RETRY does hang the requester in a busy retry mode when failures occur.

```
1          title  'Requester Network I/O System for CP/NET 1.2'
2          page   54
3
4          ;*****
5          ;*****
6          ;**
7          ;** R e q u e s t e r   N e t w o r k   I / O   S y s t e m   **
8          ;**
9          ;*****
10         ;*****
11
12         ;/*
13         ; Copyright (C) 1980, 1981, 1982
14         ; Digital Research
15         ; P.O. Box 579
16         ; Pacific Grove, CA 93950
17         ;
18         ; Revised:  October 5, 1982
19         ;*/
20
21 0000 =      false  equ    0
22 FFFF =      true   equ    not false
23
24 0000 =      cpnos  equ    false          ; cp/net system
25
26 0000 =      DSC2   equ    false
27 0000 =      DB82   equ    false
28 FFFF =      Altos  equ    true
29
30 FFFF =      always$retry  equ    true    ; force continuous retries
31
32 0000 =      modem  equ    false
33
34 0000 =      ASCII  equ    false
35
36 0000 =      debug  equ    false
37
38           CSEG
39           if      cpnos
40           extrn   BDOS
41           else
42 0005 =      BDOS   equ    0005h
43           endif
44
45           NIOS:
46           public NIOS
47           ;      Jump vector for SNIOS entry points
48 0000 C3A900  jmp     ntwrkinit    ; network initialization
49 0003 C3B800  jmp     ntwrksts      ; network status
50 0006 C3C300  jmp     cnfgtbladr    ; return config table addr
51 0009 C3C700  jmp     sendmsg       ; send message on network
52 000C C33301  jmp     receivemsg    ; receive message from network
53 000F C3DD01  jmp     ntwrkerror    ; network error
54 0012 C3DE01  jmp     ntwrkwboot    ; network warm boot
```



```

109
110          DSEG
111
112
113          ;      Slave Configuration Table
114      configtbl:
115
116          Network$status:
117      0000          ds      1          ; network status byte
118      0001          ds      1          ; slave processor ID number
119      0002          ds      2          ; A: Disk device
120      0004          ds      2          ; B:  "
121      0006          ds      2          ; C:  "
122      0008          ds      2          ; D:  "
123      000A          ds      2          ; E:  "
124      000C          ds      2          ; F:  "
125      000E          ds      2          ; G:  "
126      0010          ds      2          ; H:  "
127      0012          ds      2          ; I:  "
128      0014          ds      2          ; J:  "
129      0016          ds      2          ; K:  "
130      0018          ds      2          ; L:  "
131      001A          ds      2          ; M:  "
132      001C          ds      2          ; N:  "
133      001E          ds      2          ; O:  "
134      0020          ds      2          ; P:  "
135
136      0022          ds      2          ; console device
137
138      0024          ds      2          ; list device:
139      0026          ds      1          ;      buffer index
140      0027 00          db      0          ;      FMT
141      0028 00          db      0          ;      DID
142      0029 56          db      Slave$ID ;      SID (CP/NOS must still initialize)
143      002A 05          db      5          ;      FNC
144      002B          ds      1          ;      SIZ
145      002C          ds      1          ;      MSG(0) List number
146      002D          ds      128         ;      MSG(1) ... MSG(128)
147
148          msg$adr:
149      00AD          ds      2          ; message address
150                  if      modem
151          timeout$retries equ 0          ; timeout a max of 256 times
152                  else
153      0064 =        timeout$retries equ 100 ; timeout a max of 100 times
154                  endif
155      000A =        max$retries equ 10      ; send message max of 10 times
156          retry$count:
157      00AF          ds      1
158
159          FirstPass:
160      00B0 FF          db      0ffh
161
162          ;      Network Status Byte Equates

```

```

163      ;
164 0010 = active      equ    0001$0000b    ; slave logged in on network
165 0002 = rcvrr      equ    0000$0010b    ; error in received message
166 0001 = senderr    equ    0000$0001b    ; unable to send message
167
168      ;      General Equates
169      ;
170 0001 = SOH      equ    01h            ; Start of Header
171 0002 = STX      equ    02h            ; Start of Data
172 0003 = ETX      equ    03h            ; End of Data
173 0004 = EOT      equ    04h            ; End of Transmission
174 0005 = ENQ      equ    05h            ; Enquire
175 0006 = ACK      equ    06h            ; Acknowledge
176 000A = LF       equ    0ah            ; Line Feed
177 000D = CR       equ    0dh            ; Carriage Return
178 0015 = NAK      equ    15h            ; Negative Acknowledge
179
180 0002 = conout    equ    2              ; console output function
181 0009 = print     equ    9              ; print string function
182 0043 = rcvmsg    equ    67             ; receive message NDOS function
183 0040 = login     equ    64             ; Login NDOS function
184
185      ;      I/O Equates
186      ;
187      if      DB82
188      stati    equ    83h
189      mski     equ    08h
190      dprti    equ    80h
191
192      stato    equ    83h
193      msko     equ    10h
194      statc    equ    81h
195      mskc     equ    20h
196      dprto    equ    86h
197      endif
198
199      if      DSC2
200      if      modem
201      stati    equ    59h
202      mski     equ    02h
203      dprti    equ    58h
204
205      stato    equ    59h
206      msko     equ    01h
207      dprto    equ    58h
208      else
209      stati    equ    51h
210      mski     equ    02h
211      dprti    equ    50h
212
213      stato    equ    51h
214      msko     equ    01h
215      dprto    equ    50h
216      endif

```

```
217                                     endif
218
219                                     if      Altos
220 001F =      stati equ    1fh
221 0001 =      mski equ    01h
222 001E =      dprti equ    1eh
223
224 001F =      stato equ    1fh
225 0004 =      msko equ    04h
226 001E =      dprto equ    1eh
227                                     endif
228
229
230
231                                     page
```

```

232
233             CSEG
234             ; Utility Procedures
235             ;
236             delay:                ; delay for c[a] * 0.5 milliseconds
237 002F 3E06      mvi    a,6
238             delay1:
239 0031 0E86      mvi    c,86h
240             delay2:
241 0033 0D        dcr    c
242 0034 C23300    jnz    delay2
243 0037 3D        dcr    a
244 0038 C23100    jnz    delay1
245 003B C9        ret
246
247             if    ASCII
248             Nib$out:                ; A = nibble to be transmitted in ASCII
249                 cpi    10
250                 jnc    nibAtoF      ; jump if A-F
251                 adi    '0'
252                 mov    c,a
253                 jmp    Char$out
254             nibAtoF:
255                 adi    'A'-10
256                 mov    c,a
257                 jmp    Char$out
258             endif
259
260             Pre$Char$out:
261 003C 7A        mov    a,d
262 003D 81        add    c
263 003E 57        mov    d,a                ; update the checksum in D
264
265             nChar$out:                ; C = byte to be transmitted
266             if    Altos
267 003F 3E10      mvi    a,10h
268 0041 D31F      out    stato
269             endif
270             in    stato
271 0045 E604      ani    msko
272 0047 CA3F00    jz     nChar$out
273
274             if    DB82
275             in    statc
276             ani    mskc
277             jz    nChar$out
278             endif
279
280             if    DSC2
281             nop                    ; these NOP's make DB8/2 & DSC2
282             nop                    ; versions the same length - saves
283             nop                    ; a second listing
284             nop
285             nop

```



```

286          nop
287          nop
288          endif
289
290 004A 79      mov     a,c
291 004B D31E    out     dprto
292 004D C9      ret
293
294          ; Char$out:
295 004E CD3F00  call    nChar$out
296          if     Altos
297 0051 E3E3E3E3 xthl! xthl! xthl! xthl
298 0055 E3E3E3E3 xthl! xthl! xthl! xthl
299 0059 E3E3E3E3 xthl! xthl! xthl! xthl ;delay 54 usec
300 005D C9      ret
301          else
302          jmp     delay          ; delay after each Char sent to Mstr
303          ret
304          endif
305
306          if     ASCII
307 Nib$in:      call    Char$in          ; return nibble in A register
308          call   Char$in
309          rc
310          ani    7fh
311          sui    '0'
312          cpi    10
313          jc     Nib$in$rtn      ; must be 0-9
314          adi    ('0'-'A'+10) and 0ffh
315          cpi    16
316          jc     Nib$in$rtn      ; must be 10-15
317          lda    network$status
318          ori    rcverr
319          sta    network$status
320          mvi    a,0
321          stc          ; carry set indicating err cond
322          ret
323
324 Nib$in$rtn:  ora     a          ; clear carry & return
325          ret
326          endif
327
328
329 xChar$in:
330 005E 0664    mvi     b,100        ; 100 ms corresponds to longest possible
331 0060 C36500  jmp     char$in0     ;wait between master operations
332
333 Char$in:
334          ; return byte in A register
335          ; carry set on rtn if timeout
336          if     modem
337          mvi    b,0          ; 256 ms = 7.76 chars @ 300 baud
338          else
339          if     Altos
340          mvi    b,3          ; 3 ms = 50 chars @ 125k baud

```

```
340             else
341             mvi    b,50             ; 50 ms = 50 chars @ 9600 baud
342             endif
343             endif
344             Char$in0:
345             0065 0E5A             mvi    c,5ah
346             Char$in1:
347             if    Altos
348             0067 3E00             mvi    a,0
349             0069 D31F             out    stati
350             endif
351             006B DB1F             in     stati
352             006D E601             ani    mski
353             006F C27C00           jnz   Char$in2
354             0072 0D              dcr   c
355             0073 C26700           jnz   Char$in1
356             0076 05              dcr   b
357             0077 C26500           jnz   Char$in0
358             007A 37              stc
359             007B C9              ret
360             Char$in2:
361             007C DB1E             in     dprti
362             007E C9              ret
363                                     ; rtn with raw char and carry cleared
364             Net$out:
365                                     ; C = byte to be transmitted
366                                     ; D = checksum
367             007F 7A              mov    a,d
368             0080 81              add    c
369             0081 57              mov    d,a
370
371             if    ASCII
372             mov    a,c
373             mov    b,a
374             rar
375             rar
376             rar
377             ani    0FH           ; mask HI-L0 nibble to L0 nibble
378             call  Nib$out
379             mov    a,b
380             ani    0FH
381             jmp   Nib$out
382
383             else
384             0082 C34E00           jmp   Char$out
385             endif
386
387             Msg$in:
388                                     ; HL = destination address
389                                     ; E = # bytes to input
390             0085 CD9000           call  Net$in
391             0088 D8              rc
392             0089 77              mov    m,a
393             008A 23              inx   h
394             008B 1D              dcr   e
```

```

394 008C C28500      jnz  Msg$in
395 008F C9          ret
396
397                Net$in:                ; byte returned in A register
398                                           ; D = checksum accumulator
399
400                if  ASCII
401                call Nib$in
402                rc
403                add  a
404                add  a
405                add  a
406                add  a
407                push psw
408                call Nib$in
409                pop  b
410                rc
411                ora  b
412
413                else
414 0090 CD6300      call Char$in      ;receive byte in Binary mode
415 0093 D8          rc
416                endif
417
418                chks$in:
419 0094 47          mov  b,a
420 0095 82          add  d                ; add & update checksum accum.
421 0096 57          mov  d,a
422 0097 B7          ora  a                ; set cond code from checksum
423 0098 78          mov  a,b
424 0099 C9          ret
425
426                Msg$out:                ; HL = source address
427                                           ; E = # bytes to output
428                                           ; D = checksum
429                                           ; C = preamble byte
430 009A 1600        mvi  d,0                ; initialize the checksum
431 009C CD3C00      call Pre$Char$out    ; send the preamble character
432                Msg$out$loop:
433 009F 4E          mov  c,m
434 00A0 23          inx  h
435 00A1 CD7F00      call Net$out
436 00A4 1D          dcr  e
437 00A5 C29F00      jnz  Msg$out$loop
438 00A8 C9          ret
439
440                page

```

```

441
442         ;      Network Initialization
443 ntwrkinit:
444
445         if      cpnos          ; copy down network assignments
446         lxi    h,Initconfigtbl
447         lxi    d,configtbl
448         mvi    c,initcflen
449 initloop:
450         mov    a,m
451         stax  d
452         inx   h
453         inx   d
454         dcr   c
455         jnz   initloop          ; initialize config tbl from ROM
456
457         else
458 00A9 3E56   mvi    a,slave$ID          ;initialize slave ID byte
459 00AB 320100 sta    configtbl+1        ; in the configuration tablee
460         endif
461
462         ;      device initialization, as required
463
464         if      Altos
465 00AE 3E47   mvi    a,047h
466 00B0 D30E   out    0eh
467 00B2 3E01   mvi    a,1
468 00B4 D30E   out    0eh
469         endif
470
471         if      DSC2 and modem
472         mvi    a,0ceh
473         out    stato
474         mvi    a,027h
475         out    stato
476         endif
477
478         if      cpnos
479         call   loginpr          ; login to a master
480         endif
481
482 initok:
483 00B6 AF     xra    a          ; return code is 0=success
484 00B7 C9     ret
485
486
487         page

```

```
488
489           ;      Network Status
490 ntwrksts:
491 00B8 3A0000      lda      network$status
492 00BB 47          mov      b,a
493 00BC E6FC       ani      not (rcvrr+senderr)
494 00BE 320000     sta      network$status
495 00C1 78         mov      a,b
496 00C2 C9        ret
497
498
499           ;      Return Configuration Table Address
500 cnfgtbladr:
501 00C3 210000     lxi      h,configtbl
502 00C6 C9        ret
503
504
505
506           page
```

```

507
508           ;      Send Message on Network
509 sendmsg:   ; BC = message addr
510           mov     h,b
511           mov     l,c           ; HL = message address
512           shld   msg$adr
513 re$sendmsg:
514           mvi    a,max$retries
515           sta    retry$count   ; initialize retry count
516 send:
517           lhld   msg$adr
518           mvi    c,ENQ
519           call   Char$out      ; send ENQ to master
520           mvi    d,timeout$retries
521 ENQ$response:
522           call   Char$in
523           jnc   got$ENQ$response
524           dcr   d
525           jnz   ENQ$response
526           jmp   Char$in$timeout
527 got$ENQ$response:
528           call   get$ACK0
529           mvi    c,SOH
530           mvi    e,5
531           call   Msg$out      ; send SOH FMT DID SID FNC SIZ
532           xra   a
533           sub   d
534           mov   c,a
535           call   net$out     ; send HCS (header checksum)
536           call   get$ACK
537           dcx   h
538           mov   e,m
539           inx   h
540           inr   e
541           mvi   c,STX
542           call   Msg$out     ; send STX DB0 DB1 ...
543           mvi   c,ETX
544           call   Pre$Char$out ; send ETX
545           xra   a
546           sub   d
547           mov   c,a
548           call   Net$out    ; send the checksum
549           mvi   c,EOT
550           call   nChar$out  ; send EOT
551           call   get$ACK    ; (leave these
552           ret              ; two instructions)
553
554 get$ACK:
555           call   Char$in
556           jc    send$retry  ; jump if timeout
557 get$ACK0:
558           ani   7fh
559           sui   ACK
560           rz

```

```
561          send$retry:
562 0123 E1          pop    h          ; discard return address
563 0124 21AF00      lxi    h, retry$count
564 0127 35          dcr    m
565 0128 C2D100      jnz    send          ; send again unles max retries
566          Char$in$timeout:
567 012B 3E01      mvi    a, senderr
568
569              if    always$retry
570 012D CDD201      call   error$return
571 0130 C3CC00      jmp    re$sendmsg
572              else
573              jmp    error$return
574              endif
575
576              page
```

```
577
578 ; Receive Message from Network
579 receivemsg: ; BC = message addr
580 0133 60 mov h,b
581 0134 69 mov l,c ; HL = message address
582 0135 22AD00 shld msg$adr
583 re$receivemsg:
584 0138 3E0A mvi a,max$retries
585 013A 32AF00 sta retry$count ; initialize retry count
586 re$call:
587 013D CD4F01 call receive ; rtn from receive is receive error
588
589 receive$retry:
590 0140 21AF00 lxi h,retry$count
591 0143 35 dcr m
592 0144 C23D01 jnz re$call
593 receive$timeout:
594 0147 3E02 mvi a,rcverr
595
596 if always$retry
597 0149 CDD201 call error$return
598 014C C33801 jmp re$receivemsg
599 else
600 jmp error$return
601 endif
602
603 receive:
604 014F 2AAD00 lhld msg$adr
605 0152 1664 mvi d,timeout$retries
606 receive$firstchar:
607 0154 CD5E00 call xcharin
608 0157 D26201 jnc got$firstchar
609 015A 15 dcr d
610 015B C25401 jnz receive$firstchar
611 015E E1 pop h ; discard receive$retry rtn adr
612 015F C34701 jmp receive$timeout
613 got$firstchar:
614 0162 E67F ani 7fh
615 0164 FE05 cpi ENQ ; Enquire?
616 0166 C24F01 jnz receive
617
618 0169 0E06 mvi c,ACK
619 016B CD3F00 call nChar$out ; acknowledge ENQ with an ACK
620
621 016E CD6300 call Char$in
622 0171 D8 rc ; return to receive$retry
623 0172 E67F ani 7fh
624 0174 FE01 cpi SOH ; Start of Header ?
625 0176 C0 rnz ; return to receive$retry
626 0177 57 mov d,a ; initialize the HCS
627 0178 1E05 mvi e,5
628 017A CD8500 call Msg$in
629 017D D8 rc ; return to receive$retry
630 017E CD9000 call Net$in
```



```

631 0181 D8          rc          ; return to receive$retry
632 0182 C2CD01     jnz      bad$checksum
633 0185 CDC501     call     send$ACK
634 0188 CD6300     call     Char$in
635 018B D8          rc          ; return to receive$retry
636 018C E67F      ani      7fh
637 018E FE02      cpi      STX          ; Start of Data ?
638 0190 C0         rnz          ; return to receive$retry
639 0191 57         mov      d,a          ; initialize the CKS
640 0192 2B         dcx      h
641 0193 5E         mov      e,m
642 0194 23         inx      h
643 0195 1C         inr      e
644 0196 CD8500     call     msg$in        ; get DB0 DB1 ...
645 0199 D8          rc          ; return to receive$retry
646 019A CD6300     call     Char$in        ; get the ETX
647 019D D8          rc          ; return to receive$retry
648 019E E67F      ani      7fh
649 01A0 FE03      cpi      ETX
650 01A2 C0         rnz          ; return to receive$retry
651 01A3 82         add      d
652 01A4 57         mov      d,a          ; update CKS with ETX
653 01A5 CD9000     call     Net$in        ; get CKS
654 01A8 D8          rc          ; return to receive$retry
655 01A9 CD6300     call     Char$in        ; get EOT
656 01AC D8          rc          ; return to receive$retry
657 01AD E67F      ani      7fh
658 01AF FE04      cpi      EOT
659 01B1 C0         rnz          ; return to receive$retry
660 01B2 7A         mov      a,d
661 01B3 B7         ora      a          ; test CKS
662 01B4 C2CD01     jnz      bad$checksum
663 01B7 E1         pop      h          ; discard receive$retry rtn adr
664 01B8 2AAD00     lhld    msg$adr
665 01BB 23         inx      h
666 01BC 3A0100     lda      configtbl+1
667 01BF 96         sub      m
668 01C0 CAC501     jz      send$ACK      ; jump with A=0 if DID ok
669 01C3 3EFF      mvi      a,0ffh      ; return code shows bad DID
670
671 01C5 F5         send$ACK: push    psw          ; save return code
672 01C6 0E06      mvi      c,ACK
673 01C8 CD3F00     call     nChar$out     ; send ACK if checksum ok
674 01CB F1         pop      psw          ; restore return code
675 01CC C9         ret
676
677
678
679 01CD 0E15      bad$DID: bad$checksum: mvi      c,NAK
680 01CF C34E00     jmp      Char$out      ; send NAK on bad chksm & not max retries
681 ;          ret
682
683
684 01D2 210000     error$return: lxi     h,network$status

```

```
685 01D5 B6          ora    m
686 01D6 77          mov    m,a
687 01D7 CDDD01      call   ntwrkerror    ; perform any required device re-init.
688 01DA 3EFF        mvi    a,0ffh
689 01DC C9          ret
690
691                ntwrkerror:
692                                ; perform any required device
693 01DD C9          ret                                ; re-initialization
694
695                page
```

```
696
697      ;
698      ntwrkwboot:
699
700      ;      This procedure is called each time the CCP is
701      ;      reloaded from disk. This version prints ""
702      ;      on the console and then returns, but anything necessary
703      ;      for restart can be put here.
704
705      01DE 0E09      mvi      c,9
706      01E0 111500   lxi      d,wboot$msg
707      01E3 C30500   jmp      BD0S
708
709      page
```

```

710
711         if      cpnos
712         ;
713         ;      LOGIN to a Master
714         ;
715         ; Equates
716         ;
717         buff    equ    0080h
718
719         readbf  equ    10
720
721         active  equ    0001$0000b
722
723         loginpr:
724             mvi    c,initpasswordmsglen
725             lxi    h,initpasswordmsg
726             lxi    d,passwordmsg
727         copypassword:
728             mov    a,m
729             stax   d
730             inx   h
731             inx   d
732             dcr   c
733             jnz   copypassword
734             mvi   c,print
735             lxi   d,loginmsg
736             call  BDOS
737             mvi   c,readbf
738             lxi   d,buff-1
739             mvi   a,50h
740             stax  d
741             call  BDOS
742             lxi   h,buff
743             mov   a,m      ; get # chars in the command tail
744             ora   a
745             jz    dologin ; default login if empty command tail
746             mov   c,a      ; A = # chars in command tail
747             xra   a
748             mov   b,a      ; B will accumulate master ID
749         scanblnks:
750             inx   h
751             mov   a,m
752             cpi   ' '
753             jnz   pastblnks ; skip past leading blanks
754             dcr   c
755             jnz   scanblnks
756             jmp   prelogin ; jump if command tail exhausted
757         pastblnks:
758             cpi   '['
759             jz    scanMstrID
760             mvi   a,8
761             lxi   d,passwordmsg+5+8-1
762             xchg
763         spacefill:

```

```

764             mvi    m, ' '
765             dcx    h
766             dcr    a
767             jnz    spacefill
768             xchg
769 scanLftBrkt:
770             mov    a,m
771             cpi    '['
772             jz     scanMstrID
773             inx    d
774             stax   d      ;update the password
775             inx    h
776             dcr    c
777             jnz    scanLftBrkt
778             jmp    prelogin
779 scanMstrID:
780             inx    h
781             dcr    c
782             jz     loginerr
783             mov    a,m
784             cpi    ']'
785             jz     prelogin
786             sui    '0'
787             cpi    10
788             jc     updateID
789             adi    ('0'-'A'+10) and 0ffh
790             cpi    16
791             jnc    loginerr
792 updateID:
793             push   psw
794             mov    a,b
795             add    a
796             add    a
797             add    a
798             add    a
799             mov    b,a      ; accum * 16
800             pop   psw
801             add    b
802             mov    b,a
803             jmp    scanMstrID
804
805 prelogin:
806             mov    a,b
807
808 dologin:
809             lxi    b,passwordmsg+1
810             stax   b
811             dcx    b
812             call   sendmsg
813             inr    a
814             lxi    d,loginfailedmsg
815             jz     printmsg
816             lxi    b,passwordmsg
817             call   receivemsg

```

```

818             inr    a
819             lxi    d,loginfailedmsg
820             jz     printmsg
821             lda    passwordmsg+5
822             inr    a
823             jnz    loginOK
824             jmp    printmsg
825
826 loginerr:
827             lxi    d,loginerrmsg
828 printmsg:
829             mvi    c,print
830             call   BDOS
831             jmp    loginpr      ; try login again
832
833 loginOK:
834             lxi    h,network$status ; HL = status byte addr
835             mov    a,m
836             ori    active ; set active bit true
837             mov    m,a
838             ret
839
840             ;
841             ; Local Data Segment
842             ;
843 loginmsg:
844             db     cr,lf
845             db     'LOGIN='
846             db     '$'
847
848 initpasswordmsg:
849             db     00h      ; FMT
850             db     00h      ; DID Master ID #
851             db     slave$ID ;SID
852             db     40h      ; FNC
853             db     7        ; SIZ
854             db     'PASSWORD' ; password
855 initpasswordmsglen equ $-initpasswordmsg
856
857
858 loginerrmsg:
859             db     lf
860             db     'Invalid LOGIN'
861             db     '$'
862
863 loginfailedmsg:
864             db     lf
865             db     'LOGIN Failed'
866             db     '$'
867
868             DSEG
869 passwordmsg:
870             ds     1        ; FMT
871             ds     1        ; DID

```

CP/M RMAC ASSEM 1.1 #021 REQUESTER NETWORK I/O SYSTEM FOR CP/NET 1.2

```
872          ds      1      ; SID
873          ds      1      ; FNC
874          ds      1      ; SIZ
875          ds      8      ; DAT = password
876          endif
877
878 01E6      end
```

CP/M RMAC ASSEM 1.1 #022 REQUESTER NETWORK I/O SYSTEM FOR CP/NET 1.2

ACK	0006	175#	559	618	672					
ACTIVE	0010	164#	721#	836						
ALTOS	FFFF	28#	62	219	266	296	338	347	464	
ALWAYSRETRY	FFFF	30#	569	596						
ASCII	0000	34#	247	306	370	400				
BADCHECKSUM	01CD	632	662	678#						
BADDID	01CD	677#								
BDOS	0005	40	42#	707	736	741	830			
CHARIN	0063	308	333#	414	522	555	621	634	646	655
CHARIN0	0065	331	344#	357						
CHARIN1	0067	346#	355							
CHARIN2	007C	353	360#							
CHARINTIMEOUT	012B	526	566#							
CHAROUT	004E	253	257	294#	384	519	680			
CHKSIN	0094	418#								
CNFGTBLADR	00C3	50	501#							
CONFIGTBL	0000	114#	447	459	502	666				
CONOUT	0002	180#								
CPNOS	0000	24#	39	66	445	478	711			
CR	000D	177#	844							
DB82	0000	27#	56	187	274					
DEBUG	0000	36#								
DEFAULTMASTER	0000	97#								
DELAY	002F	236#	302							
DELAY1	0031	238#	244							
DELAY2	0033	240#	242							
DPRTI	001E	190#	203#	211#	222#	361				
DPRT0	001E	196#	207#	215#	226#	291				
DSC2	0000	26#	59	199	280	471				
ENQ	0005	174#	518	615						
ENQRESPONSE	00DB	521#	525							
EOT	0004	173#	549	658						
ERRORRETURN	01D2	570	573	597	600	683#				
ETX	0003	172#	543	649						
FALSE	0000	21#	22	24	26	27	32	34	36	
FIRSTPASS	00B0	159#								
GETACK	0118	536	551	554#						
GETACK0	011E	528	557#							
GOTENQRESPONSE	00E8	523	527#							
GOTFIRSTCHAR	0162	608	613#							
INITOK	00B6	482#								
LF	000A	176#	844	859	864					
LOGIN	0040	183#								
MAXRETRIES	000A	155#	514	584						
MODEM	0000	32#	150	200	335	471				
MSGADR	00AD	148#	512	517	582	604	664			
MSGIN	0085	387#	394	628	644					
MSGOUT	009A	426#	531	542						
MSGOUTLOOP	009F	432#	437							
MSKI	0001	189#	202#	210#	221#	352				
MSKO	0004	193#	206#	214#	225#	271				
NAK	0015	178#	679							
NCHAROUT	003F	265#	272	277	295	550	619	673		
NETIN	0090	389	397#	630	653					

CP/M RMAC ASSEM	1.1	#023	REQUESTER	NETWORK	I/O	SYSTEM	FOR CP/NET	1.2
NETOUT	007F	364#	435	535	548			
NETWORKERRMSG	0021	103#						
NETWORKSTATUS	0000	116#	317	319	491	494	684	834
NIOS	0000	45#	46					
NTWRKERROR	01DD	53	687	691#				
NTWRKINIT	00A9	48	443#					
NTWRKSTS	00B8	49	490#					
NTWRKWBOOT	01DE	54	698#					
PRECHAROUT	003C	260#	431	544				
PRINT	0009	181#	734	829				
RCVERR	0002	165#	318	493	594			
RCVMSG	0043	182#						
RECALL	013D	586#	592					
RECEIVE	014F	587	603#	616				
RECEIVEFIRSTCHAR	0154	606#	610					
RECEIVMSG	0133	52	579#	817				
RECEIVERTRY	0140	589#						
RECEIVETIMEOUT	0147	593#	612					
RERECEIVMSG	0138	583#	598					
RESENDMSG	00CC	513#	571					
RETRYCOUNT	00AF	156#	515	563	585	590		
SEND	00D1	516#	565					
SENDACK	01C5	633	668	670#				
SENDERR	0001	166#	493	567				
SENDMSG	00C7	51	509#	812				
SENDRETRY	0123	556	561#					
SLAVEID	0056	57#	60#	63#	70	92	142	458 851
SOH	0001	170#	529	624				
STATI	001F	188#	201#	209#	220#	349	351	
STATO	001F	192#	205#	213#	224#	268	270	473 475
STX	0002	171#	541	637				
TIMEOUTRETRIES	0064	151#	153#	520	605			
TRUE	FFFF	22#	28	30				
WBOOTMSG	0015	99#	706					
XCHARIN	005E	329#	607					

Listing E-1: Request Network I/O System

E.5 Modifying the NETWRKIF

The NETWRKIF, designed for an Altos ACS 8000-10, is also easy to modify. The NETWRKIF implements the protocol by checking for the first character of an incoming message through one of the XIOS CONIN routines. After receiving the first character and validating it, the NETWRKIF disables interrupts and reads the rest of the message in under direct process control. If an XIOS CONIN routine does not exist for the port to be used for the network, you must write one.

To modify this NETWRKIF, follow these steps:

1. Set the NMB\$SLVS equate to the number of requesters to be supported. If more than four must be supported, you must add more Process Descriptors and queues.
2. If the server can only transmit or receive one message at a time, then the NETWRKIF supports a mutual exclusion queue to prevent collisions. To use this queue, set MUTEXIN or MUTEXOUT to true.
3. If the server is running on a Z80 processor, set Z80 to true for more efficient implementation of character I/O.
4. If all or some of the network RS-232 ports support only 7 bit ASCII, modify the BINARYASCII table by setting the appropriate entries to 0.
5. Modify the network port definitions. CONSOLE4\$STATUS through PRINTER2\$STATUS must be modified. Also, CHARIOTBL must be modified, so that the console numbers associated with the ports listed in STATUS\$PORTS match.
6. I/O port numbers in the routines CHAR\$OUT and CHAR\$IN might have to be modified. You might

have to implement a I/O port table similar to STATUS\$PORTS. This implementation relies on the fact that the Altos ACS 8000-10 always positions its I/O ports at a fixed offset from its status ports.

The sample NETWRKIF contains a debug conditional assembly flag that permits generation of a NETWRKIF.COM file. The NETWRKIF.COM version can debug a single requester, as follows:

1. Perform a GENSYMS in which the SERVER.RSP is included; do not include a NETWRKIF.RSP. During the GENSYMS, do not specify bank-switched memory.
2. Execute the MPM.SYS produced from GENSYMS, and load the NETWRKIF.COM file with DDT, SID, or ZSID.
3. Use DDT, SID, or ZSID to debug the NETWRKIF process. This works only for a single requester.

```

CP/M RMAC ASSEM 1.1    #001    MASTER NETWORK I/F MODULE

1          title    'Master Network I/F Module'
2          page     54
3
4          ;*****
5          ;*****
6          ;**
7          ;**      S e r v e r   N e t w o r k   I / F   M o d u l e      **
8          ;**
9          ;*****
10         ;*****
11
12         ;/*
13         ; Copyright (C) 1980
14         ; Digital Research
15         ; P.O. Box 579
16         ; Pacific Grove, CA 93950
17         ;
18         ; Modified October 5, 1982
19         ;
20         ;*/
21
22
23 0000 =    false   equ    0
24 FFFF =    true    equ    not false
25
26 FFFF =    z80     equ    true
27
28 0000 =    debug   equ    false
29 0000 =    modem   equ    false
30
31 0000 =    WtchDg  equ    false          ; include watch dog timer
32
33 0000 =    mutexin equ    false          ; provide mutual exclusion on input
34 0000 =    mutexout equ    false        ; provide mutual exclusion on output
35
36
37         if      debug
38
39         NmbSlvs equ    1                ;debug only one requester
40
41         lxi     sp,NtwrkIS0+2eh
42         mvi     c,145
43         mvi     e,64
44         call    bdos                    ; set priority to 64
45         lxi     h,UQCBNtwrkQI0        ; initialize reentrant variables
46         lxi     d,UQCBNtwrkQ00
47         lxi     b,BufferQ0
48         mvi     a,00h
49         ret
50
51         bdosadr:
52         dw      0005h
53
54         else

```

```

CP/M RMAC ASSEM 1.1    #002    MASTER NETWORK I/F MODULE

```

```

55
56 0002 =      NmbSlvs equ      2      ;RSP is configured for two requesters
57
58      bdosadr:
59 0000 0000      dw      $-$      ;XDOS entry point for RSP version
60
61      endif
62
63      ; Network Interface Process #0
64
65      NtwrkIP0:
66 0002 0000      dw      0      ; link
67 0004 00      db      0      ; status
68 0005 40      db      64      ; priority
69 0006 6400      dw      NtwrkIS0+46      ; stack pointer
70 0008 4E7477726B      db      'NtwrkIP0'      ; name
71 0010 00      db      0      ; console
72 0011 FF      db      0ffh      ; memseg
73 0012      ds      2      ; b
74 0014      ds      2      ; thread
75 0016      ds      2      ; buff
76 0018      ds      1      ; user code & disk slct
77 0019      ds      2      ; dcnt
78 001B      ds      1      ; searchl
79 001C      ds      2      ; searcha
80 001E      ds      2      ; active drives
81 0020 0000      dw      0      ; HL'
82 0022 0000      dw      0      ; DE'
83 0024 0000      dw      0      ; BC'
84 0026 0000      dw      0      ; AF'
85 0028 0000      dw      0      ; IY
86 002A 0000      dw      0      ; IX
87 002C 8000      dw      UQCBNtwrkQI0      ; HL
88 002E A000      dw      UQCBNtwrkQ00      ; DE
89 0030 A600      dw      BufferQ0      ; BC
90 0032 0000      dw      0      ; AF, A = ntwkif console dev #
91 0034      ds      2      ; scratch
92
93      NtwrkIS0:
94 0036 C7C7C7C7      dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
95 003E C7C7C7C7      dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
96 0046 C7C7C7C7      dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
97 004E C7C7C7C7      dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
98 0056 C7C7C7C7      dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
99 005E C7C7C7C7      dw      0c7c7h,0c7c7h,0c7c7h
100 0064 4206      dw      setup
101
102      QCBNtwrkQI0:
103 0066      ds      2      ; link
104 0068 4E7477726B      db      'NtwrkQI0'      ; name
105 0070 0200      dw      2      ; msglen
106 0072 0100      dw      1      ; nmbmsgs
107 0074      ds      2      ; dqph
108 0076      ds      2      ; nqph

```

CP/M RMAC ASSEM 1.1 #003 MASTER NETWORK I/F MODULE

```

109 0078      ds      2      ; msgin
110 007A      ds      2      ; msgout
111 007C      ds      2      ; msgcnt
112 007E      ds      2      ; buffer
113
114      UQCBNtwrkQI0:
115 0080 6600      dw      QCBNtwrkQI0      ; pointer
116 0082 8400      dw      BufferQI0Addr      ; msgadr
117      BufferQI0Addr:
118 0084 A600      dw      BufferQ0
119
120      QCBNtwrkQ00:
121 0086      ds      2      ; link
122 0088 4E7477726B      db      'NtwrkQ00'      ; name
123 0090 0200      dw      2      ; msglen
124 0092 0100      dw      1      ; nmbmsgs
125 0094      ds      2      ; dqph
126 0096      ds      2      ; nqph
127 0098      ds      2      ; msgin

```

```

128 009A          ds      2          ; msgout
129 009C          ds      2          ; msgcnt
130 009E          ds      2          ; buffer
131
132              UQCBNtwrkQ00:
133 00A0 8600      dw      QCBNtwrkQ00    ; pointer
134 00A2 A400      dw      BufferQ00Addr    ; msgadr
135              BufferQ00Addr:
136 00A4          ds      2
137
138              BufferQ0:
139 00A6          ds      1          ; FMT
140 00A7          ds      1          ; DID
141 00A8          ds      1          ; SID
142 00A9          ds      1          ; FNC
143 00AA          ds      1          ; SIZ
144 00AB          ds      257       ; MSG
145
146              ;      Network Interface Process #1
147
148              if      NmbSlvs GE 2
149              NtwrkIPI1:
150
151              if      NmbSlvs GE 3
152              dw      NtwrkIP2    ; link
153              else
154 01AC 0000      dw      0          ; link
155              endif
156
157 01AE 00        db      0          ; status
158 01AF 40        db      64         ; priority
159 01B0 0E02      dw      NtwrkIS1+46    ; stack pointer
160 01B2 4E7477726B db      'NtwrkIP1'    ; name
161 01BA 00        db      0          ; console
162 01BB FF        db      0ffh     ; memseg

```

CP/M RMAC ASSEM 1.1 #004 MASTER NETWORK I/F MODULE

```

163 01BC          ds      2          ; b
164 01BE          ds      2          ; thread
165 01C0          ds      2          ; buff
166 01C2          ds      1          ; user code & disk slct
167 01C3          ds      2          ; dcnt
168 01C5          ds      1          ; searchl
169 01C6          ds      2          ; searcha
170 01C8          ds      2          ; active drives
171 01CA 0000      dw      0          ; HL'
172 01CC 0000      dw      0          ; DE'
173 01CE 0000      dw      0          ; BC'
174 01D0 0000      dw      0          ; AF'
175 01D2 0000      dw      0          ; IY
176 01D4 0000      dw      0          ; IX
177 01D6 2A02      dw      UQCBNtwrkQI1 ; HL
178 01D8 4A02      dw      UQCBNtwrkQ01 ; DE
179 01DA 5002      dw      BufferQ1    ; BC
180 01DC 0001      dw      0100h     ; AF, A = ntwkif console dev #
181 01DE          ds      2          ; scratch
182
183              NtwrkIS1:
184 01E0 C7C7C7C7  dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
185 01E8 C7C7C7C7  dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
186 01F0 C7C7C7C7  dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
187 01F8 C7C7C7C7  dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
188 0200 C7C7C7C7  dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
189 0208 C7C7C7C7  dw      0c7c7h,0c7c7h,0c7c7h
190 020E 6906      dw      init
191
192              QCBNtwrkQI1:
193 0210          ds      2          ; link
194 0212 4E7477726B db      'NtwrkQI1'    ; name
195 021A 0200      dw      2          ; msglen
196 021C 0100      dw      1          ; nmbmsgs
197 021E          ds      2          ; dqph
198 0220          ds      2          ; nqph
199 0222          ds      2          ; msgin
200 0224          ds      2          ; msgout
201 0226          ds      2          ; msgcnt

```

```

202 0228          ds      2          ; buffer
203
204          UQCBNtwrkQI1:
205 022A 1002      dw      QCBNtwrkQI1  ; pointer
206 022C 2E02      dw      BufferQI1Addr ; msgadr
207          BufferQI1Addr:
208 022E 5002      dw      BufferQ1
209
210          QCBNtwrkQ01:
211 0230          ds      2          ; link
212 0232 4E7477726B db      'NtwrkQ01'  ; name
213 023A 0200      dw      2          ; msglen
214 023C 0100      dw      1          ; nmbmsg
215 023E          ds      2          ; dqph
216 0240          ds      2          ; nqph

```

CP/M RMAC ASSEM 1.1 #005 MASTER NETWORK I/F MODULE

```

217 0242          ds      2          ; msgin
218 0244          ds      2          ; msgout
219 0246          ds      2          ; msgcnt
220 0248          ds      2          ; buffer
221
222          UQCBNtwrkQ01:
223 024A 3002      dw      QCBNtwrkQ01  ; pointer
224 024C 4E02      dw      BufferQ01Addr ; msgadr
225          BufferQ01Addr:
226 024E          ds      2
227
228          BufferQ1:
229 0250          ds      1          ; FMT
230 0251          ds      1          ; DID
231 0252          ds      1          ; SID
232 0253          ds      1          ; FNC
233 0254          ds      1          ; SIZ
234 0255          ds      257        ; MSG
235          endif
236
237          ; Network Interface Process #2
238
239          if NmbSlvs GE 3
240          NtwrkIP2:
241
242          if NmbSlvs GE 4
243          dw      NtwrkIP3          ; link
244          else
245          dw      0                  ; link
246          endif
247
248          db      0                  ; status
249          db      64                 ; priority
250          dw      NtwrkIS2+46        ; stack pointer
251          db      'NtwrkIP2'        ; name
252          db      0                  ; console
253          db      0ffh              ; memseg
254          ds      2                  ; b
255          ds      2                  ; thread
256          ds      2                  ; buff
257          ds      1                  ; user code & disk slct
258          ds      2                  ; dcnt
259          ds      1                  ; searchl
260          ds      2                  ; searcha
261          ds      2                  ; active drives
262          dw      0                  ; HL'
263          dw      0                  ; DE'
264          dw      0                  ; BC'
265          dw      0                  ; AF'
266          dw      0                  ; IY
267          dw      0                  ; IX
268          dw      UQCBNtwrkQI2      ; HL
269          dw      UQCBNtwrkQ02      ; DE
270          dw      BufferQ2           ; BC

```

CP/M RMAC ASSEM 1.1 #006 MASTER NETWORK I/F MODULE

```

271          dw      0200h            ; AF, A = ntwkif console dev #

```

```

272             ds      2             ; scratch
273
274     NtwrkIS2:
275             dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
276             dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
277             dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
278             dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
279             dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
280             dw      0c7c7h,0c7c7h,0c7c7h
281             dw      init
282
283     QCBNtwrkQI2:
284             ds      2             ; link
285             db      'NtwrkQI2'   ; name
286             dw      2             ; msglen
287             dw      1             ; nmbmsgs
288             ds      2             ; dqph
289             ds      2             ; nqph
290             ds      2             ; msgin
291             ds      2             ; msgout
292             ds      2             ; msgcnt
293             ds      2             ; buffer
294
295     UQCBNtwrkQI2:
296             dw      QCBNtwrkQI2   ; pointer
297             dw      BufferQI2Addr  ; msgadr
298     BufferQI2Addr:
299             dw      BufferQ2
300
301     QCBNtwrkQ02:
302             ds      2             ; link
303             db      'NtwrkQ02'   ; name
304             dw      2             ; msglen
305             dw      1             ; nmbmsgs
306             ds      2             ; dqph
307             ds      2             ; nqph
308             ds      2             ; msgin
309             ds      2             ; msgout
310             ds      2             ; msgcnt
311             ds      2             ; buffer
312
313     UQCBNtwrkQ02:
314             dw      QCBNtwrkQ02   ; pointer
315             dw      BufferQ02Addr  ; msgadr
316     BufferQ02Addr:
317             ds      2
318
319     BufferQ2:
320             ds      1             ; FMT
321             ds      1             ; DID
322             ds      1             ; SID
323             ds      1             ; FNC
324             ds      1             ; SIZ

```

CP/M RMAC ASSEM 1.1 #007 MASTER NETWORK I/F MODULE

```

325             ds      257          ; MSG
326     endif
327
328     ;      Network Interface Process #3
329
330     if      NmbSlvs GE 4
331     NtwrkIP3:
332             dw      0             ; link
333             db      0             ; status
334             db      64            ; priority
335             dw      NtwrkIS3+46   ; stack pointer
336             db      'NtwrkIP3'   ; name
337             db      0             ; console
338             db      0ffh         ; memseg
339             ds      2             ; b
340             ds      2             ; thread
341             ds      2             ; buff
342             ds      1             ; user code & disk slct
343             ds      2             ; dcnt
344             ds      1             ; searchl
345             ds      2             ; searcha

```

```

346          ds      2          ; active drives
347          dw      0          ; HL'
348          dw      0          ; DE'
349          dw      0          ; BC'
350          dw      0          ; AF'
351          dw      0          ; IY
352          dw      0          ; IX
353          dw      UQCBNtwrkQI3 ; HL
354          dw      UQCBNtwrkQ03 ; DE
355          dw      BufferQ3      ; BC
356          dw      0300h        ; AF, A = ntwkif console dev #
357          ds      2          ; scratch
358
359          NtwrkIS3:
360          dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
361          dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
362          dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
363          dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
364          dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
365          dw      0c7c7h,0c7c7h,0c7c7h
366          dw      init
367
368          QCBNtwrkQI3:
369          ds      2          ; link
370          db      'NtwrkQI3'  ; name
371          dw      2          ; msglen
372          dw      1          ; nmbmsgs
373          ds      2          ; dqph
374          ds      2          ; nqph
375          ds      2          ; msgin
376          ds      2          ; msgout
377          ds      2          ; msgcnt
378          ds      2          ; buffer

```

CP/M RMAC ASSEM 1.1 #008 MASTER NETWORK I/F MODULE

```

379
380          UQCBNtwrkQI3:
381          dw      QCBNtwrkQI3  ; pointer
382          dw      BufferQI3Addr ; msgadr
383          BufferQI3Addr:
384          dw      BufferQ3
385
386          QCBNtwrkQ03:
387          ds      2          ; link
388          db      'NtwrkQ03'  ; name
389          dw      2          ; msglen
390          dw      1          ; nmbmsgs
391          ds      2          ; dqph
392          ds      2          ; nqph
393          ds      2          ; msgin
394          ds      2          ; msgout
395          ds      2          ; msgcnt
396          ds      2          ; buffer
397
398          UQCBNtwrkQ03:
399          dw      QCBNtwrkQ03  ; pointer
400          dw      BufferQ03Addr ; msgadr
401          BufferQ03Addr:
402          ds      2
403
404          BufferQ3:
405          ds      1          ; FMT
406          ds      1          ; DID
407          ds      1          ; SID
408          ds      1          ; FNC
409          ds      1          ; SIZ
410          ds      257        ; MSG
411          endif
412
413
414          if      WtchDg
415          ; Watchdog Timer Process
416          ;
417          WatchDogPD:
418
419          if      NmbSlvs GT 1

```

```

420          dw      NtwrkIP1      ; link to the remaining NETWRKIF PD's
421          else
422          dw      0              ; link
423          endif
424
425          db      0              ; status
426          db      64             ; priority
427          dw      WatchDogSTK+46 ; stack pointer
428          db      'WatchDog'     ; name
429          db      0              ; console
430          db      0ffh           ; memseg
431          ds      2              ; b
432          ds      2              ; thread

```

CP/M RMAC ASSEM 1.1 #009 MASTER NETWORK I/F MODULE

```

433          ds      2              ; buff
434          ds      1              ; user code & disk slct
435          ds      2              ; dcnt
436          ds      1              ; searchl
437          ds      2              ; searcha
438          ds      2              ; active drives
439          dw      0              ; HL'
440          dw      0              ; DE'
441          dw      0              ; BC'
442          dw      0              ; AF'
443          dw      0              ; IY
444          dw      0              ; IX
445          dw      0              ; HL
446          dw      0              ; DE
447          dw      0              ; BC
448          dw      0              ; AF
449          ds      2              ; scratch
450
451          WatchDogSTK:
452          dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
453          dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
454          dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
455          dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
456          dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
457          dw      0c7c7h,0c7c7h,0c7c7h
458          dw      WatchDog
459
460          WatchDogTime:
461          dw      $-$            ; one-second counter
462
463          WatchDogTable:
464          ;      Waiting Timeout   Start   Flag   Requester
465          db      0,    0,    0,0,    0ah    ; #0
466          db      0,    0,    0,0,    0bh    ; #1
467          db      0,    0,    0,0,    0fh    ; #2
468          db      0,    0,    0,0,    0dh    ; #3
469          endif
470
471          if      mutexin or mutexout
472          QCBMXSXmitq:          ; MX queue for requester transmitting
473
474          ds      2              ; link
475          db      'MXSXmitq'    ; name
476          dw      0              ; msglen
477          dw      1              ; nmbmsgs
478          ds      2              ; dqph
479          ds      2              ; nqph
480          ds      2              ; msgin
481          ds      2              ; msgout
482          ds      2              ; msgcnt
483          ds      2              ; buffer (owner PD)
484
485          UQCBMXSXmitq:
486          dw      QCBMXSXmitq

```

CP/M RMAC ASSEM 1.1 #010 MASTER NETWORK I/F MODULE

```

487          ;      dw      0              ; no message, since it's an MX queue
488          ;      db      'MXSXmitq'    ; no name, since the QCB pointer is resolved
489          ;      endif

```



```

490
491           ;      Server Configuration Table
492
493      configtbl:
494  0356 00      db      0           ; Server status byte
495  0357 00      db      0           ; Server ID
496  0358 02      db      NmbSlvs    ; Maximum number of requesters supported
497  0359 00      db      0           ; Number of requesters currently logged-in
498  035A 0000    dw      0000h      ; 16 bit vector of logged in requesters
499  035C         ds      16          ; Requester ID's currently logged-in
500  036C 5041535357 db      'PASSWORD' ; login password
501
502  0001 =      nmsg      equ      1     ; number of messages buffered
503  0096 =      slave$stk$len equ      96h ; server process stack size
504
505           if      NmbSlvs GE 2
506      slave1$stk:
507  0374         ds      slave$stk$len-2
508  0408 0A04    dw      Slave1
509
510           endif
511
512           if      NmbSlvs GE 3
513      slave2$stk:
514         ds      slave$stk$len-2
515         dw      Slave2
516           endif
517
518           if      NmbSlvs GE 4
519      slave3$stk:
520         ds      slave$stk$len-2
521         dw      Slave3
522           endif
523
524           if      NmbSlvs GE 2
525      Slave1:
526  040A         ds      52           ; SERVR1PR processor descriptor
527           endif
528
529           if      NmbSlvs GE 3
530      Slave2:
531         ds      52           ; SERVR2PR processor descriptor
532           endif
533
534           if      NmbSlvs GE 4
535      Slave3:
536         ds      52           ; SERVR3PR processor descriptor
537           endif
538
539           ;      Local Data Segment
540

```

CP/M RMAC ASSEM 1.1 #011 MASTER NETWORK I/F MODULE

```

541      BinaryASCII:
542  043E FF      db      0ffh        ; Requester #0: 0=7 bit ASCII, FF=8 bit binary
543  043F FF      db      0ffh        ;      #1
544  0440 FF      db      0ffh        ;      #2
545  0441 FF      db      0ffh        ;      #3
546
547      Networkstatus:
548  0442 00      db      0           ; Slave #0 network status byte
549  0443 00      db      0           ;      #1
550  0444 00      db      0           ;      #2
551  0445 00      db      0           ;      #3
552
553  0446 0000    conin: dw      $-$    ; save area for XIOS routine address
554
555  000A =      max$retries equ      10 ; maximum send message retries
556
557           ;
558           ;      The following tables are for use in the ALTOS i/o routines.
559           ;      Note that this program MUST be used with an XIOS which allows
560           ;      using the second printer port as a console port - Accessed as console
561           ;      #4
562  002B =      Console4$status equ      02bh
563  002F =      Console3$status equ      02fh

```

```

564 002D = Console2$status equ 02dh
565 0029 = Printer2$status equ 029h ; ALSO CONSOLE #4
566
567 if z80
568 ;
569 ; ENTRIES IN THE FOLLOWING TWO TABLES MUST MATCH !!!!
570
571 status$ports:
572 0448 2B db Console4$status ; Console 4 (Requester 0) status port
573 0449 2F db Console3$status ; Console 3 (Requester 1) status port
574 044A 2D db Console2$status ; Console 2 (Requester 2) status port
575 044B 29 db Printer2$status ; Printer 2 (Requester 3) status port
576 endif
577
578 chariotbl: ; Relationship between requesters and consoles
579 044C 03 db 3
580 044D 02 db 2
581 044E 01 db 1
582 044F 04 db 4
583
584 ; Network Status Byte Equates
585
586 0080 = ntwrktxrdy equ 10000000b ; NETWRKIF ready to send msg
587 0010 = active equ 00010000b ; requester logged into network
588 0008 = msgerr equ 00001000b ; error in received message
589 0004 = ntwrk equ 00000100b ; network alive
590 0002 = msgovr equ 00000010b ; message overrun
591 0001 = ntwrkrxrdy equ 00000001b ; NETWRKIF has rcvd msg
592
593 ; BDOS and XDOS Equates
594

```

CP/M RMAC ASSEM 1.1 #012 MASTER NETWORK I/F MODULE

```

595 0085 = flagset equ 133 ; flag set
596 0086 = makeq equ 134 ; make queue
597 0089 = readq equ 137 ; read queue
598 008B = writeq equ 139 ; write queue
599 008D = delay equ 141 ; delay
600 008E = dsptch equ 142 ; dispatch
601 0090 = createp equ 144 ; create process
602 009A = sydatad equ 154 ; system data page address
603 0083 = poll equ 083h ; Poll device
604
605 ; General Equates
606
607 0001 = SOH equ 01h ; Start of Header
608 0002 = STX equ 02h ; Start of Data
609 0003 = ETX equ 03h ; End of Data
610 0004 = EOT equ 04h ; End of Transmission
611 0005 = ENQ equ 05h ; Enquire
612 0006 = ACK equ 06h ; Acknowledge
613 000A = LF equ 0ah ; Line Feed
614 000D = CR equ 0dh ; Carriage Return
615 0015 = NAK equ 15h ; Negative Acknowledge
616
617 0010 = printer2 equ 10h ; special poll device number for second
618 ; printer port
619
620 ; Utility Procedures
621
622 bdos:
623 0450 2A0000 lhld bdosadr ; get XDOS entry point from RSP start
624 0453 E9 pchl
625
626 Nibout: ; A = nibble to be transmitted in ASCII
627 0454 FE0A cpi 10
628 0456 D25F04 jnc nibatof ; jump if A-F
629 0459 C630 adi '0'
630 045B 4F mov c,a
631 045C C36804 jmp Charout
632
633 045F C637 nibatof: adi 'A' -10
634 0461 4F mov c,a
635 0462 C36804 jmp Charout
636
637 PreCharout:

```

```

638 0465 7A      mov    a,d
639 0466 81      add    c
640 0467 57      mov    d,a          ; update the checksum
641
642          if    z80          ; Z80 version, using OUT A,(C) instruction
643 char$out:
644
645          ;      Character output routine for network i/o
646          ;      using the ALTOS SIO ports
647          ;
648          ;      Z80 version: this can use indirect port numbers in a clean,

```

CP/M RMAC ASSEM 1.1 #013 MASTER NETWORK I/F MODULE

```

649          ;      reentrant fashion
650          ;
651          ;      Entry: C register contains 8 bit value to transmit
652          ;      Entry : Slave number in register b
653
654 0468 E5      push   h
655 0469 D5      push   d
656 046A C5      push   b
657 046B 51      mov    d, c          ; save the character
658 046C 214804  lxi   h, status$ports
659 046F 48      mov    c, b
660 0470 0600    mvi   b, 0          ; set (BC) = (b)
661 0472 09      dad    b
662 0473 4E      mov    c,m
663
664          ;      Now C contains the address of the correct status port
665
666 outputloop:
667 0474 3E10    mvi   a,10h
668
669          ;      out    (c),a
670 0476 ED79    db    0edh,79h
671
672          ;      in     a,(c)
673 0478 ED78    db    0edh,78h
674
675 047A E604    ani   04h          ; wait for TXready
676 047C CA7404  jz    outputloop
677
678          ;      In the Altos system, data registers are one below status registers...
679
680 047F 0D      dcr    c
681
682          ;      out    (c),d
683 0480 ED51    db    0edh,51h
684
685 0482 C1      pop    b
686 0483 D1      pop    d
687 0484 E1      pop    h
688 0485 C9      ret
689
690          else
691
692 char$out:
693
694          ;      Character output routine for network I/O
695          ;      using ALTOS SIO ports
696          ;
697          ;      8080 version: This has to dispatch and then use direct port I/O
698          ;      --extremely messy to do reentrantly
699          ;
700          ;      Entry: C = character to transmit
701          ;      B = slave id byte
702

```

CP/M RMAC ASSEM 1.1 #014 MASTER NETWORK I/F MODULE

```

703          push   h
704          push   d
705          push   b
706
707          lxi   d,out0          ; dispatch address =

```

```

708      mov     l,b           ; out0 + slaveid*16
709      mvi     h,0
710      dad     h
711      dad     h
712      dad     h
713      dad     h
714      dad     d
715      mvi     a,10h        ;load "get transmit status" value
716      pchl                    ;dispatch
717
718      out0:
719      out     Console4$status ;wait for TXready status
720      in      Console4$status
721      ani     4
722      jz     out0
723
724      mov     a,c
725      out     Console4$status-1 ;write the character
726      pop     b
727      pop     d
728      pop     h
729      ret
730
731      out1:  out     Console3$status
732      in      Console3$status
733      ani     4
734      jz     out1
735
736      mov     a,c
737      out     Console3$status-1
738      pop     b
739      pop     d
740      pop     h
741      ret
742
743      out2:  out     Console2$status
744      in      Console2$status
745      ani     4
746      jz     out2
747
748      mov     a,c
749      out     Console2$status-1
750      pop     b
751      pop     d
752      pop     h
753      ret
754
755      out3:  out     Printer2$status
756      in      Printer2$status

```

CP/M RMAC ASSEM 1.1 #015 MASTER NETWORK I/F MODULE

```

757      ani     4
758      jz     out3
759
760      mov     a,c
761      out     Printer2$status-1
762      pop     b
763      pop     d
764      pop     h
765      ret
766
767      endif
768
769
770      Nibin:
771      call    Charin          ; return nibble in A register
772      rc
773      ani     07fh
774      sui     '0'
775      cpi     10
776      jc     Nibin$return    ; must be 0-9
777      adi     ('0'-'A'+10) and 0ffh
778      cpi     16
779      jc     Nibin$return    ; must be 10-15
780      lda     networkstatus
781      ori     F608

```

```

782 049F 324204      sta   networkstatus
783 04A2 3E00        mvi   a,0
784 04A4 37          stc
785 04A5 C9          ret
786
787                Nibin$return:
788 04A6 B7          ora   a
789 04A7 C9          ret
790
791                xChar$in:
792                                ; Get the first character using polled
793                                ; console I/O. Note that the rest of the
794                                ; message will be received using direct
795                                ; port I/O with interrupts disabled.
796                                ; OVERRUNS ARE NOT POSSIBLE USING THIS SCHEME
797 04A8 E5          push  h
798 04A9 C5          push  b
799 04AA 21BA04      lxi   h, Charin$return
800 04AD E5          push  h
801 04AE 48          mov   c,b
802 04AF 0600        mvi   b,0
803 04B1 214C04      lxi   h, chariotbl
804 04B4 09          dad   b
805 04B5 56          mov   d, m      ; Get the console number
806 04B6 2A4604      lhld  conin
807 04B9 E9          pchl          ; vector off
808
809                Charin$return:
810 04BA C1          pop   b

```

CP/M RMAC ASSEM 1.1 #016 MASTER NETWORK I/F MODULE

```

811 04BB E1          pop   h
812 04BC C9          ret
813
814
815                if      z80
816                char$in:
817
818                ;      Character input routine for network i/o
819                ;      using the ALTOS SIO ports at 125k baud
820                ;
821                ;      Z80 Version uses indirect port addresses loaded into register C
822                ;
823                ;      Entry : Slave number in register b
824                ;      Exit  : Character in register a
825                ;
826 04BD E5          push  h
827 04BE C5          push  b
828 04BF 214804      lxi   h, status$ports
829 04C2 48          mov   c, b
830 04C3 0600        mvi   b, 0      ; set (BC) = (b)
831 04C5 09          dad   b
832 04C6 4E          mov   c,m
833
834                ;      Now C contains the address of the correct status port
835
836 04C7 2E50        mvi   l, 80
837
838                inputloop1:
839 04C9 2D          dcr   l
840 04CA CADA04      jz    retout
841
842                ;      in      a,(c)
843 04CD ED78        db    0edh,78h
844
845 04CF E601        ani   01h      ; wait for RXready
846 04D1 CAC904      jz    inputloop1
847
848                ;      In the Altos system, data registers are one below status registers...
849
850 04D4 0D          dcr   c
851
852                ;      in      a,(c)
853 04D5 ED78        db    0edh,78h      ;get the character
854
855 04D7 C1          pop   b

```

```

856 04D8 E1      pop    h
857 04D9 C9      ret
858
859                retout:
860 04DA 37      stc                ;set carry => error flag
861 04DB C1      pop    b
862 04DC E1      pop    h
863 04DD C9      ret
864

```

CP/M RMAC ASSEM 1.1 #017 MASTER NETWORK I/F MODULE

```

865                else
866
867 char$in:
868
869                ; Character input routine for network I/O
870                ; using ALTOS SIO ports
871                ;
872                ; 8080 Version uses same nasty dispatch mechanism that the output
873                ; routine used
874                ;
875                ; Entry: B = Slave ID
876                ; Exit: A = character input
877
878                push    h
879                push    d
880                push    b
881                lxi    d,in0                ; HL = in0 + 17*slaveid
882                mov    l,b
883                mvi    h,0
884                xchg
885                dad    d
886                xchg
887                dad    h
888                dad    h
889                dad    h
890                dad    h
891                dad    d
892
893                mvi    c,80                ; load status retry count
894                pchl                ; dispatch
895
896 in0:
897                dcr    c
898                jz    retout                ; error return if retry timeout
899
900                in    Console4$status ; wait for RXready
901                ani    1
902                jz    in0
903
904                in    Console4$status-1    ; get the character
905                pop    b
906                pop    d
907                pop    h
908                ret
909
910 in1:
911                dcr    c
912                jz    retout
913
914                in    Console3$status
915                ani    1
916                jz    in1
917
918                in    Console3$status-1

```

CP/M RMAC ASSEM 1.1 #018 MASTER NETWORK I/F MODULE

```

919                pop    b
920                pop    d
921                pop    h
922                ret
923
924 in2:
925                dcr    c

```

```

926                jz      retout
927
928                in      Console2$status
929                ani     1
930                jz      in2
931
932                in      Console2$status-1
933                pop     b
934                pop     d
935                pop     h
936                ret
937                in3:
938                dcr     c
939                jz      retout
940
941                in      Printer2$status
942                ani     1
943                jz      in3
944
945                in      Printer2$status-1
946                pop     b
947                pop     d
948                pop     h
949                ret
950
951                retout:                ; error return (carry=1)
952                stc
953                pop     b
954                pop     d
955                pop     h
956                ret
957
958                endif
959
960
961                Netout:                ; C = byte to be transmitted
962                mov     a,d
963                add     c
964                mov     d,a
965                lda     BinaryASCII
966                ora     a
967                jnz     Charout        ; transmit byte in Binary mode
968                mov     a,c
969                push    psw
970                rar
971                rar
972                rar

```

CP/M RMAC ASSEM 1.1 #019 MASTER NETWORK I/F MODULE

```

973                rar
974                ani     0FH          ; Shift HI nibble to LO nibble
975                call    Nibout
976                pop     psw
977                ani     0FH
978                jmp     Nibout
979
980                Netin:                ; byte returned in A register
981                ; D = checksum accumulator
982                lda     BinaryASCII
983                ora     a
984                jz      ASCIIin
985                call    charin        ;receive byte in Binary mode
986                rc
987                jmp     chksin
988
989                ASCIIin:
990                call    Nibin
991                rc
992                add     a
993                add     a
994                add     a
995                add     a
996                push    psw
997                call    Nibin
998                rc
999                xthl

```

```

1000 0515 B4      ora    h
1001 0516 E1      pop    h
1002                chksin:
1003 0517 B7      ora    a
1004 0518 F5      push  psw
1005 0519 82      add    d                ; add & update checksum accum.
1006 051A 57      mov    d,a
1007 051B F1      pop    psw
1008 051C C9      ret
1009
1010                Msgin:
1011                ; HL = destination address
1012                ; E = # bytes to input
1012 051D CDF904   call   Netin
1013 0520 D8      rc
1014 0521 77      mov    m,a
1015 0522 23      inx    h
1016 0523 1D      dcr    e
1017 0524 C21D05   jnz    Msgin
1018 0527 C9      ret
1019
1020                Msgout:
1021                ; HL = source address
1022                ; E = # bytes to output
1023                ; D = checksum
1024                ; C = preamble character
1024 0528 1600     mvi    d,0
1025 052A CD6504   call   PreCharout
1026

```

CP/M RMAC ASSEM 1.1 #020 MASTER NETWORK I/F MODULE

```

1027                Msgoutloop:
1028 052D 4E      mov    c,m
1029 052E 23      inx    h
1030 052F CDDE04   call   Netout
1031 0532 1D      dcr    e
1032 0533 C22D05   jnz    Msgoutloop
1033 0536 C9      ret
1034
1035                ; Network Initialization
1036
1037                nwinit:
1038
1039                ; device initialization, as required
1040
1041
1042 0537 3E47     mvi    a,047h          ;sets up CTC for baud rate of 125k
1043 0539 D331     out    031h
1044
1045                if nmbslvs ge 3 ;initialize only the ports that are needed
1046                out    030h
1047                endif
1048
1049                if nmbslvs ge 4
1050                out    032h
1051                endif
1052
1053 053B 3E01     mvi    a,1              ;count of one => max speed
1054 053D D331     out    031h
1055
1056                if nmbslvs ge 3
1057                out    030h
1058                endif
1059
1060                if nmbslvs ge 4
1061                out    032h
1062                endif
1063
1064
1065                ; Find address of XI0S console output routine
1066
1067 053F 2A0100   lhld   0001h          ; get warmstart entry in the XI0S jump table
1068 0542 23      inx    h
1069 0543 5E      mov    e, m
1070 0544 23      inx    h
1071 0545 5E      mov    d, m
1072 0546 210600   lxi    h, 0006h      ; Offset for conin routine
1073 0549 19      dad    d

```



```

1074 054A 224604      shld  conin      ; save the address
1075 054D AF          xra    a          ; return code is 0=success
1076 054E C9          ret
1077
1078
1079                ;      Network Status
1080

```

CP/M RMAC ASSEM 1.1 #021 MASTER NETWORK I/F MODULE

```

1081                nwstat:                ; C = Slave #
1082 054F 0600          mvi    b,0
1083 0551 214204        lxi    h,networkstatus
1084 0554 09            dad    b
1085 0555 7E            mov    a,m
1086 0556 47            mov    b,a
1087 0557 E6F5          ani    not (msgerr+msgovr)
1088 0559 77            mov    m,a
1089 055A 78            mov    a,b
1090 055B C9            ret
1091
1092
1093                ;      Return Configuration Table Address
1094
1095                cfgadr:
1096 055C 215603        lxi    h,configtbl
1097 055F C9            ret
1098
1099
1100                ;      Send Message on Network
1101
1102                sndmsg:                ; DE = message addr
1103                                ; C = Slave #
1104 0560 41            mov    b,c
1105 0561 3E0A          mvi    a,max$retries ; A = max$retries
1106
1107                send:
1108 0563 F5            push   psw
1109
1110                if      mutexout
1111
1112                ;      Use mutual exclusion if it is possible for some unsolicited input
1113                ;      to stomp on your output (This is nice is you;re running some sort
1114                ;      of multi-drop protocol)
1115
1116                push   b
1117                push   d
1118                mvi    c,readq
1119                lxi    d,UQCBMXSXmitq
1120                call   bdos          ; obtain mutual exclusion token
1121                pop    d
1122                pop    b
1123                endif
1124
1125 0564 EB            xchg
1126 0565 E5            push   h
1127 0566 F3            di          ; disable interrupts to avoid underrun
1128 0567 0E05          mvi    c,ENQ
1129 0569 CD6804        call   Charout      ; send ENQ
1130 056C CDA005        call   getACK       ; won't return on an error
1131 056F 1E05          mvi    e,5
1132 0571 0E01          mvi    c,SOH
1133 0573 CD2805        call   Msgout       ; send SOH FMT DID SID FNC SIZ
1134 0576 AF            xra    a

```

CP/M RMAC ASSEM 1.1 #022 MASTER NETWORK I/F MODULE

```

1135 0577 92            sub    d
1136 0578 4F            mov    c,a
1137 0579 CDDE04        call   Netout       ; send HCS (header checksum)
1138 057C CDA005        call   getACK       ; won't return on an error
1139 057F 2B            dcx   h
1140 0580 5E            mov    e,m
1141 0581 23            inx   h
1142 0582 1C            inr   e
1143 0583 0E02          mvi    c,STX

```

```

1144 0585 CD2805      call    Msgout          ; send STX DB0 DB1 ...
1145 0588 0E03        mvi     c,ETX
1146 058A CD6504      call    PreCharout     ; send ETX
1147 058D AF          xra     a
1148 058E 92          sub     d
1149 058F 4F          mov     c,a
1150 0590 CDDE04      call    Netout         ; send CKS
1151 0593 0E04        mvi     c,EOT
1152 0595 CD6504      call    PreCharout     ; send EOT
1153 0598 CDA005      call    getACK         ; won't return on an error
1154 059B D1          pop     d              ; discard message address
1155 059C F1          pop     psw           ; discard retry counter
1156
1157                if     mutexout
1158                call   release$MX
1159                endif
1160
1161 059D FB          ei              ; return from suspended animation
1162 059E AF          xra     a
1163 059F C9          ret              ; A = 0, successful send message
1164
1165                getACK:
1166 05A0 CDBD04      call    Charin
1167 05A3 DAAB05      jc     getACK$timeout ; receive timeout-->start error recovery
1168 05A6 E67F        ani     7fh
1169 05A8 D606        sui     ACK
1170 05AA C8          rz
1171
1172                getACK$timeout:
1173 05AB D1          pop     d              ; discard return address
1174
1175                if     mutexout
1176                push   b
1177                call   release$MX
1178                pop    b
1179                endif
1180
1181 05AC D1          pop     d              ; DE = message address
1182 05AD F1          pop     psw           ; A = retry count
1183 05AE 3D          dcr     a
1184 05AF C26305      jnz    send          ; continue if retry count non-zero
1185 05B2 3D          dcr     a              ; else-->we're dead-->A = 0ffh
1186 05B3 C9          ret              ; failed to send message
1187
1188                if     mutexin or mutexout

```

CP/M RMAC ASSEM 1.1 #023 MASTER NETWORK I/F MODULE

```

1189
1190                release$MX:          ; send back requester transmit MX message
1191                mvi     c,writeq
1192                lxi     d,UQCBMXSXmitq
1193                jmp     bdos
1194                endif
1195
1196                ;      Receive Message from Network
1197
1198                rcvmsg:          ; DE = message addr
1199                ;      C = Slave #
1200 05B4 41          mov     b,c
1201
1202                receive:
1203 05B5 EB          xchg
1204 05B6 E5          push   h
1205 05B7 CDBF05      call   get$ENQ
1206
1207                ;      a return to this point indicates an error
1208
1209                receive$retry:
1210 05BA FB          ei              ; re-enable other processes
1211
1212                if     mutexin
1213                push   b
1214                call   release$MX
1215                pop    b
1216                endif
1217

```

```

1218 05BB D1          pop    d
1219 05BC C3B505      jmp    receive
1220
1221                get$ENQ:                ; get first character of message using
1222                                ; polled console I/O
1223 05BF CDA804        call   xCharin
1224 05C2 DABF05        jc     get$ENQ
1225 05C5 E67F          ani    7fh
1226 05C7 FE05          cpi    ENQ                ; Start of Message ?
1227 05C9 C2BF05        jnz   get$ENQ
1228
1229                if     mutexin
1230
1231                ; Don't get too involved with receiving a message if some other
1232                ; NETWRKIF process is going to stomp you by sending a message along
1233                ; the same line
1234
1235                push   b
1236                push   h
1237                mvi    c,readq
1238                lxi    d,UQCBMXSXmitq
1239                call   bdos
1240                pop    h
1241                pop    b
1242                endif

```

CP/M RMAC ASSEM 1.1 #024 MASTER NETWORK I/F MODULE

```

1243
1244 05CC 0E06          mvi    c,ACK
1245 05CE F3            di                    ; requester in gear now serve only him
1246
1247 05CF CD6804        call   charout          ; send ACK to requester, allowing transmit
1248 05D2 CDBD04        call   Charin
1249 05D5 D8            rc
1250 05D6 E67F          ani    7fh
1251 05D8 FE01          cpi    SOH
1252 05DA C0            rnz
1253 05DB 57            mov    d,a              ; initialize the HCS
1254 05DC 1E05          mvi    e,5
1255 05DE CD1D05        call   Msgin
1256 05E1 D4F904        cnc    Netin
1257 05E4 D8            rc
1258 05E5 7A            mov    a,d
1259 05E6 B7            ora    a
1260 05E7 C21406        jnz   sendNAK          ; jmp & send NAK if HCS <= 0
1261 05EA 0E06          mvi    c,ACK
1262 05EC CD6804        call   Charout
1263 05EF CDBD04        call   Charin
1264 05F2 D8            rc
1265 05F3 E67F          ani    7fh
1266 05F5 FE02          cpi    STX
1267 05F7 C0            rnz
1268 05F8 57            mov    d,a              ; initialize the CKS
1269 05F9 2B            dcx   h
1270 05FA 5E            mov    e,m
1271 05FB 23            inx   h
1272 05FC 1C            inr   e
1273 05FD CD1D05        call   msgin
1274 0600 D4BD04        cnc    Charin
1275 0603 D8            rc
1276 0604 E67F          ani    7fh
1277 0606 FE03          cpi    ETX
1278 0608 C0            rnz
1279 0609 82            add   d
1280 060A 57            mov    d,a
1281 060B CDF904        call   Netin          ; get Checksum byte
1282 060E D8            rc
1283 060F 7A            mov    a,d
1284 0610 B7            ora    a                ; should be zero
1285 0611 CA1906        jz     sendACK          ; jump if checksum OK
1286
1287                sendNAK:                ; else-->refuse the message
1288 0614 0E15          mvi    c,NAK
1289 0616 C36804        jmp    Charout          ; send NAK and return to receive$retry
1290
1291                sendACK:                ; come here if message was received properly

```

```

1292 0619 CDBD04      call   Charin          ; get EOT
1293 061C D8          rc
1294 061D E67F        ani    7fh
1295 061F FE04        cpi    EOT
1296 0621 C0          rnz

```

CP/M RMAC ASSEM 1.1 #025 MASTER NETWORK I/F MODULE

```

1297 0622 0E06        mvi    c,ACK
1298 0624 CD6804       call   Charout         ; send ACK if checksum ok
1299 0627 D1           pop    d               ; discard return address
1300 0628 D1           pop    d               ; discard message address
1301 0629 FB           ei                    ; Dispense with the Rip Van Winkle act
1302
1303                  if    mutexin
1304                  call   release$MX
1305                  endif
1306
1307 062A AF           xra    a
1308 062B C9           ret
1309
1310
1311                restore:
1312
1313                ;      This routine allows N copies of NtwrkIPx to run reentrantly.
1314                ;      It takes the values that were pre-initialized in the process
1315                ;      descriptor and later saved on the stack and loads them into
1316                ;      the registers, leaving the stack image untouched. All variables
1317                ;      intrinsic to the process therefore always reside on the
1318                ;      process-dependent stack
1319
1320 062C F3           di                    ; this is a real critical region
1321 062D E1           pop    h
1322 062E 224006       shld   rtnadr
1323 0631 E1           pop    h
1324 0632 D1           pop    d
1325 0633 C1           pop    b
1326 0634 F1           pop    psw
1327 0635 F5           push   psw
1328 0636 C5           push   b
1329 0637 D5           push   d
1330 0638 E5           push   h
1331 0639 E5           push   h
1332 063A 2A4006       lhld   rtnadr
1333 063D E3           xthl
1334 063E FB           ei
1335 063F C9           ret
1336
1337 0640                rtnadr: ds    2
1338
1339                if    WtchDg
1340
1341                ;      WatchDog Timer Process
1342                ;      This process needs adjunct processes to handle the timeout flags
1343                ;      that it sets. They might possibly abort the offending NtwrkIPx
1344                ;      process, recreate it, and allow it to re-initialize its queues
1345
1346                WatchDog:
1347                mvi    c,Delay
1348                lxi    d,60          ; delay for 1 second
1349                call   bdos
1350                lhld   WatchDogTime

```

CP/M RMAC ASSEM 1.1 #026 MASTER NETWORK I/F MODULE

```

1351                inx    h
1352                shld   WatchDogTime
1353                lxi    h,WatchDogTable-5
1354                mvi    c,NmbSlvs
1355
1356                WatchDogLoop:
1357                lxi    d,0005h
1358                dad    d
1359                mov    a,m
1360                ora    a
1361                jz     WatchDogDec

```

```

1362          inx      h
1363          ana      m
1364          dcx      h
1365          jnz      WatchDogDec      ; waiting & timeout set
1366          push     h                  ; save HL -> WDT.waiting
1367          inx      h
1368          inx      h
1369          di
1370          mov      e,m
1371          inx      h
1372          mov      d,m
1373          ei
1374          lhld     WatchDogTime
1375          mov      a,l
1376          sub      e
1377          mov      l,a
1378          mov      a,h
1379          sbb      d
1380          mov      h,a
1381          mvi      a,10              ; # seconds since started Charin
1382          sub      l
1383          mvi      a,0
1384          sbb      h
1385          pop      h
1386          jnc      WatchDogDec
1387          push     h
1388          inx      h
1389          mvi      m,0ffh           ; WDT.timeout = 0ffh
1390          inx      h
1391          inx      h
1392          inx      h
1393          push     b
1394          mov      e,m              ; E = Flag #
1395          mvi      c,Flagset
1396          call     bdos
1397          pop      b
1398          pop      h
1399
1400 WatchDogDec:
1401          dcr      c
1402          jnz      WatchDogLoop
1403
1404          jmp      WatchDog

```

CP/M RMAC ASSEM 1.1 #027 MASTER NETWORK I/F MODULE

```

1405          endif
1406
1407
1408          ;      Setup code for Network Interface Procedures
1409
1410          Setup:
1411          0642 F5      push     psw          ;create stack image of all reentrant variables
1412          0643 C5      push     b
1413          0644 D5      push     d
1414          0645 E5      push     h
1415          0646 CD3705  call     nwinit
1416
1417          if      mutexin or mutexout
1418          mvi      c,makeq          ; make the mutual exclusion queue
1419          lxi      d,QCBMXSXmitq
1420          call     bdos
1421
1422          mvi      c,writq          ; leave a token in the queue
1423          lxi      d,UQCBMXSXmitq
1424          call     bdos
1425          endif
1426
1427          if      WtchDg
1428          lxi      d,WatchDogPD      ;since this process is linked to all other
1429                                     ;NtwrkIPx processes, creating it creates all
1430                                     ;of the others
1431          mvi      c,createp
1432          call     bdos
1433
1434          else
1435

```

```

1436                                     if      NmbSlvs GE 2
1437 0649 11AC01                          lxi     d,NtwrkIP1      ;this will create all the other NtwrkIPx
1438                                     ;processes if there's no watchdog
1439 064C 0E90                             mvi     c,createp
1440 064E CD5004                          call    bdos
1441                                     endif
1442                                     endif
1443
1444 0651 0E8E                             mvi     c,dsptch       ;give everything a chance to create its queues
1445 0653 CD5004                          call    bdos
1446
1447 0656 0E9A                             mvi     c,sydatad
1448 0658 CD5004                          call    bdos
1449 065B 110900                          lxi     d,9
1450 065E 19                               dad     d
1451 065F 115603                          lxi     d,configtbl
1452 0662 73                               mov     m,e
1453 0663 23                               inx     h
1454 0664 72                               mov     m,d             ; sysdatpage(9&10) = co.configtbl
1455                                     ; filling in the config tbl address is the
1456                                     ; the server processes' cue to start
1457
1458                                     if      modem

```

CP/M RMAC ASSEM 1.1 #028 MASTER NETWORK I/F MODULE

```

1459                                     ;      Initialize the modem
1460
1461                                     mvi     c,CR
1462                                     mvi     b,slvmodem
1463                                     call    Charout
1464                                     mvi     c,'Z'
1465                                     call    Charout
1466                                     mvi     c,CR
1467                                     call    Charout
1468
1469      WtSpace:
1470                                     call    Charin
1471                                     jc      SetupDone
1472                                     ani     07fh
1473                                     cpi     ' '
1474                                     jnz    WtSpace
1475                                     mvi     c,'A'
1476                                     call    Charout
1477
1478      SetupDone:
1479                                     endif
1480
1481 0665 E1                               pop     h
1482 0666 D1                               pop     d
1483 0667 C1                               pop     b
1484 0668 F1                               pop     psw
1485
1486                                     ; Network Interface Reentrant Procedure
1487
1488      Init:
1489 0669 F5                               push    psw             ; A = network i/f console dev #
1490 066A C5                               push    B               ; BC= buffer address
1491 066B D5                               push    D               ; DE= UQCB ntwrk queue out
1492 066C E5                               push    H               ; HL= UQCB ntwrk queue in
1493 066D 5E                               mov     e,m
1494 066E 23                               inx     h
1495 066F 56                               mov     d,m
1496 0670 0E86                             mvi     c,makeq
1497 0672 CD5004                          call    bdos           ; make the ntwrk queue in
1498 0675 CD2C06                          call    restore
1499 0678 EB                               xchg
1500 0679 5E                               mov     e,m
1501 067A 23                               inx     h
1502 067B 56                               mov     d,m
1503 067C 0E86                             mvi     c,makeq
1504 067E CD5004                          call    bdos           ; make the ntwrk queue out
1505
1506      Loop:
1507 0681 CD2C06                          call    restore
1508 0684 50                               mov     d,b
1509 0685 59                               mov     e,c

```

```

1510
1511 0686 4F          mov    c,a
1512 0687 CDB405      call   rcvmsg

```

CP/M RMAC ASSEM 1.1 #029 MASTER NETWORK I/F MODULE

```

1513
1514 068A CD2C06      call   restore
1515 068D EB          xchg
1516 068E 0E8B        mvi   c,writeq
1517 0690 CD5004      call   bdos
1518
1519 0693 CD2C06      call   restore
1520 0696 0E89        mvi   c,readq
1521 0698 CD5004      call   bdos
1522
1523 069B CD2C06      call   restore
1524 069E 50          mov    d,b
1525 069F 59          mov    e,c
1526
1527 06A0 4F          mov    c,a
1528 06A1 CD6005      call   sndmsg
1529
1530 06A4 C38106      jmp    Loop
1531
1532 06A7              end

```

CP/M RMAC ASSEM 1.1 #030 MASTER NETWORK I/F MODULE

```

ACK          0006 612# 1169 1244 1261 1297
ACTIVE       0010 587#
ASCIIIN     0507 984 989#
BDOS        0450 44 622# 1120 1193 1239 1349 1396 1420 1424 1432
            1440 1445 1448 1497 1504 1517 1521
BDOSADR     0000 51# 58# 623
BINARYASCII 043E 541# 965 982
BUFFERQ0    00A6 47 89 118 138#
BUFFERQ1    0250 179 208 228#
BUFFERQI0ADDR 0084 116 117#
BUFFERQI1ADDR 022E 206 207#
BUFFERQ00ADDR 00A4 134 135#
BUFFERQ01ADDR 024E 224 225#
CFGADR      055C 1095#
CHARIN      04BD 771 816# 867# 985 1166 1248 1263 1274 1292 1470
CHARINRETURN 04BA 799 809#
CHARIOTBL   044C 578# 803
CHAROUT     0468 631 635 643# 692# 967 1129 1247 1262 1289 1298
            1463 1465 1467 1476
CHKSIN      0517 987 1002#
CONFIGTBL   0356 493# 1096 1451
CONIN       0446 553# 806 1074
CONSOLE2STATUS 002D 564# 574 743 744 749 928 932
CONSOLE3STATUS 002F 563# 573 731 732 737 914 918
CONSOLE4STATUS 002B 562# 572 719 720 725 900 904
CR          000D 614# 1461 1466
CREATEP     0090 601# 1431 1439
DEBUG       0000 28# 37
DELAY       008D 599# 1347
DSPTCH      008E 600# 1444
ENQ         0005 611# 1128 1226
EOT         0004 610# 1151 1295
ETX         0003 609# 1145 1277
FALSE       0000 23# 24 28 29 31 33 34
FLAGSET     0085 595# 1395
GETACK      05A0 1130 1138 1153 1165#
GETACKTIMEOUT 05AB 1167 1172#
GETENQ      05BF 1205 1221# 1224 1227
INIT        0669 190 281 366 1488#
INPUTLOOP1  04C9 838# 846
LF          000A 613#
LOOP        0681 1506# 1530
MAKEQ       0086 596# 1418 1496 1503
MAXRETRIES  000A 555# 1105
MODEM       0000 29# 1458
MSGERR      0008 588# 781 1087
MSGIN       051D 1010# 1017 1255 1273

```

MSGOUT	0528	1020#	1133	1144							
MSGOUTLOOP	052D	1027#	1032								
MSGOVR	0002	590#	1087								
MUTEXIN	0000	33#	471	1188	1212	1229	1303	1417			
MUTEXOUT	0000	34#	471	1110	1157	1175	1188	1417			
NAK	0015	615#	1288								
NETIN	04F9	980#	1012	1256	1281						

CP/M RMAC ASSEM 1.1		#031	MASTER NETWORK I/F MODULE								
NETOUT	04DE	961#	1030	1137	1150						
NETWORKSTATUS	0442	547#	780	782	1083						
NIBATOF	045F	628	632#								
NIBIN	0486	770#	990	997							
NIBINRETURN	04A6	776	779	787#							
NIBOUT	0454	626#	975	978							
NMSLVS	0002	39#	56#	148	151	239	242	330	419	496	505
		512	518	524	529	534	1045	1049	1056	1060	1354
		1436									
NMSG	0001	502#									
NTWRK	0004	589#									
NTWRKIP0	0002	65#									
NTWRKIP1	01AC	149#	420	1437							
NTWRKIS0	0036	41	69	93#							
NTWRKIS1	01E0	159	183#								
NTWRKRXYDY	0001	591#									
NTWRKTXRDY	0080	586#									
NWINIT	0537	1037#	1415								
NWSTAT	054F	1081#									
OUTPUTLOOP	0474	666#	676								
POLL	0083	603#									
PRECHAROUT	0465	637#	1025	1146	1152						
PRINTER2	0010	617#									
PRINTER2STATUS	0029	565#	575	755	756	761	941	945			
QCBNTWRKQI0	0066	102#	115								
QCBNTWRKQI1	0210	192#	205								
QCBNTWRKQ00	0086	120#	133								
QCBNTWRKQ01	0230	210#	223								
RCVMSG	05B4	1198#	1512								
READQ	0089	597#	1118	1237	1520						
RECEIVE	05B5	1202#	1219								
RECEIVERETRY	05BA	1209#									
RESTORE	062C	1311#	1498	1507	1514	1519	1523				
RETOUT	04DA	840	859#	898	912	926	939	951#			
RTNADR	0640	1322	1332	1337#							
SEND	0563	1107#	1184								
SENDACK	0619	1285	1291#								
SENDNAK	0614	1260	1287#								
SETUP	0642	100	1410#								
SLAVE1	040A	508	525#								
SLAVE1STK	0374	506#									
SLAVE1STKLEN	0096	503#	507	514	520						
SNDMSG	0560	1102#	1528								
SOH	0001	607#	1132	1251							
STATUSPORTS	0448	571#	658	828							
STX	0002	608#	1143	1266							
SYDATAD	009A	602#	1447								
TRUE	FFFF	24#	26								
UQCBNTWRKQI0	0080	45	87	114#							
UQCBNTWRKQI1	022A	177	204#								
UQCBNTWRKQ00	00A0	46	88	132#							
UQCBNTWRKQ01	024A	178	222#								
WRITEQ	008B	598#	1191	1422	1516						
WTCHDG	0000	31#	414	1339	1427						

CP/M RMAC ASSEM 1.1		#032	MASTER NETWORK I/F MODULE								
XCHARIN	04A8	791#	1223								
Z80	FFFF	26#	567	642	815						

Listing E-2: Server Network I/F Module

Appendix F

A CP/NET System for use with ULCnet

F.1 Overview of ULCnet

ULCnet® (Universal Low Cost Network) is a local area network system designed specifically for microcomputers in the CP/M and MP/M II operating system environments. ULCnet was introduced by Orange Compuco, Inc. in June 1982 as a low cost method of sharing resources and data among microcomputers of varying manufacture and architecture. ULCnet, in combination with CP/NET, creates a cost effective method for the development of shared data base applications among single user microcomputers. ULCnet architecture readily supports CP/NET implementation.

The ULCnet connector adaptor box can be connected to any computer that has a spare RS-232 port. ULCnet employs a multidrop topology with carrier sense, multiple-access design. Contention between network nodes is arbitrated using a full-duplex collision detection mechanism.

ULCnet is available to OEMs on a private label basis and through licensing. Keybrook Business Systems, Inc., Hayward, California, a licensee of ULCnet, produces the FileServer™ system. This system uses CP/NET to drive ULCnet. For more information on ULCnet, contact

Orange Compuco, Inc.
17801-G South East Main Street
Irvine, California 92714
(714) 957-8075

Orange Compuco distributes ULCnet connector adaptor hardware with a variety of release software, including the example programs in this appendix. In addition, Orange Compuco provides documentation detailing the installation and operation of ULCnet and logical structure of the data-link layer software. This documentation includes

- details on the installation and configuration of ULCnet
- a detailed description of the linkage between the proprietary data-link software and the user-definable Network I/O Drivers (NIOD)
- a detailed description of the interface between higher-level software and data-link software
- a description of the data-link interface (DLIF) between the data-link software and higher-level layers

F.2 Customizing a ULCnet SNIOS for the Requester

The CP/NET requester listing, SNIOS for ULCnet, that appears at the end of this section, is contained in a file called ULCNIO.SASM on the CP/NET release disk and is designed to run ULCnet in a polled environment on a Xerox® 820 computer, now called the Xerox R820-IIS. The listing uses the ULCnet short format. This means that virtual circuit numbers must be agreed upon before the requester and the server can communicate. This version assumes that the server ID is always 0, and that up to four requesters, ID 1 through 4, are on the network. The virtual circuit number and the requester ID are always the same.

This SNIOS combines the two sections of the ULCnet protocol that are user configurable, the data-link interface (DLIF) and the network I/O drivers (NIOD). The DLIF acts as a transport layer between the NDOS and the data-link routines. The NIOD contains the physical device drivers use to communicate with the ULCnet network adaptor box. The bulk of the data-link protocol is contained in a module called PBMAIN.REL. This module is proprietary to Orange Compuco, and is therefore distributed only in REL file format by Orange Compuco.

When the NDOS instructs the SNIOS to send a message, the SNIOS first converts the CP/NET message format into ULCnet short format. The SNIOS then calls the TRANSMIT routine in PBMAIN to send the message, followed by the GETTCODE routine to discover the status of the message. If the send was successful, the SNIOS returns to the NDOS. If it was not successful, the SNIOS continues to try to send the message. No timeout is included in this routine to halt transmission.

To receive a message, the SNIOS calls RECEIVE, followed by GETRCODE to check the status of the message. If the status shows success, the message is converted from ULCnet format back into CP/NET format and returns to the NDOS. If the status shows an error, the SNIOS attempts to receive the message again.

To modify the SNIOS for a requester other than a Xerox 820, follow these steps:

1. Decide whether to make the requester operate in a polled or interrupt-driven environment. If you want interrupts, set the INTERRUPTS assembly switch to TRUE, and link the module using IPBMAIN instead of PBMAIN.
2. If your ULCnet connector adaptor has been modified for self clocked operation, set the assembly switch SLFCLKD to TRUE. Application notes detailing how to modify the connector adaptor for self-clocked operation are available from Orange Compuco.
3. Determine your requester's transmission speed capabilities. Set the baud rate masks BAUDSL and BAUDSH to reflect these values. Enter values for the requester's baud rate generator into the table BAUDTBL.
4. Modify the port numbers for the baud rate generator and the UART to reflect those used by your requester.
5. Modify the NIOD to run on your requester. The NIOD is currently set up to drive a Z80 SIO chip. If your requester has an SIO, it needs little modification. The routine PGMUART, which sets up the network port for ULCnet operation, might have to be modified. In an interrupt driven system, interrupt vectors must be set up here.
6. Assemble and link the SNIOS by performing

```
A>RMAC ULCN10S
A>LINK SNIOS=ULCN10S,PBMAIN[0S]
```

If the requester is interrupt-driven, perform

```
A>LINK SNIOS=ULCN10S,IPBKAIN[0S]
```

to link the module. The module is then ready for installation on the CP/NET requester system disk.

```
CP/M RMAC ASSEM 1.1      #001  REQUESTER NETWORK I/O SYSTEM FOR ULCNET

1          title 'Requester Network I/O System for ULCnet'
2          page   54
3
4          ;*****
5          ;*****
6          ;**
7          ;**          SNIOS FOR ULCNET          **
8          ;**
9          ;*****
10         ;*****
11
12         ;          Developed jointly by:
13         ;
14         ;          Digital Research, Inc.
15         ;          P.O. Box 579
```

```

16      ; Pacific Grove, CA 93950
17      ; and
18      ; Keybrook Business Systems, Inc.
19      ; 2035 National Avenue
20      ; Hayward, CA 94545
21
22      ; This SNIOS was written for a Xerox 820 attached to Orange
23      ; Compuco's ULCnet network adaptor. This module transports
24      ; messages between the NDOS and the low-level data-link software
25      ; provided by Orange Compuco. It also contains the physical drivers
26      ; usually contained in the NIOD module. This version is not
27      ; interrupt-driven and must be linked with PBMAIN.REL.
28
29
30
31 0000 = false equ 0
32 FFFF = true equ not false
33
34 0000 = interrupts equ false ; false=pollled, true=interrupt-driven
35 FFFF = netstats equ true ; switch to gather network statistics
36 FFFF = slfclkd equ true ; supports self-clocked operation
37
38 ; Linkage information
39
40 public setbaud,xmit,rcv,initu ; NIOD routines called by IPBMAIN
41 public inituart,pgmuart
42 public chkstat,netidle,initrcv
43 public wait,restuart,csniod
44 public dsblxmit
45 public dllbau,netadr
46
47 if interrupts
48 public enblrcv,dsblrcv
49 endif
50
51 extrn transmit,receive ; IPBMAIN routines and objects
52 extrn gettcode,getrcode
53 extrn csdll,dllon,regshrt
54 extrn terrcnt,parcntr,ovrcntr

```

CP/M RMAC ASSEM 1.1 #002 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

55 extrn frmcntr,incntr
56
57 if interrupts
58 extrn rtmochk ; IPBMAIN interrupt routines
59 extrn dlisr,reisr,niisr
60 endif
61
62
63 ; Hardware definitions for the Z80-SIO channel A - For the Xerox 820.
64
65 0003 = baudsl equ 03h ; Usable baud rates: 9600, 19.2K asynch.,
66 002A = baudsh equ 2ah ; 76.8K, 153.6K, 307.2K self-clocked
67
68 ; baud rate capability mask
69 2A03 = bauds equ (baudsh*100h)+baudsl
70
71 0000 = baudgen equ 0 ; External baud rate generator register
72 0006 = siocmd equ 6 ; Command/Mode register
73 0006 = siostat equ 6 ; Status register
74 0004 = sioxmit equ 4 ; Transmit register
75 0004 = siorecv equ 4 ; Receive register
76
77 0002 = xrdybit equ 2 ; Transmit buffer empty status bit
78 0004 = xrdysk equ 4 ; transmit buffer empty status mask
79 0000 = rrdybit equ 0 ; Receive buffer full status bit
80 0001 = rrdysk equ 1 ; receive buffer full status mask
81 0003 = carbit equ 3 ; Net Idle detect bit position
82 0008 = carmsk equ 8 ; Net Idle detect mask
83 0030 = errrst equ 030h ; Error flag reset
84 0070 = errbits equ 070h ; Error bit position mask
85 0004 = pbit equ 4 ; Parity error bit position
86 0010 = pmsk equ 10h ; parity error mask
87 0005 = obit equ 5 ; Overrun error bit position
88 0020 = omsk equ 20h ; overrun error mask
89 0006 = fbit equ 6 ; Framing error bit position

```

```

90 0040 = fmsk equ 40h ; framing error mask
91 0003 = selfbit equ 3 ; Self clock bit position
92 0008 = selfmsk equ 8 ; self clock bit mask
93 00EA = dtron equ 0eah ; Turn on DTR
94 006A = dtroff equ 06ah ; Turn off DTR
95 00C1 = enarcv equ 0c1h ; Enable receive-clock
96 00C0 = disrcv equ 0c0h ; Disable receive clock
97 000F = enaslf equ 00fh ; Enable Self-clock mode
98 004F = disslf equ 04fh ; Disable Self-clock mode
99
100 ; SIO Mode 2 interrupts vector table
101
102 FF08 = siov4 equ 0ff08h ; SIO port A xmit buffer empty
103 FF0A = siov5 equ 0ff0ah ; SIO port A external status change
104 FF0C = siov6 equ 0ff0ch ; SIO port A receive
105 FF0E = siov7 equ 0ff0eh ; SIO port A special receive condition
106
107
108 ; Message Buffer Offsets

```

CP/M RMAC ASSEM 1.1 #003 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

109
110 0000 = fmt equ 0 ; format
111 0001 = did equ fmt+1 ; destination ID
112 0002 = sid equ did+1 ; source ID
113 0003 = fnc equ sid+1 ; server function number
114 0004 = siz equ fnc+1 ; size of message (normalized to 0)
115 0005 = msg equ siz+1 ; message
116 0106 = buf$len equ msg+257 ; length of total message buffer
117
118 ; ULCnet Packet Offsets
119
120 0000 = ulc$fmt equ 0 ; packet format
121 0001 = ulc$v$circ equ ulc$fmt+1 ; virtual circuit number
122 0002 = ulc$len$lo equ ulc$v$circ+1 ; low order of length
123 0003 = ulc$len$hi equ ulc$len$lo+1 ; high order of length
124 0004 = ulc$fnc equ ulc$len$hi+1 ; start of message: function code
125 0005 = ulc$msg equ ulc$fnc+1 ; CP/NET message
126
127 ; Network Status Byte Equates
128
129 0010 = active equ 0001$0000b ; slave logged in on network
130 0002 = rcvrr equ 0000$0010b ; error in received message
131 0001 = senderr equ 0000$0001b ; unable to send message
132
133
134
135 CSEG
136 0005 = BDOS equ 0005h
137
138 NIOS:
139 public NIOS
140
141 ; Jump vector for SNIOS entry points
142
143 0000 C3E100 jmp ntwrkinit ; network initialization
144 0003 C3EE00 jmp ntwrksts ; network status
145 0006 C3F600 jmp cnfgtbladr ; return config table addr
146 0009 C30401 jmp sendmsg ; send message on network
147 000C C32001 jmp receivmsg ; receive message from network
148 000F C3FA00 jmp ntwrkerror ; network error
149 0012 C30301 jmp ntwrkboot ; network warm boot
150
151
152 0001 = rqstr$id equ 1 ; requester ID: must be between 1 and 4
153 004B = fmt$byte equ 4bh ; format byte: short format with data-link
154 ; acknowledge, 153.6K baud self-clocked
155
156 DSEG
157
158 ; Transport Layer Data
159
160 network$error$msg:
161

```

```
162 0000 0D0A          db      0dh,0ah
```

```
CP/M RMAC ASSEM 1.1    #004    REQUESTER NETWORK I/O SYSTEM FOR ULCNET
```

```
163 0002 4E6574776F      db      'Network Error'
164 000F 0D0A          db      0dh,0ah
165 0011 24            db      '$'
166
167
168                ; Requester Configuration Table
169
170            configtbl:
171            Network$status:
172
173 0012                ds      1                ; network status byte
174 0013 01            db      rqstr$id        ; slave processor ID number
175 0014                ds      2                ; A: Disk device
176 0016                ds      2                ; B: "
177 0018                ds      2                ; C: "
178 001A                ds      2                ; D: "
179 001C                ds      2                ; E: "
180 001E                ds      2                ; F: "
181 0020                ds      2                ; G: "
182 0022                ds      2                ; H: "
183 0024                ds      2                ; I: "
184 0026                ds      2                ; J: "
185 0028                ds      2                ; K: "
186 002A                ds      2                ; L: "
187 002C                ds      2                ; M: "
188 002E                ds      2                ; N: "
189 0030                ds      2                ; O: "
190 0032                ds      2                ; P: "
191 0034                ds      2                ; console device
192 0036                ds      2                ; list device:
193
194                ; List Buffer Data
195
196 0038                ds      1                ;      buffer index
197
198 0039 00            db      0                ;      FMT
199 003A 00            db      0                ;      DID
200 003B 01            db      rqstr$id        ;      SID
201 003C 05            db      5                ;      FNC
202 003D                ds      1                ;      SIZ
203 003E                ds      1                ;      MSG(0) List number
204 003F                ds      128           ;      MSG(1) ... MSG(128)
205
206
207                ; ULCnet Data Definitions
208
209 00BF                netadr: ds      3                ;ULCnet network address
210 00C2                dllbau: ds      2                ;baud rate mask
211
212 0016 =                timeval equ      22                ; WAIT routine time constant
213                ; 12 for 2.5 megahertz Z80
214                ; 22 for 4.0 megahertz Z80
215
216 00C4 FF            curbaud db      0ffh                ; Current baud rate
```

```
CP/M RMAC ASSEM 1.1    #005    REQUESTER NETWORK I/O SYSTEM FOR ULCNET
```

```
217
218
219                ; table to convert baud number codes
220                ;      into a bit mask
221
222 00C5 0102040810btbl: db      1,2,4,8,16,32,64,128
223
224
225                baudtbl:                ; async baud rate table
226
227 00CD 0E            db      0eh                ; 9600 Baud
228 00CE 0F            db      0fh                ; 19200
229
230                scbaudt:                ; self-clock baud rate table
231
```

```

232 00CF 00          db      0          ; 62500 Baud - Not implemented
233 00D0 0D          db      0dh         ; 76800 Baud
234 00D1 00          db      0          ; 125000 Baud - Not implemented
235 00D2 0E          db      0eh         ; 153600 Baud
236 00D3 00          db      0          ; 250000 Baud - Not implemented
237 00D4 0F          db      0fh         ; 307200 Baud
238
239                if      interrupts
240                sioiblk db      030h,14h,4fh,15h,06ah,13h,0c1h,11h,01h,10h,10h,30h
241                else
242 00D5 30144F156Asioiblk db      030h,14h,4fh,15h,06ah,13h,0c1h,11h,00h,10h,10h,30h
243                endif
244
245 000C =          sioilen equ    $-sioiblk
246
247
248                page

```

CP/M RMAC ASSEM 1.1 #006 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

249
250                ;      Network Initialization Routine
251
252                ntwrkinit:
253
254 00E1 CD0000          call   csdll          ; cold start the data link
255 00E4 CD0000          call   dllon         ; initialize the SIO drivers
256 00E7 3E01          mvi    a,rqstr$id    ; register the id with the data link
257 00E9 CD0000          call   regshrt
258 00EC AF            xra    a          ; return with no error
259 00ED C9            ret
260
261
262                ;      Return network status byte
263
264                ntwrksts:
265
266 00EE 3A1200          lda    network$status
267 00F1 47            mov    b,a
268 00F2 E6FC          ani    not (rcverr or senderr)
269 00F4 78            mov    a,b
270 00F5 C9            ret
271
272
273                ;      Return configuration table address
274
275                cnfgtbladr:
276
277 00F6 211200          lxi    h,configtbl
278 00F9 C9            ret
279
280                ;      Network error routine
281
282                ntwrkerror:
283
284 00FA 0E09          mvi    c,9
285 00FC 110000          lxi    d,network$error$msg
286 00FF CD0500          call   bdos
287
288
289 0102 C9            ret
290
291                ;      Network Warm Boot Routine
292
293                ntwrkwboot:                ; this entry is unused in this version
294
295 0103 C9            ret
296
297
298                ;      Send a Message on the Network
299                ;      Input:
300                ;      BC=pointer to message buffer
301                ;      Output:
302                ;      A = 0 if successful

```

CP/M RMAC ASSEM 1.1 #007 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

303                ;                1 if failure
304
305                sendmsg:
306
307                0104 C5                push    b
308                0105 60                mov     h,b
309                0106 69                mov     l,c
310
311                0107 364B                mvi    m,fmt$byte                ;set ulc$net format byte
312
313                0109 23                inx    h                ;reformat source to virtual circuit
314                010A 23                inx    h
315                010B 56                mov     d,m
316                010C 2B                dcx    h
317                010D 72                mov     m,d
318
319
320                010E 23                inx    h
321                010F 23                inx    h
322                0110 46                mov     b,m                ;save function
323
324                0111 23                inx    h
325                0112 5E                mov     e,m                ;get size
326                0113 70                mov     m,b                ;function=msg(0) in ULC format
327
328                0114 1600                mvi    d,0
329                0116 13                inx    d
330                0117 13                inx    d                ;normalize CP/NET to ULC sizes
331
332                0118 2B                dcx    h
333                0119 72                mov     m,d
334                011A 2B                dcx    h
335                011B 73                mov     m,e
336
337                011C C1                pop     b                ;restore buffer pointer
338
339                011D C34A01                jmp     dl$send                ;blast away
340
341
342                ;                Receive a Message on the Network
343                ;
344                ;                This routine calls the data-link routine to receive the message,
345                ;                then converts it into ULCnet format.
346                ;
347                ;                Input:
348                ;                BC = pointer to buffer to receive the message
349                ;                Output:
350                ;                A = 0 if successful
351                ;                1 if failure
352
353                receivemsg:
354
355                0120 C5                push    b                ;save buffer pointer
356

```

```

CP/M RMAC ASSEM 1.1    #008    REQUESTER NETWORK I/O SYSTEM FOR ULCNET

357                0121 CD3701                call    dl$receive                ;slurp the message
358
359                0124 E1                pop     h
360                0125 3601                mvi    m,1                ;FMT = 0 (requester to server)
361
362                0127 23                inx    h                ;DID already = virtual circuit #
363
364                0128 23                inx    h                ;get length
365                0129 5E                mov     e,m
366                012A 23                inx    h
367                012B 56                mov     d,m
368
369                012C 1B                dcx    d
370                012D 1B                dcx    d                ;normalize ULC to CP/NET format
371
372                012E 23                inx    h
373                012F 7E                mov     a,m                ;save FNC
374
375                0130 73                mov     m,e                ;format SIZ (<256)

```

```

376
377 0131 2B          dcx    h
378 0132 77          mov    m,a          ;format FNC
379
380 0133 2B          dcx    h
381 0134 AF          xra    a          ;set success
382 0135 77          mov    m,a          ;assume server always 0
383
384 0136 C9          ret      ;CP/NET message formatted form ULCnet
385
386
387
388          ; Data Link Interface Routines
389
390
391          ; DL$RECEIVE: Network Receive Function.
392          ;      Input:
393          ;          BC = Buffer address
394
395
396          dl$receive:
397
398 0137 50          mov    d,b          ; Buffer address in DE for data link
399 0138 59          mov    e,c
400
401          rretry:
402
403 0139 AF          xra    a          ; Packet mode
404 013A 010101      lxi    b,257       ; Buffer size
405 013D 210000      lxi    h,0         ; Infinite wait
406 0140 D5          push   d          ; Save buffer address for retry
407
408 0141 CD7801      call   psrecv     ; Initiate Receive and wait for completion
409
410 0144 D1          pop    d          ; Restore buffer address

```

CP/M RMAC ASSEM 1.1 #009 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

411 0145 B7          ora    a
412 0146 C8          rz          ; Return if no error
413
414 0147 C33901      jmp    rretry     ; Jump to try again if error
415
416
417          ; DL$SEND: Network Transmit Function
418          ;      Input:
419          ;          BC = Buffer address
420
421          dl$send:
422
423 014A 50          mov    d,b          ; Buffer address in DE for data link
424 014B 59          mov    e,c
425
426          tretry:
427
428 014C AF          xra    a          ; Packet mode, wait for Net Idle
429 014D D5          push   d          ; Save buffer address for retry
430
431 014E CD5701      call   psxmit     ; Initiate Transmit, wait for completion
432
433 0151 D1          pop    d          ; Restore buffer address
434 0152 B7          ora    a
435 0153 C8          rz          ; Return if no error
436
437 0154 C34C01      jmp    tretry     ; Jump to retry if error
438
439          ; PSXMIT: Transmit the packet pointed at by DE. If carry flag is set
440          ;          then don't wait for the Net to become idle.
441          ;
442          ;          Returns the completion code in A
443          ;          0 - Transmission ok and Data Link Ack Received
444          ;          (In the case of multicast, no Ack required)
445          ;          2 - Transmission OK but no Data Link Ack received.
446          ;
447          ;          4 - Other error.
448
449          psxmit:

```



```

450
451 0157 CD0000      call    transmit          ; This will transmit, set return code
452
453                twait:
454
455 015A CD0000      call    gettcode          ; A := GETTCODE - Xmit return code
456 015D 5F          mov     e,a
457 015E 1600        mvi    d,0
458 0160 216901      lxi    h,trtbl           ; dispatch on the return code
459 0163 19          dad     d
460 0164 5E          mov     e,m
461 0165 23          inx    h
462 0166 66          mov     h,m
463 0167 6B          mov     l,e
464 0168 E9          pchl

```

CP/M RMAC ASSEM 1.1 #010 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

465
466                trtbl:
467
468 0169 7701        dw     psxret            ; Good transmission
469 016B 7701        dw     psxret            ; No Data Link Ack
470 016D 7701        dw     psxret            ; Too many collisions
471 016F 7701        dw     psxret            ; Transmitter is disabled
472 0171 5A01        dw     twait             ; Transmitter is idle
473 0173 5A01        dw     twait             ; Transmitter is in progress
474 0175 5A01        dw     twait             ; Transmitter is waiting for ack
475
476                psxret:
477
478 0177 C9          ret
479
480                ; PSRECV: Receive a packet into buffer pointed at by DE. Length of
481                ; packet must be less than length of buffer in BC. HL is the receive
482                ; timeout count.
483                ;
484                ; Upon return clear the carry bit if a packet received and ACKed.
485                ; Set the carry flag if any error occurred.
486
487                psrecv:
488
489 0178 CD0000      call    receive          ; Receive. Return code will be set
490
491                rwait:
492
493 017B CD0000      call    getrcode         ; A := GETRCODE
494
495 017E 5F          mov     e,a
496 017F 1600        mvi    d,0
497 0181 218A01      lxi    h,rrtbl           ; dispatch on the return code
498 0184 19          dad     d
499 0185 5E          mov     e,m
500 0186 23          inx    h
501 0187 66          mov     h,m
502 0188 6B          mov     l,e
503 0189 E9          pchl
504
505                rrtbl:
506
507 018A 9601        dw     rgood             ; Good receive
508 018C 9801        dw     rbad             ; Bad receive
509 018E 9801        dw     rbad             ; Disabled
510
511                if not interrupts
512 0190 9801        dw     rbad             ; Still idle after timeout
513                else
514                dw     ridle             ; Idle
515                endif
516
517 0192 7B01        dw     rwait            ; Inprogress
518 0194 7B01        dw     rwait            ; In progress and for us.

```

CP/M RMAC ASSEM 1.1 #011 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

520             if      interrupts
521             ridle:
522
523             call    rtmockh          ; Check for timeout
524             jc      ridle1          ; Jump if timeout
525             call    wait1           ; Wait 1 ms
526             jmp     rwait           ; Continue to wait if no timeout
527
528             ridle1:
529
530             call    dsblrecv         ; Disable the receiver
531             stc
532             ret                ; Return with error
533             endif
534
535             rgood:
536
537             0196 A7             ana    a
538             0197 C9             ret
539
540             rbad:
541
542             0198 37             stc                ; Indicate error
543             0199 C9             ret
544             page

```

CP/M RMAC ASSEM 1.1 #012 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

545
546
547             ; NIOD routines
548
549
550
551             ; SETBAUD: Set the baud rate based on the baud rate code in A. Do special
552             ;          logic for self-clocked mode.
553             ;
554             ;          0 = 9600 baud
555             ;          1 = 19200 baud
556             ;          9 = 76800 baud self-clock
557             ;          11= 153600 baud self-clock
558             ;          13= 307200 baud self-clock
559             ;
560             ; If this station cannot handle the requested baud rate, then set
561             ; the carry flag.
562
563             setbaud:
564
565             019A E60F             ani    0fh          ; mask all but the baud bits
566             019C 21C400          lxi    h,curbaud   ; are we at the current baud rate?
567             019F BE              cmp    m
568             01A0 C8              rz                ; yes-->all done
569
570             01A1 47              mov    b,a          ; else-->get baud rate generator value
571             01A2 E607             ani    7
572             01A4 5F              mov    e,a
573             01A5 1600            mvi    d,0
574
575             01A7 21C500          lxi    h,btbl      ; point to vertical-to-horizontal decode
576             01AA 19              dad    d            ; table
577
578             if      slfclkd
579             01AB 78              mov    a,b
580             01AC E608             ani    selfmsk     ; is this a self-clocked value?
581             01AE C2D601          jnz   selfclkd
582             endif
583
584             01B1 3E03            mvi    a,baudsl    ; get legal baud rate mask
585             01B3 A6              ana    m
586             01B4 37              stc
587             01B5 C8              rz                ; return with error if its an illegal rate
588
589             if      slfclkd
590             01B6 3E05            mvi    a,5         ; else-->switch off possible self-clock mode
591             01B8 D306             out   siocmd
592             01BA 3E6A            mvi    a,dtroff    ; disable DTR in SIO register 5
593             01BC D306             out   siocmd

```

```

594
595 01BE 3E04          mvi    a,4          ; disable sync mode in register 4
596 01C0 D306          out     siocmd
597 01C2 3E4F          mvi    a,disslf
598 01C4 D306          out     siocmd

```

CP/M RMAC ASSEM 1.1 #013 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

599                      endif
600
601 01C6 21CD00        lxi    h,baudtbl    ; point to async baud rate table
602
603                      outbau:
604
605 01C9 19             dad     d            ; get async baud rate value
606 01CA 7E             mov     a,m
607 01CB D300          out     baudgen     ; load it into the baud rate generator
608                                     ; NOTE: This is not a CTC
609
610 01CD 21C400        lxi    h,curbaud
611 01D0 70             mov     m,b         ; set current baud byte
612
613 01D1 CDA702        call   wait         ; allow the system to reach equilibrium
614
615 01D4 A7             ana     a            ; return success
616 01D5 C9             ret
617
618                      if      slfclkd
619                      ; Throw SIO into self-clocked mode
620
621                      selfclkd:
622
623 01D6 3E2A          mvi    a,baudsh     ; Is this a legal rate?
624 01D8 A6             ana     m
625 01D9 37             stc
626 01DA C8             rz                ; return an error if not
627
628 01DB 3E04          mvi    a,4          ; enable sync mode in register 4
629 01DD D306          out     siocmd
630 01DF 3E0F          mvi    a,enaslf
631 01E1 D306          out     siocmd
632
633 01E3 3E05          mvi    a,5          ; enable DTR in register 5
634 01E5 D306          out     siocmd
635 01E7 3EEA          mvi    a,dtron
636 01E9 D306          out     siocmd
637
638 01EB 21CF00        lxi    h,scbaudt    ; point to baud rate table for self-clock mode
639 01EE C3C901        jmp     outbau      ; program the baud rate generator
640                      endif
641
642
643                      ; DSBLXMIT: Disable the transmitter if in self clocked mode
644
645                      dsblxmit:
646
647                      if      slfclkd
648 01F1 3AC400        lda     curbaud     ; are we in self-clocked mode?
649 01F4 E608          ani     selfmsk
650 01F6 C8             rz                ; no-->don't bother
651
652 01F7 3E05          mvi    a,5          ; disable SIO from transmitting by disabling

```

CP/M RMAC ASSEM 1.1 #014 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

653 01F9 D306          out     siocmd     ; DTR in register 5
654 01FB 3E6A          mvi    a,dtroff
655 01FD D306          out     siocmd
656
657 01FF 3E05          mvi    a,5          ; Enable receive by re-enabling DTR
658 0201 D306          out     siocmd
659 0203 3EEA          mvi    a,dtron
660 0205 D306          out     siocmd
661                      endif
662
663 0207 C9             ret

```

```

664
665
666           ; XMIT: Transmit the byte in A on network A.
667
668
669           xmit:
670
671                 if      not interrupts
672 0208 F5         push    psw
673
674           xmit1:
675
676 0209 DB06       in      siostat      ; don't overrun the transmitter if we're
677 020B E604       ani     xrdymask     ; interrupt-driven; wait for TxReady
678 020D CA0902    jz      xmit1
679
680 0210 F1         pop     psw
681               endif
682
683 0211 D304       out     sioxmit      ; blast that byte
684 0213 C9         ret
685
686
687           ; RECV: Receive a byte from Network A. Set the carry flag if there was
688           ; a receive error.
689           ;
690           ; For Z80-SIO receive errors are handled by the special receive
691           ; condition interrupts.
692
693           recv:
694
695                 if      not interrupts
696 0214 CD5D02     call   netidle
697 0217 DA2702    jc      rto          ; set error condition if the net went idle
698
699 021A DB06       in      siostat      ; else-->wait until a character is in the
700 021C E601       ani     rrdymask     ; buffer
701 021E CA1402    jz      recv
702
703 0221 CD2A02     call   chkstat      ; check for receive errors
704
705                 else
706 0221 CD2A02     ana     a            ; clear carry flag

```

```

CP/M RMAC ASSEM 1.1   #015   REQUESTER NETWORK I/O SYSTEM FOR ULCNET

707               endif
708
709 0224 DB04       in      siorecv      ; input the character
710 0226 C9         ret
711
712           rto:          ; set an error
713
714 0227 AF         xra     a
715 0228 37         stc
716 0229 C9         ret
717
718
719           ; CHKSTAT: Check error status bits of a receive error. If not error then
720           ; clear the carry flag and return. Otherwise figure out which
721           ; error occurred and increment its counter and set the carry flag.
722           ; Issue an error reset command to the UART.
723
724
725           chkstat:
726
727 022A 3E01       mvi     a,1          ; get error status from SIO read register 1
728 022C D306       out     siocmd
729 022E DB06       in      siostat
730
731 0230 E670       ani     errbits
732 0232 C8         rz          ; no error occurred-->all done
733
734                 if      netstats     ; gather statistics on the type of error
735 0233 47         mov     b,a
736 0234 E610       ani     pmsk
737 0236 CA3F02    jz      np            ; not a parity error

```

```

738
739 0239 210000      lxi  h,parcntr      ; else-->
740 023C CD0000      call incntr         ; increment parity error counter
741
742                np:
743
744 023F 78           mov  a,b
745 0240 E605         ani  obit
746 0242 CA4B02      jz   no             ; not an overrun
747
748 0245 210000      lxi  h,ovrcntr     ; else-->
749 0248 CD0000      call incntr         ; increment overrun counter
750
751                no:
752
753 024B 78           mov  a,b
754 024C E606         ani  fbit
755 024E CA5702      jz   nf             ; not a framing error
756
757 0251 210000      lxi  h,frmcntr     ; else-->
758 0254 CD0000      call incntr         ; increment framing error counter
759
760                nf:

```

CP/M RMAC ASSEM 1.1 #016 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

761                endif
762
763 0257 3E30         mvi  a,errst       ; reset error condition
764 0259 D306         out  siocmd
765 025B 37          stc
766 025C C9          ret
767
768
769
770                ; NETIDLE: See if network A is idle. If idle then set the carry flag.
771
772                netidle:
773
774 025D 3E10         mvi  a,10h         ; reset interrupts
775 025F D306         out  siocmd
776 0261 D306         out  siocmd         ; do it twice to reject glitches on DCD
777
778 0263 DB06         in   siostat       ; is there a data-carrier detect?
779 0265 E608         ani  carmsk
780 0267 C8          rz
781
782 0268 AF          xra  a
783 0269 CD9A01      call setbaud       ; net is idle-->reset to hailing rate (9600)
784 026C 37          stc
785 026D C9          ret
786
787
788                if      interrupts
789
790                ; ENBLRECV: Enable the channel A receiver interrupts.
791
792                enblrcv:
793
794                mvi  a,1             ; enable interrupts on all characters
795                out  siocmd
796                mvi  a,011h         ; NOTE: This mask would have to be 015h on
797                out  siocmd         ; channel B
798                ret
799
800                ; DSBLRECV: Disable the channel A receiver interrupts.
801
802                dsblrcv:
803
804                mvi  a,1             ; Disable interrupts on received characters
805                out  siocmd         ; (Keep status interrupts enabled)
806                out  siocmd         ; NOTE: Channel B mask is 05h
807                ret
808
809                endif
810
811

```

```

812 ; PGMUART: Program the Network UART channel
813
814 pgmuart:

```

```

CP/M RMAC ASSEM 1.1 #017 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

```

815
816 if interrupts
817 ; The 820 already has the SIO vector address
818 ; programmed from channel B. Other
819 ; implementations will have to provide linkage
820 ; to the vector area in the main XIOS, and
821 ; load the vector offset into SIO write
822 ; register 2
823
824 lxi h,niisr ; load status interrupt service routine vector
825 shld siov5
826 lxi h,dlsir ; load transmit ISR vector
827 shld siov6
828 lxi h,reisr ; load receiv ISR vector
829 shld siov7
830 endif
831
832 026E 21D500 lxi h,sioiblk ; point to SIO initialization block
833 0271 060C mvi b,sioilen ; length of block
834 0273 F3 di
835
836 pgm1:
837
838 0274 7E mov a,m ; output the block to the SIO
839 0275 D306 out siocmd
840 0277 23 inx h
841 0278 05 dcr b
842 0279 C27402 jnz pgm1
843
844 027C FB ei
845 027D AF xra a ; set up hailing baud rate = 9600
846 027E CD9A01 call setbaud
847 0281 C9 ret
848
849
850 ; INITUART: Initialize the uart for network A by issuing a reset command
851 ; and clearing out the receive buffer.
852
853 inituart:
854
855 0282 3E03 mvi a,3 ; disable the receiver through register 3
856 0284 D306 out siocmd
857 0286 3EC0 mvi a,disrcv
858 0288 D306 out siocmd
859
860 028A DB06 in siostat ; is there a garbage byte?
861 028C E601 ani rrdymask
862 028E CA9602 jz initu ; no-->continue initialization
863
864 0291 DB04 in siorecv ; else-->eat the character
865 0293 C38202 jmp inituart ; try again
866
867 initu:
868

```

```

CP/M RMAC ASSEM 1.1 #018 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

```

869 0296 3E30 mvi a,errst ; reset error conditions
870 0298 D306 out siocmd
871
872 029A 3E03 mvi a,3 ; re-enable the receiver
873 029C D306 out siocmd
874 029E 3EC1 mvi a,enarcv
875 02A0 D306 out siocmd
876
877 02A2 C9 ret
878
879 ; INITRECV: Initialize a receive operation
880
881 initrecv:

```

```

882
883 02A3 CD8202          call    inituart
884
885                    if    interrupts
886                    call  enblrecv      ; enable receiver interrupts
887                    endif
888
889 02A6 C9              ret
890
891
892                    ; WAIT - Wait 100 micro seconds
893
894                    wait:
895
896 02A7 3E16            mvi    a,timeval
897
898                    w:
899
900 02A9 3D              dcr    a          ; 04
901 02AA A7              ana    a          ; 04
902 02AB C2A902         jnz    w          ; 12
903                    ; ---
904 02AE C9              ret          ; 30 T-States total
905
906
907                    ; RESTUART: Reinitialize the UART to the way it was in the
908                    ;          original BIOS after completing the network operations
909
910
911                    restuart:
912 02AF C9              ret          ; UART not used except by network
913
914
915                    ; CSNIOD: Do any cold start initialization which is necessary.
916                    ;          Must at least return the value of BAUDS
917                    ;          If the network uses the printer port then set theh carry flag
918                    ;          otherwise clear it.
919
920                    csniod:
921
922 02B0 01032A         lxi    b,bauds    ; return the legal baud rates

```

CP/M RMAC ASSEM 1.1 #019 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

923 02B3 B7            ora    a          ; not using a printer port
924 02B4 C9            ret
925
926
927 02B5                end

```

CP/M RMAC ASSEM 1.1 #020 REQUESTER NETWORK I/O SYSTEM FOR ULCNET

```

ACTIVE          0010 129#
BAUDGEN         0000 71# 607
BAUDS           2A03 69# 922
BAUDSH         002A 66# 69 623
BAUDSL         0003 65# 69 584
BAUDTBL        00CD 225# 601
BDOS           0005 136# 287
BTBL           00C5 222# 575
BUFLEN         0106 116#
CARBIT         0003 81#
CARMSK         0008 82# 779
CHKSTAT        022A 42 703 725#
CNFGTBLADR     00F6 145 275#
CONFIGTBL      0012 170# 277
CSDLL          0000 53 254
CSNIOD         02B0 43 920#
CURBAUD        00C4 216# 566 610 648
DID            0001 111# 112
DISRCV         00C0 96# 857
DISSLF         004F 98# 597
DLLBAU         00C2 45 210#
DLLON          0000 53 255
DLRECEIVE      0137 357 396#
DLSEND         014A 339 421#

```


SIOIBLK	00D5	240#	242#	245	832		
SIOILEN	000C	245#	833				
SIORECV	0004	75#	709	864			
SIOSTAT	0006	73#	676	699	729	778	860
SIOV4	FF08	102#					
SIOV5	FF0A	103#	825				
SIOV6	FF0C	104#	827				
SIOV7	FF0E	105#	829				
SIOXMIT	0004	74#	683				
SIZ	0004	114#	115				
CP/M RMAC ASSEM 1.1	#022	REQUESTER	NETWORK	I/O	SYSTEM	FOR	ULCNET
SLFCLKD	FFFF	36#	578	589	618	647	
TERRCNT	0000	54					
TIMEVAL	0016	212#	896				
TRANSMIT	0000	51	451				
TRETRY	014C	426#	437				
TRTBL	0169	458	466#				
TRUE	FFFF	32#	35	36			
TWAIT	015A	453#	472	473	474		
ULCFMT	0000	120#	121				
ULCFNC	0004	124#	125				
ULCLENHI	0003	123#	124				
ULCLENLO	0002	122#	123				
ULCMSG	0005	125#					
ULCVCIRC	0001	121#	122				
W	02A9	898#	902				
WAIT	02A7	43	613	894#			
XMIT	0208	40	669#				
XMIT1	0209	674#	678				
XRDYBIT	0002	77#					
XRDYMSK	0004	78#	677				

Listing F-1: Requester Network I/O System for ULCnet

F.3 Creating the ULCnet Server

The server communications software is contained in the modules XIOSNET.ASM and ULCIF.ASM. XIOSNET.ASM contains modifications to MP/M II's XIOS. ULCIF.ASM is the equivalent of the NETWRKIF transport processes.

ULCIF.ASM uses only two processes, one for input and one for output. To use ULCIF.ASM with the module SERVER.RSP, you must patch SERVER.RSP to write all message responses to a single output queue named NtwrkQ00. This patch is detailed in *CP/NET V1.2 Application Note #2* dated 11-11-82.

The communications interface is interrupt driven, servicing each character as it is received by the network port. ULCIF.ASM requests the network resource through a set of dummy console I/O calls to the XIOS. A call to CONST initializes the network. Calls to CONIN and CONOUT receive and send messages on the network. The communications interface checks network status through a set of poll calls.

The ULCIF input transport process is dispatched at MP/M II cold start. This process makes all necessary queues, creates the ULCIF output process, initializes the network, and writes the configuration table address into the system data page. ULCIF then goes into a loop where it perpetually performs the following actions:

1. Allocates a buffer for an incoming message. If no buffer is available, ULCIF repeats the allocation process until a buffer becomes available.
2. Receives a message by placing the dummy console number in register D, a pointer to the message buffer just allocated in register pair BC, and calling CONIN in the XIOS.
3. Converts the ULCnet format message into CP/NET format. To do this, ULCnet assumes that the virtual circuit number and the requester source ID are identical.

4. Matches the requester ID with a requester control block. If no server is allocated to this requester and the message is a login, ULCIF allocates a server if one is available. Otherwise, ULCIF writes an extended error message to the output queue, NtwrkQO0.
5. Using the requester control block, ULCIF writes the address of the message buffer to the appropriate input queue, NtwrkQI.
6. Repeats.

The output process performs the following actions:

1. Reads the output queue, NtwrkQI0.
2. If the message is a LOGOFF function, frees the appropriate requester control block entry.
3. Converts the message response from CP/NET format into ULCnet format. To do this, ULCnet uses the requester destination ID as the virtual circuit number.
4. Places the dummy console number into register D, the message buffer address into register pair BC, and calls CONOUT in the XIOS.
5. Repeats.

The ULCnet modules DLIF and NIOD are contained in the module XIOSNET.ASM. This module must be incorporated into the server's XIOS. XIOSNET.ASM handles four XIOS jump vector entries, CONST, CONIN, CONOUT, and POLLDEVICE. The jump vector in the XIOS must be modified to point to these routines. XIOSNET contains a linkage to the real XIOS routines for these functions, in this case renamed NCONST, NCONIN, NCONOUT, and POLDEV. The XIOS's interrupt vector might also have to be modified to support the SIO interrupt service routines in IPBMAIN.

When the console I/O routines are entered, they immediately check to see if the dummy console number has been supplied.

Note: you must define a console number that does not conflict with real consoles. Make the dummy console number at least larger than the number of requesters to be supported, since each server process pretends to attach to a unique console ID. If a dummy console number has not been supplied, these routines jump into the real console routines. If the dummy number has been supplied, the routines take the following steps.

CONST:

1. performs network initialization.
2. registers the expected Requester ID's as virtual circuit numbers by repeatedly calling REGSHRT.
3. returns to the ULCIF. This routine is called only once.

CONIN:

1. Calls RECEIVE, using the buffer pointer passed from ULCIF
2. Executes the MP/M II poll function, specifying a poll device routine that repeatedly performs the GETRCODE function until its status shows that a message has been received properly.
3. Returns to the ULCIF.

CONOUT:

1. Calls TRANSMIT, using the buffer pointer passed from ULCIF.
2. Executes the poll function, specifying a poll device routine that repeatedly performs the GETTCODE function until the message has been sent and received by the destination without error.
3. Returns to the ULCIF.

The POLLDEVICE routine behaves almost like the console I/O routines. POLLDEVICE checks for specific poll device numbers to perform network status functions. If these numbers are not detected, control passes to the real POLDEV routine. If network status functions are detected, POLLDEVICE performs the appropriate status check. If the check is successful, a hexadecimal 0FF is returned in register A. If not successful, a 0 is

returned.

The MP/M II dispatcher calls POLLDEVICE when it is entered. If the status returned is 0, MP/M II maintains the poll device number on a list and continues to call POLLDEVICE every time it is entered. When the returned status is FF, the dispatcher removes the device number from its list and returns control to the code that originally performed the poll function call, in this case either CONIN or CONOUT. In this manner, the communications interface operates completely transparently, requiring very little CPU resource.

The XIOSNET is designed to be interrupt driven. The IPBMAIN.REL module performs the actual data-link. This module is identical to the IPBMAIN.REL used in the SNIOS. An interrupt-driven protocol is strongly recommended. If you use the polled version, PBMAIN, calls to TRANSMIT and RECEIVE do not return until the requested operation has been performed. This means communications software uses up enormous amounts of CPU time, suspending only when a clock tick interrupts them and forces the dispatcher to be entered. This results in poor server performance.

The interrupt-driven IPBMAIN module sets up the requested operation only when TRANSMIT and RECEIVE are called. The actual protocol is driven by the arrival or departure of each character of the message. This interrupt-driven protocol consumes considerably less CPU time.

To modify the modules ULCIF and XIOSNET for your own server:

1. Patch the module SERVER.RSP to write all of its outputs to a single queue, as described in an application note.
2. Only three parameters must be modified in the ULCIF if four or fewer requesters are to be supported.

Set NMB\$RQSTRS to the number of requesters supported.

Set NMB\$BUFS to the number of requesters, plus one. This extra buffer permits the transmission of LOGIN error messages to the output process, even when all SERVER processes are busy. Having fewer buffers limits the burden on the server at any one time.

Set CONSOLE\$NUM to the dummy console number. The sample listing uses the arbitrarily large number hex 20. This number should be sufficient.

3. If more than four requesters are supported, you must provide extra QCBs, requester control blocks, stack space, and Process Descriptor areas.
4. Modify the XIOS jump vector to jump into the XIOSNET routines CONST, CONIN, CONOUT, and POLLDEVICE. You might have to make additional PUBLIC and EXTRN declarations.
5. Include linkage access to the XIOS interrupt vector. If the XIOS has no interrupt vector, create one.
6. Make sure the false console number specified by the ULCIF module agrees with the one used by XIOSNET.
7. Make sure the device numbers CONIN and CONOUT use in their poll calls do not conflict with other device numbers used by the XIOS.
8. Customize the NIOD section of XIOSNET the same way you customized this section in ULCNIOS.ASM.
9. Create a resident or banked XIOS by linking the regular XIOS module with the network interface:

```
A>LINK RESXIOS=<regular XIOS modules>,XIOSNET,IPBMAIN[05]
```

If you are creating a banked system, all of XIOSNET must reside in common memory.

10. Build the ULCIF.RSP module:

```
A>RMAC ULCIF
```



```

40 ;*****
41 ;
42 ;       This software was developed jointly by
43 ;
44 ;       Digital Research, Inc.
45 ;       P.O. Box 579
46 ;       Pacific Grove, CA 93950
47 ;       and
48 ;       Keybrook Business Systems, Inc.
49 ;       2035 National Avenue
50 ;       Hayward, CA 94545
51 ;
52 ;
53 bdosadr:
54 0000 0000          dw    $-$          ; RSP XDOS entry point

```

CP/M RMAC ASSEM 1.1 #002 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

55 ;
56 ; User-Configurable Parameters (These should be the only changes needed)
57 ;
58 0002 = nmb$rqstrs equ 2 ; Number of requesters supported at one time
59 0003 = nmb$bufs equ 3 ; Number of message buffers
60 0020 = console$num equ 20h ; Pseudo-console number
61 004B = fmt$byte equ 4bh ; Format byte: short format with acknowledge,
62 ; 153.6K baud self-clocked
63 ;
64 ; Message Buffer Offsets
65 ;
66 0000 = fmt equ 0 ; format
67 0001 = did equ fmt+1 ; destination ID
68 0002 = sid equ did+1 ; source ID
69 0003 = fnc equ sid+1 ; server function number
70 0004 = siz equ fnc+1 ; size of message (normalized to 0)
71 0005 = msg equ siz+1 ; message
72 0106 = buf$len equ msg+257 ; length of total message buffer
73 ;
74 ; ULCnet Packet Offsets
75 ;
76 0000 = ulc$fmt equ 0 ; packet format
77 0001 = ulc$v$circ equ ulc$fmt+1 ; virtual circuit number
78 0002 = ulc$len$lo equ ulc$v$circ+1 ; low order of length
79 0003 = ulc$len$hi equ ulc$len$lo+1 ; high order of length
80 0004 = ulc$fnc equ ulc$len$hi+1 ; start of message: function code
81 0005 = ulc$msg equ ulc$fnc+1 ; CP/NET message
82 ;
83 ; Requester Control Block Offsets
84 ;
85 0000 = rqstr$id equ 0 ; requester ID for this server
86 0001 = uqcb equ rqstr$id+1 ; uqcb to queue to this server
87 0005 = buf$ptr equ uqcb+4 ; queue message <-> msg buffer ptr
88 0007 = rcb$len equ buf$ptr+2 ; length of requester control block
89 ;
90 ;
91 ; NETWRKIF Process Descriptors and Stack Space
92 ;
93 networkin: ; Receiver Process
94 ;
95 0002 0000 dw 0 ; link
96 0004 00 db 0 ; status
97 0005 42 db 66 ; priority
98 0006 6400 dw netstkin+46 ; stack pointer
99 0008 4E4554752 db 'NETWRKIN' ; name
100 0010 00 db 0 ; console
101 0011 FF db 0ffh ; memseg
102 0012 ds 2 ; b
103 0014 ds 2 ; thread
104 0016 ds 2 ; buff
105 0018 ds 1 ; user code & disk slct
106 0019 ds 2 ; dcnt
107 001B ds 1 ; searchl
108 001C ds 2 ; searcha

```

CP/M RMAC ASSEM 1.1 #003 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

109 001E ds 2 ; active drives

```

```

110 0020 0000          dw      0          ; HL'
111 0022 0000          dw      0          ; DE'
112 0024 0000          dw      0          ; BC'
113 0026 0000          dw      0          ; AF'
114 0028 0000          dw      0          ; IY
115 002A 0000          dw      0          ; IX
116 002C 0000          dw      0          ; HL
117 002E 0000          dw      0          ; DE
118 0030 0000          dw      0          ; BC
119 0032 0000          dw      0          ; AF, A = ntwkif console dev #
120 0034              ds      2          ; scratch
121
122                    netstkin:
123 0036 C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
124 003E C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
125 0046 C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
126 004E C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
127 0056 C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
128 005E C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h
129 0064 B405           dw      setup
130
131                    networkout:          ; Transmitter Process
132
133 0066 0000          dw      0          ; link
134 0068 00           db      0          ; status
135 0069 42           db      66         ; priority
136 006A C800          dw      netstkou+46 ; stack pointer
137 006C 4E45545752   db      'NETWRKOU' ; name
138 0074 00           db      0          ; console
139 0075 FF           db      0ffh       ; memseg
140 0076              ds      2          ; b
141 0078              ds      2          ; thread
142 007A              ds      2          ; buff
143 007C              ds      1          ; user code & disk slct
144 007D              ds      2          ; dcnt
145 007F              ds      1          ; searchl
146 0080              ds      2          ; searcha
147 0082              ds      2          ; active drives
148 0084 0000          dw      0          ; HL'
149 0086 0000          dw      0          ; DE'
150 0088 0000          dw      0          ; BC'
151 008A 0000          dw      0          ; AF'
152 008C 0000          dw      0          ; IY
153 008E 0000          dw      0          ; IX
154 0090 0000          dw      0          ; HL
155 0092 0000          dw      0          ; DE
156 0094 0000          dw      0          ; BC
157 0096 0000          dw      0          ; AF, A = ntwkif console dev #
158 0098              ds      2          ; scratch
159
160                    netstkou:
161 009A C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
162 00A2 C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h

```

CP/M RMAC ASSEM 1.1 #004 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

163 00AA C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
164 00B2 C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
165 00BA C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h,0c7c7h
166 00C2 C7C7C7C7       dw      0c7c7h,0c7c7h,0c7c7h
167 00C8 8606           dw      output
168
169
170                    ; Input queue control blocks
171
172                    qcb$in$0:
173 00CA              ds      2          ; link
174 00CC 4E7477726B     db      'NtwrkQI0' ; name
175 00D4 0200          dw      2          ; msglen
176 00D6 0100          dw      1          ; nmbmsgs
177 00D8              ds      2          ; dqph
178 00DA              ds      2          ; nqph
179 00DC              ds      2          ; msgin
180 00DE              ds      2          ; msgout
181 00E0              ds      2          ; msgcnt
182 00E2              ds      2          ; buffer
183

```

```

184             if      nmb$rqstrs ge 2
185             qcb$in$1:
186 00E4             ds      2           ; link
187 00E6 4E7477726B db      'NtwrkQI1'      ; name
188 00EE 0200       dw      2           ; msglen
189 00F0 0100       dw      1           ; nmbmsgs
190 00F2             ds      2           ; dqph
191 00F4             ds      2           ; nqph
192 00F6             ds      2           ; msgin
193 00F8             ds      2           ; msgout
194 00FA             ds      2           ; msgcnt
195 00FC             ds      2           ; buffer
196             endif
197
198             if      nmb$rqstrs ge 3
199             qcb$in$2:
200             ds      2           ; link
201             db      'NtwrkQI2'      ; name
202             dw      2           ; msglen
203             dw      1           ; nmbmsgs
204             ds      2           ; dqph
205             ds      2           ; nqph
206             ds      2           ; msgin
207             ds      2           ; msgout
208             ds      2           ; msgcnt
209             ds      2           ; buffer
210             endif
211
212             if      nmb$rqstrs ge 4
213             qcb$in$3:
214             ds      2           ; link
215             db      'NtwrkQI3'      ; name
216             dw      2           ; msglen

```

CP/M RMAC ASSEM 1.1 #005 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

217             dw      1           ; nmbmsgs
218             ds      2           ; dqph
219             ds      2           ; nqph
220             ds      2           ; msgin
221             ds      2           ; msgout
222             ds      2           ; msgcnt
223             ds      2           ; buffer
224             endif
225
226             ; Output queue control blocks
227
228             qcb$out$0:
229 00FE             ds      2           ; link
230 0100 4E7477726B db      'NtwrkQ00'      ; name
231 0108 0200       dw      2           ; msglen
232 010A 0300       dw      nmb$bufs    ; nmbmsgs
233 010C             ds      2           ; dqph
234 010E             ds      2           ; nqph
235 0110             ds      2           ; msgin
236 0112             ds      2           ; msgout
237 0114             ds      2           ; msgcnt
238 0116             ds      2*nmb$bufs+1 ; buffer
239
240             ; Requester Management Table
241
242             rqstr$table:
243
244             ;requester 0 control block
245
246 011D FF         db      0ffh        ; requester ID (marked not in use)
247 011E CA00       dw      qcb$in$0    ; UQCB: QCB pointer
248 0120 2201       dw      $+2        ; pointer to queue message
249 0122 0000       dw      $-$        ; pointer to msg buffer (loaded on receive)
250
251             if      nmb$rqstrs ge 2
252             ;requester 1 control block
253
254 0124 FF         db      0ffh        ; requester ID (marked not in use)
255 0125 E400       dw      qcb$in$1    ; UQCB: QCB pointer
256 0127 2901       dw      $+2        ; pointer to queue message
257 0129 0000       dw      $-$        ; pointer to msg buffer (loaded on receive)

```

```

258             endif
259
260             if      nmb$rqstrs ge 3
261 ;requester 2 control block
262
263             db      0ffh           ; requester ID (marked not in use)
264             dw      qcb$in$2      ; UQCB:  QCB pointer
265             dw      $+2           ;           pointer to queue message
266             dw      $-$           ;           pointer to msg buffer (loaded on receive)
267             endif
268
269             if      nmb$rqstrs ge 4
270 ;requester 3 control block

```

CP/M RMAC ASSEM 1.1 #006 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

271
272             db      0ffh           ; requester ID (marked not in use)
273             dw      qcb$in$3      ; UQCB:  QCB pointer
274             dw      $+2           ;           pointer to queue message
275             dw      $-$           ;           pointer to msg buffer (loaded on receive)
276             endif
277
278 ; Output user queue control block
279
280 uqcb$out$0:
281 012B FE00      dw      qcb$out$0      ; pointer
282 012D 2F01      dw      out$buffer$ptr ; pointer to queue message
283
284 out$buffer$ptr:
285 012F          ds      2              ; a queue read will return the message
286                                     ; buffer pointer in this location
287
288 ; UQCB for flagging errors from receive process to send process
289
290 uqcb$in$out$0:
291 0131 FE00      dw      qcb$out$0      ; pointer
292 0133 3501      dw      in$out$buffer$ptr
293                                     ; pointer to queue message
294
295 in$out$buffer$ptr:
296 0135          ds      2              ; this pointer used by input process to
297                                     ; to output "server not logged in" errors
298
299 ; Server Configuration Table
300
301 configtbl:
302 0137 00      db      0              ; Server status byte
303 0138 00      db      0              ; Server processor ID
304 0139 02      db      nmb$rqstrs    ; Max number of requesters supported at once
305 013A 00      db      0              ; Number of currently logged in requesters
306 013B 0000    dw      0000h         ; 16 bit vector of logged in requesters
307 013D          ds      16           ; Logged In Requester processor ID's
308 014D 5041535357 db      'PASSWORD' ; login password
309
310 ; Stacks for server processes. A pointer to the associated process
311 ; descriptor area must reside on the top of each stack. The stack for
312 ; SERVROPR is internal to SERVER.RSP, and is consequently omitted from the
313 ; NETWRKIF module.
314
315 0096 =      srvr$stk$len equ      96h ; server process stack size
316
317             if      nmb$rqstrs ge 2
318 0155      srvr$stk$1: ds      srvr$stk$len-2
319 01E9 EB01      dw      srvr$1$pd
320             endif
321
322             if      nmb$rqstrs ge 3
323      srvr$stk$2: ds      srvr$stk$len-2
324             dw      srvr$2$pd

```

CP/M RMAC ASSEM 1.1 #007 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

325             endif
326
327             if      nmb$rqstrs ge 4

```



```

328          srvr$stk$3:    ds      srvr$stk$len-2
329                                dw      srvr$3$pd
330                                endif
331
332          ; Memory allocation for server process descriptor copydown
333          ; All server process descriptor allocation must be contiguous
334
335          01EB          srvr$1$pd:    if      nmb$rqstrs ge 2
336                                ds      52
337                                endif
338
339          01EB          srvr$2$pd:    if      nmb$rqstrs ge 3
340                                ds      52
341                                endif
342
343          01EB          srvr$3$pd:    if      nmb$rqstrs ge 4
344                                ds      52
345                                endif
346
347
348          ; Buffer Control Block: 0 indicates buffer is free for receiving a message
349          ; 0ffh indicates that the buffer is in use
350
351          buf$cb:        rept      nmb$bufs
352                                db      0
353                                endm
354          021F+00        DB      0
355          0220+00        DB      0
356          0221+00        DB      0
357
358          ; Message Buffer Storage Area
359
360          msg$buffers:   rept      nmb$bufs
361                                ds      buf$len
362                                endm
363          0222+          DS      BUF$LEN
364          0328+          DS      BUF$LEN
365          042E+          DS      BUF$LEN
366
367          ; save area for XIOS routine addresses
368
369          conin$jmp:
370          0534 C3        db      jmp
371          0535 0000      conin:  dw      $-$
372
373          conout$jmp:
374          0537 C3        db      jmp
375          0538 0000      conout: dw      $-$
376
377          constat$jmp:
378          053A C3        db      jmp

```

CP/M RMAC ASSEM 1.1 #008 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

379          constat:
380          053B 0000      dw      $-$
381
382
383
384
385          ; NETWRKIF Utility Routines
386
387          ; Operating system linkage routine
388
389          monx:
390
391          053D 2A0000      lhld    bdos$adr
392          0540 E9        pchl
393
394
395          ; Double word subtract: DE = HL - DE
396
397          dw$sub:
398          0541 7D        mov     a,l
399          0542 93        sub     e
400          0543 5F        mov     e,a
401          0544 7C        mov     a,h

```

```

402 0545 9A          sbb    d
403 0546 57          mov    d,a
404 0547 C9          ret
405
406                ; Routine to scan requester control blocks for a match with the received
407                ; source ID.
408                ;
409                ; Input:  A = Source ID to Match
410                ;
411                ; Output:
412                ;   success:  HL = pointer to requester control block
413                ;           A <> 0FFh
414                ;   no match, but a free control block found:
415                ;           HL = pointer to RCB
416                ;           A = 0FFh
417                ;           CY = 0
418                ;   no match and no available RCB's:
419                ;           A = 0FFh
420                ;           CY = 1
421
422                scan$table:
423
424 0548 211D01        lxi    h,rqstr$table          ;point to the start of the RCB table
425 054B 0602        mvi    b,nmb$rqstrs
426 054D 110700        lxi    d,rcb$len             ;size of RCB's for scanning the table
427
428                sc$t1:
429
430 0550 BE          cmp    m                      ;RCB ID = SID?
431 0551 C8          rz                          ;yes--> a match--> return
432

```

CP/M RMAC ASSEM 1.1 #009 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

433 0552 19          dad    d                      ;else-->check next entry
434 0553 05          dcr    b
435 0554 C25005        jnz    sc$t1
436
437 0557 211D01        lxi    h,rqstr$table          ;no match-->look for a free entry
438 055A 0602        mvi    b,nmb$rqstrs
439
440                sc$t2:
441
442 055C 7E          mov    a,m
443 055D 3C          inr    a
444 055E CA6A05        jz     sc$t3                  ;an unoccupied entry has been found
445
446 0561 19          dad    d                      ;else-->keep looking
447 0562 05          dcr    b
448 0563 C25C05        jnz    sc$t2
449
450 0566 3EFF        mvi    a,0ffh                ;outa luck-->set the big error
451 0568 37          stc
452 0569 C9          ret
453
454                sc$t3:
455                ;no match, but found a free entry
456 056A 3D          dcr    a                      ;A=0FFh
457 056B B7          ora    a                      ;CY=0
458 056C C9          ret
459
460
461                ; This routine free up a requester control block for somebody else who
462                ; might want to Log In.
463                ;
464                ;   Input:  A = source ID that just logged off
465
466                free$rqstr$tbl:
467
468 056D 211D01        lxi    h,rqstr$table
469 0570 110700        lxi    d,rcb$len
470
471                fr$t1:
472
473 0573 BE          cmp    m                      ;RCB ID <> SID-->keep scanning
474 0574 C27A05        jnz    fr$t2
475

```

```

476 0577 36FF          mvi    m,0ffh          ;else-->mark it as unoccupied
477 0579 C9           ret                      ; and bug out
478
479                 fr$t2:
480
481 057A 19           dad     d
482 057B C37305       jmp     fr$t1          ;keep going--it's in there somewhere
483
484
485
486                 ; Routine to send a message on the network

```

CP/M RMAC ASSEM 1.1 #010 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

487                 ; Input: HL = pointer to message buffer
488
489                 send$msg:
490
491 057E E5           push   h
492 057F 364B       mvi    m,fmt$byte     ;set ulc$net format byte
493
494 0581 23           inx    h               ;virtual circuit = requester ID
495
496 0582 23           inx    h
497 0583 23           inx    h
498
499 0584 46           mov    b,m            ;save function number
500
501 0585 23           inx    h               ;get SIZ
502 0586 5E           mov    e,m
503
504 0587 1600        mvi    d,0            ;normalize CP/NET to ULCnet length
505 0589 13           inx    d
506 058A 13           inx    d
507
508 058B 70           mov    m,b            ;put FNC in first message byte
509
510 058C 2B           dcx    h               ;store length
511 058D 72           mov    m,d
512 058E 2B           dcx    h
513 058F 73           mov    m,e
514
515 0590 C1           pop    b               ;restore buffer pointer
516 0591 1620        mvi    d,console$num  ;set up fake console number for xios
517 0593 C33705     jmp    conout$jmp     ;blast that packet
518
519
520                 ; Routine to receive a message on the network
521                 ; Input: DE = pointer to buffer
522
523                 rcv$message:
524
525 0596 42           mov    b,d
526 0597 4B           mov    c,e
527 0598 C5           push   b               ;save buffer pointer
528 0599 1620        mvi    d,console$num  ;receive the message
529 059B CD3405     call   conin$jmp
530
531 059E E1           pop    h
532 059F 3600        mvi    m,0            ;FMT = 0 (requester to server)
533
534 05A1 23           inx    h
535 05A2 46           mov    b,m            ;save rqstr ID = virtual circuit
536
537 05A3 3A3801     lda    configtbl+1
538 05A6 77           mov    m,a            ;DID = server ID
539
540 05A7 23           inx    h

```

CP/M RMAC ASSEM 1.1 #011 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

541 05A8 5E           mov    e,m            ;get low order length
542
543 05A9 70           mov    m,b            ;SID = requester ID
544
545 05AA 23           inx    h

```

```

546 05AB 56          mov    d,m          ;get hi order length
547
548 05AC 1B          dcx    d
549 05AD 1B          dcx    d          ;normalize ULCnet to CP/NET SIZ
550
551 05AE 23          inx    h
552 05AF 46          mov    b,m          ;get FNC
553
554 05B0 73          mov    m,e          ;store SIZ
555
556 05B1 2B          dcx    h
557 05B2 70          mov    m,b          ;store FNC
558
559 05B3 C9          ret
560
561
562
563
564
565                ; Network I/F Receiver Process
566
567
568                setup:
569
570                ;initialize NETWRKIF
571                ;loop counter for making n+1 queues
572                ;make queue function code
573                ;make all input and output queue(s)
574                ;makeq:
575
576 05BB C5          push   b
577 05BC D5          push   d
578 05BD CD3D05     call  monx
579
580 05C0 E1          pop    h
581 05C1 111A00     lxi    d,26
582 05C4 19          dad    d
583 05C5 EB          xchg
584
585 05C6 C1          pop    b
586 05C7 05          dcr    b
587 05C8 C2BB05     jnz   makeq
588
589 05CB 0E9A       mvi    c,154
590 05CD CD3D05     call  monx
591
592 05D0 110900     lxi    d,9          ;write configuration table address
593 05D3 19          dad    d          ; into system data page, allowing
594 05D4 113701     lxi    d,configtbl ; server initialization to proceed

```

CP/M RMAC ASSEM 1.1 #012 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

595 05D7 F3          di
596 05D8 73          mov    m,e
597 05D9 23          inx    h
598 05DA 72          mov    m,d
599 05DB FB          ei
600
601 05DC 2B          dcx    h          ;point to XIOS jump table page
602 05DD 2B          dcx    h
603 05DE 2B          dcx    h
604 05DF 66          mov    h,m
605 05E0 2E00       mvi    l,0
606
607 05E2 110600     lxi    d,6
608 05E5 19          dad    d          ;point to constat
609 05E6 223B05     shld  constat
610
611 05E9 23          inx    h
612 05EA 23          inx    h
613 05EB 23          inx    h          ;point to conin
614 05EC 223505     shld  conin
615
616 05EF 23          inx    h
617 05F0 23          inx    h
618 05F1 23          inx    h
619 05F2 223805     shld  conout      ;point to conout

```

```

620
621 05F5 1620      mvi    d,console$num
622 05F7 CD3A05    call   constat$jmp      ;use constat to initialize ulcnet
623
624 05FA 116600    lxi    d,networkout    ;create network I/F output process
625 05FD 0E90      mvi    c,144
626 05FF CD3D05    call   monx
627
628          input:      ;input process loop
629
630          ; Find a free buffer
631
632 0602 211F02    lxi    h,buf$cb        ;point to buffer control block
633 0605 112202    lxi    d,msg$buffers   ;point to base of buffer area
634 0608 0603      mvi    b,nmb$buffers   ;get total number of buffers
635
636          input2:
637
638 060A 7E        mov    a,m
639 060B 3C        inr    a
640 060C C22306    jnz    input3          ;we found a free buffer-->use it
641
642 060F E5        push   h               ;point to next buffer
643 0610 210601    lxi    h,buf$len
644 0613 19        dad    d
645 0614 EB        xchg
646
647 0615 E1        pop    h               ;point to next buffer control field
648 0616 23        inx    h

```

CP/M RMAC ASSEM 1.1 #013 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

649
650 0617 05        dcr    b               ;have we scanned all the buffers?
651 0618 C20A06    jnz    input2
652
653 061B 0E8E      mvi    c,142          ;uh oh, we're all clogged up
654 061D CD3D05    call   monx           ;dispatch and go sleepy bye for a bit
655 0620 C30206    jmp    input          ;try again
656
657          input3:
658
659 0623 36FF      mvi    m,0ffh        ;found a buffer-->mark it used
660
661 0625 D5        push   d
662
663          ; Receive the message
664
665 0626 CD9605    call   rcv$message
666
667 0629 E1        pop    h
668 062A E5        push   h
669
670 062B 23        inx    h               ;check requester table to see
671 062C 23        inx    h               ; whether the source requester
672 062D 7E        mov    a,m            ; is logged-in
673 062E CD4805    call   scan$table
674
675 0631 3C        inr    a
676 0632 CA4A06    jz     input4         ;not logged-in-->go check for login
677
678          input6:
679
680 0635 110500    lxi    d,buf$ptr      ;else-->update message buffer pointer
681 0638 19        dad    d
682
683 0639 D1        pop    d
684 063A 73        mov    m,e
685 063B 23        inx    h
686 063C 72        mov    m,d
687
688 063D 11FBFF    lxi    d,uqcb-buf$ptr-1 ;point to the uqcb for this requester
689 0640 19        dad    d
690 0641 EB        xchg
691
692 0642 0E8B      mvi    c,139          ;write the message to the queue
693 0644 CD3D05    call   monx

```

```

694
695 0647 C30206      jmp    input          ;round and round we go
696
697                input4:                ;else-->requester not logged-in
698
699 064A D1           pop    d
700 064B 13          inx   d
701 064C 13          inx   d
702 064D 13          inx   d

```

CP/M RMAC ASSEM 1.1 #014 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

703 064E DA6006      jc    input5          ;bomb the message if there's no
704                                     ; table entries left
705
706 0651 1A          ldax  d
707 0652 FE40        cpi   64              ;is it a login?
708 0654 C26006      jnz   input5
709
710 0657 1B          dcx   d               ;yes-->mark the control block with
711 0658 1A          ldax  d               ; the source ID
712 0659 77          mov   m,a
713
714 065A 1B          dcx   d               ;go do the queue write
715 065B 1B          dcx   d
716 065C D5          push  d
717 065D C33506      jmp   input6
718
719                input5:                ;flag a "not logged in" extended error
720
721 0660 EB          xchg  h
722 0661 23          inx   h
723 0662 3601        mvi   m,1             ;set SIZ=1
724 0664 23          inx   h
725 0665 36FF        mvi   m,0ffh         ;set return code to error
726 0667 23          inx   h
727 0668 360C        mvi   m,0ch          ;flag extended error 12
728
729 066A 11FAFF      lxi   d,fmt-msg-1
730 066D 19          dad   d               ;point back at message start
731 066E 3601        mvi   m,1             ;format = 1
732
733 0670 23          inx   h               ;swap DID and SID
734 0671 7E          mov   a,m
735 0672 23          inx   h
736 0673 46          mov   b,m
737 0674 77          mov   m,a
738 0675 2B          dcx   h
739 0676 70          mov   m,b
740 0677 2B          dcx   h
741
742 0678 223501      shld  in$out$buffer$ptr ;write buffer pointer to queue msg buf
743
744 067B 113101      lxi   d,uqcb$in$out$0 ;write to the queue
745 067E 0E8B        mvi   c,139
746 0680 CD3D05      call  monx
747 0683 C30206      jmp   input          ;try again
748
749
750
751                ; Network I/F transmitter process
752
753                output:
754
755 0686 112B01      lxi   d,uqcb$out$0    ;read the output queue-->go sleepy
756 0689 0E89        mvi   c,137           ; bye until some server process

```

CP/M RMAC ASSEM 1.1 #015 NETWRKIF FOR SYSTEMS RUNNING ULCNET

```

757 068B CD3D05      call  monx            ; sends a response
758
759 068E 2A2F01      lhld  out$buffer$ptr
760 0691 EB          xchg  d
761 0692 D5          push  d               ;save message pointer
762
763 0693 210300      lxi   h,fnc           ;get message function code

```

```

764 0696 19          dad    d
765 0697 7E          mov    a,m
766 0698 2B          dcx    h
767
768 0699 FE41         cpi    65          ;is it a logoff?
769 069B C2A206      jnz    output2
770
771 069E 7E          mov    a,m          ;load SID
772 069F CC6D05      cz     free$rqstr$tbl ;yes-->free up the server process
773
774                output2:
775
776 06A2 E1          pop    h
777 06A3 E5          push   h
778 06A4 CD7E05      call   send$msg     ;send the message
779
780 06A7 E1          pop    h          ;retrieve message pointer
781
782 06A8 112202      lxi    d,msg$buffers ;DE = pointer - message buffer base
783 06AB CD4105      call   dw$sub
784
785 06AE 011F02      lxi    b,buf$cb    ;BC = DE/buf$len + buf$cb
786
787                output3:
788
789 06B1 7B          mov    a,e
790 06B2 B2          ora    d
791 06B3 CAC106      jz     output4
792
793 06B6 EB          xchg
794 06B7 110601      lxi    d,buf$len
795 06BA CD4105      call   dw$sub
796 06BD 0C          inr    c
797 06BE C3B106      jmp    output3
798
799                output4:
800
801 06C1 AF          xra    a
802 06C2 02          stax  b          ;free the buffer for re-use
803
804 06C3 C38606      jmp    output     ;transmission without end, amen
805
806 06C6             end

```

CP/M RMAC ASSEM 1.1 #016 NETWRKIF FOR SYSTEMS RUNNING ULCNET

BDOSADR	0000	53#	391						
BUFCB	021F	351#	632	785					
BUFLEN	0106	72#	361	363	364	365	643	794	
BUFPTR	0005	87#	88	680	688				
CONFIGTBL	0137	301#	537	594					
CONIN	0535	371#	614						
CONINJMP	0534	369#	529						
CONOUT	0538	375#	619						
CONOUTJMP	0537	373#	517						
CONSOLENUM	0020	60#	516	528	621				
CONSTAT	053B	379#	609						
CONSTATJMP	053A	377#	622						
DID	0001	67#	68						
DWSUB	0541	397#	783	795					
FMT	0000	66#	67	729					
FMTBYTE	004B	61#	492						
FNC	0003	69#	70	763					
FREERQSTRTBL	056D	466#	772						
FRT1	0573	471#	482						
FRT2	057A	474	479#						
INOUTBUFFERPTR	0135	292	295#	742					
INPUT	0602	628#	655	695	747				
INPUT2	060A	636#	651						
INPUT3	0623	640	657#						
INPUT4	064A	676	697#						
INPUT5	0660	703	708	719#					
INPUT6	0635	678#	717						
MAKEQ	05BB	574#	587						
MONX	053D	389#	578	590	626	654	693	746	757
MSG	0005	71#	72	729					
MSGBUFFERS	0222	360#	633	782					

NETSTKIN	0036	98	122#									
NETSTKOU	009A	136	160#									
NETWORKIN	0002	93#										
NETWORKOUT	0066	131#	624									
NMBBIFS	0003	59#	232	238	351	360	634					
NMBRQSTRS	0002	58#	184	198	212	251	260	269	304	317	322	
		327	335	339	343	425	438	570				
OUTBUFFERPTR	012F	282	284#	759								
OUTPUT	0686	167	753#	804								
OUTPUT2	06A2	769	774#									
OUTPUT3	06B1	787#	797									
OUTPUT4	06C1	791	799#									
QCBIN0	00CA	172#	247	572								
QCBIN1	00E4	185#	255									
QCBOUT0	00FE	228#	281	291								
RCBLEN	0007	88#	426	469								
RCVMESAGE	0596	523#	665									
RQSTRID	0000	85#	86									
RQSTRTABLE	011D	242#	424	437	468							
SCANTABLE	0548	422#	673									
SCT1	0550	428#	435									
SCT2	055C	440#	448									
SCT3	056A	444	454#									

CP/M RMAC ASSEM	1.1	#017	NETWRKIF FOR SYSTEMS RUNNING ULCNET									
SENDMSG	057E	489#	778									
SETUP	05B4	129	568#									
SID	0002	68#	69									
SIZ	0004	70#	71									
SRVR1PD	01EB	319	336#									
SRVRSTK1	0155	318#										
SRVRSTKLEN	0096	315#	318	323	328							
ULCFMT	0000	76#	77									
ULCFNC	0004	80#	81									
ULCLENHI	0003	79#	80									
ULCLENLO	0002	78#	79									
ULCMSG	0005	81#										
ULCVCIRC	0001	77#	78									
UQCB	0001	86#	87	688								
UQCBINOUT0	0131	290#	744									
UQCBOUT0	012B	280#	755									

Listing F-2: NETWRKIF for Systems Running ULCnet

CP/M RMAC ASSEM	1.1	#001	ULCNET DATA LINK LAYER MP/M XIOS MODULE									
1			title 'ULCNET Data Link Layer MP/M XIOS Module'									
2			page 54									
3												
4			;*****									
5			;* This module must be linked into the server's XIOS. It is designed to *									
6			;* run under MP/M for the Xerox 820, but should be easily customized. It *									
7			;* contains the ULCnet interface modules DLIF and NIOD. The DLIF is an *									
8			;* interface between the transport software contained in ULCIF.RSP and the *									
9			;* data-link software contained in IPBMAIN.REL. The NIOD contains the actual*									
10			;* hardware drivers required to run ULCnet. The module IPBMAIN.REL must also*									
11			;* be linked into the XIOS. *									
12			;*****									
13												
14			; This software is the result of a joint effort between									
15			;									
16			; Digital Research, Inc.									
17			; P.O. Box 579									
18			; Pacific Grove, CA 93950									
19			; and									
20			; Keybrook Business Systems, Inc.									
21			; 2035 National Avenue									
22			; Hayward, CA 94545									
23												
24			; Conditional assembly control									
25												
26	FFFF =	true	equ	0ffffh								
27	0000 =	false	equ	not true								
28												


```

29 FFFF =      interrupts      equ      true      ; false=pollled, true=interrupt-driven
30 FFFF =      netstats       equ      true      ; switch to gather network statistics
31 FFFF =      slfclkd        equ      true      ; supports self-clocked operation
32
33           ; Linkage information
34
35           public nconst,nconin,nconout ; XIOS console jump table entries
36           public polldevice           ; XIOS polling routine
37           public setbaud,xmit,recv,initu ; NIOD routines called by IPBMAIN
38           public inituart,pgmuart
39           public chkstat,netidle,initrecv
40           public wait,restuart,csniod
41           public dsblxmit
42           public dllbau,netadr
43
44           if      interrupts
45           public enblrecv,dsblrecv
46           endif
47
48           extrn  transmit,receive      ; IPBMAIN routines and objects
49           extrn  gettcode,getrcode
50           extrn  csdll,dllon,regshrt
51           extrn  terrcnt,parcntr,ovrcntr
52           extrn  frmcntr,incntr
53           extrn  xdos,const,conin,conout ; linkage back to the rest of XIOS
54           extrn  poldev

```

CP/M RMAC ASSEM 1.1 #002 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

55
56           if      interrupts
57           extrn  rtmochk           ; IPBMAIN interrupt routines
58           extrn  dlisr,reisr,niisr
59           endif
60
61           ; Hardware definitions for the Z80-SIO channel A - For the Xerox 820.
62
63           0003 =      baudsl  equ      03h      ; Usable baud rates: 9600, 19.2K asynch.,
64           002A =      baudsh  equ      2ah      ; 76.8K, 153.6K, 307.2K self-clocked
65
66           ; baud rate capability mask
67           2A03 =      bauds   equ      (baudsh*100h)+baudsl
68
69           0000 =      baudgen  equ      0        ; External baud rate generator register
70           0006 =      siocmd  equ      6        ; Command/Mode register
71           0006 =      siostat  equ      6        ; Status register
72           0004 =      sioxmit  equ      4        ; Transmit register
73           0004 =      siorecv  equ      4        ; Receive register
74
75           0002 =      xrdybit  equ      2        ; Transmit buffer empty status bit
76           0004 =      xrdymask equ      4        ; transmit buffer empty status mask
77           0000 =      rrdybit  equ      0        ; Receive buffer full status bit
78           0001 =      rrdymask equ      1        ; receive buffer full status mask
79           0003 =      carbit   equ      3        ; Net Idle detect bit position
80           0008 =      carmask  equ      8        ; Net Idle detect mask
81           0030 =      errrst   equ      030h     ; Error flag reset
82           0070 =      errbits  equ      070h     ; Error bit position mask
83           0004 =      pbit     equ      4        ; Parity error bit position
84           0010 =      pmsk     equ      10h      ; parity error mask
85           0005 =      obit     equ      5        ; Overrun error bit position
86           0020 =      omsk     equ      20h      ; overrun error mask
87           0006 =      fbit     equ      6        ; Framing error bit position
88           0040 =      fsmk     equ      40h      ; framing error mask
89           0003 =      selfbit  equ      3        ; Self clock bit position
90           0008 =      selfmsk  equ      8        ; self clock bit mask
91           00EA =      dtron    equ      0eah     ; Turn on DTR
92           006A =      dtroff   equ      06ah     ; Turn off DTR
93           00C1 =      enarcv   equ      0c1h     ; Enable receive-clock
94           00C0 =      disrcv   equ      0c0h     ; Disable receive clock
95           000F =      enaslf   equ      00fh     ; Enable Self-clock mode
96           004F =      disslf   equ      04fh     ; Disable Self-clock mode
97
98           ; SIO Mode 2 interrupts vector table
99
100          101 FF08 =      siov4   equ      0ff08h   ; SIO port A xmit buffer empty
101          102 FF0A =      siov5   equ      0ff0ah   ; SIO port A external status change

```

```

103 FF0C =      siov6 equ    0ff0ch      ; SIO port A receive
104 FF0E =      siov7 equ    0ff0eh      ; SIO port A special receive condition
105
106 0020 =      netcon equ    20h        ; fake console number called by ULCIF for
107                                           ; network operations
108

```

CP/M RMAC ASSEM 1.1 #003 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

109           ; polling equates
110
111 0020 =      ulctx equ    20h          ; transmission poll number
112 0021 =      ulcrx equ    21h          ; receive poll number
113
           page

```

CP/M RMAC ASSEM 1.1 #004 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

114
115
116
117           ; ULCnet Data Definitions
118
119 0000        netadr: ds      3          ;ULCnet network address
120 0003        dllbau: ds      2          ;baud rate mask
121
122 0016 =      timeval equ    22          ; WAIT routine time constant
123                                           ; 12 for 2.5 megahertz Z80
124                                           ; 22 for 4.0 megahertz Z80
125
126           dev$table:          ;polling device table
127
128 0005 9800          dw      twait          ;receive poll wait
129 0007 D300          dw      rwait          ;transmit poll wait
130 0002 =      num$devices equ    ($-dev$table)/2
131
132 0009        tcode: ds      1          ; Transmit Return code
133 000A        rcode: ds      1          ; Receive Return code
134
135 000B FF        curbaud db    0ffh      ; Current baud rate
136
137
138 000C 0102040810btbl: db    1,2,4,8,16,32,64,128 ; table to convert baud number codes
139                                           ; into a bit mask
140
141           baudtbl:          ; async baud rate table
142
143 0014 0E        db      0eh            ; 9600 Baud
144 0015 0F        db      0fh            ; 19200
145
146           scbaudt:          ; self-clock baud rate table
147
148 0016 00        db      0              ; 62500 Baud - Not implemented
149 0017 0D        db      0dh            ; 76800 Baud
150 0018 00        db      0              ; 125000 Baud - Not implemented
151 0019 0E        db      0eh            ; 153600 Baud
152 001A 00        db      0              ; 250000 Baud - Not implemented
153 001B 0F        db      0fh            ; 307200 Baud
154
155           if      interrupts
156 001C 30144F156Asioiblk db    030h,14h,4fh,15h,06ah,13h,0c1h,11h,01h,10h,10h,30h
157           else
158           sioiblk db    030h,14h,4fh,15h,06ah,13h,0c1h,11h,00h,10h,10h,30h
159           endif
160
161 000C =      sioilen equ    $-sioiblk
162
163           page

```

CP/M RMAC ASSEM 1.1 #005 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

164
165
166
167           ; ULCnet data-link interface code

```

```

168
169
170         ; POLLDEVICE: Device polling routine.
171         ;       Input:
172         ;           C = device number to poll
173         ;       Output:
174         ;           A = 0   if not ready
175         ;           0ffh if ready
176
177     polldevice:
178
179     0028 79             mov     a,c           ; if not a network poll, go to the real
180     0029 D620          sui     ulctx        ; routine
181     002B DA0000        jc      poldev
182
183     002E FE02          cpi     num$devices ; check for poll number in bounds
184     0030 DA3600        jc      devok
185
186     0033 3E00          mvi     a,0           ; out-of-bounds-->don't do anything
187     0035 C9           ret
188
189     devok:
190
191     0036 6F           mov     l,a
192     0037 2600          mvi     h,0
193     0039 29           dad     h             ; multiply index by 2
194
195     003A 110500        lxi     d,dev$table  ; index into the poll routine table
196     003D 19           dad     d
197
198     003E 5E           mov     e,m
199     003F 23           inx     h
200     0040 56           mov     d,m           ; get the routine address
201
202     0041 EB           xchg
203     0042 E9           pchl
204
205
206
207     ;
208     ; NCONST: Console status entry point. If register D = fake network
209     ; console ID, do network initialization. Otherwise, go back to
210     ; the real console routines.
211
212     nconst:
213
214     0043 3E20          mvi     a,netcon     ; Check if network call
215     0045 BA           cmp     d
216     0046 C20000        jnz     const        ; Jump to normal CONST if not network
217

```

CP/M RMAC ASSEM 1.1

#006

ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

218     0049 CD0000        call    csdll         ; Cold start the data link
219     004C CD0000        call    dllon        ; Initialize the SIO Drivers
220     004F AF           xra     a             ; Initialize all the short addresses
221
222     nxtadd:
223
224     0050 3C           inr     a
225     0051 FE05          cpi     5             ; Check for last address
226     0053 C8           rz
227     0054 F5           push   psw
228     0055 CD0000        call    regshrt
229     0058 F1           pop    psw
230     0059 C35000        jmp     nxtadd        ; Jump to process next address
231
232
233     ; NCONIN: Console In entry point. If register D = the fake network ID
234     ; then receive a network message, using polled status checks of
235     ; an interrupt-driven data-link. Otherwise, go back to the real
236     ; CONIN routine.
237
238     nconin:
239
240     005C 3E20          mvi     a,netcon     ; Check for network call
241     005E BA           cmp     d

```

```

242 005F C20000      jnz  conin      ; Jump to normal CONIN if not network
243
244 0062 50          mov  d,b        ; Setup for PSRECEIVE
245 0063 59          mov  e,c
246
247                rretry:
248
249 0064 AF          xra  a          ; Packet mode
250 0065 010101     lxi  b,257     ; Buffer size
251 0068 210000     lxi  h,0       ; Infinite wait
252 006B D5          push d         ; Save buffer address for retry
253 006C CDC100     call psrecv    ;
254 006F D1          pop  d         ; Restore buffer address
255 0070 B7          ora  a
256 0071 C8          rz           ; Return if no error
257
258 0072 C36400     jmp  rretry    ; Jump to try again if error
259
260
261                ; NCONOUT: Console out entry point. If D = fake console ID, send a network
262                ; message. Otherwise, just head for the real CONOUT routine.
263
264
265                nconout:
266
267 0075 3E20       mvi  a,netcon  ; Check for network call
268 0077 BA         cmp  d
269 0078 C20000     jnz  conout    ; Jump to normal CONOUT if not network
270
271 007B 50         mov  d,b        ; Setup for PSXMIT

```

CP/M RMAC ASSEM 1.1 #007 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

272 007C 59          mov  e,c
273
274                tretry:
275
276 007D AF          xra  a          ; Packet mode, wait for Net Idle
277 007E D5          push d         ; Save buffer address for retry
278 007F CD8800     call psxmit    ;
279 0082 D1          pop  d         ; Restore buffer address
280 0083 B7          ora  a
281 0084 C8          rz           ; Return if no error
282
283 0085 C37D00     jmp  tretry    ; Jump to retry if error
284
285
286                ; PSXMIT: Transmit the packet pointed at by DE. If carry flag is set
287                ; then don't wait for the Net to become idle.
288                ;
289                ; Returns the completion code in A:
290                ;
291                ; 0 - Transmission ok and Data Link Ack Received
292                ; (In the case of multicast, no Ack required)
293                ; 2 - Transmission OK but no Data Link Ack received.
294                ;
295                ; 4 - Other error.
296
297                psxmit:
298
299 0088 CD0000     call transmit  ; TRETCODE := TRANSMIT(TBUFPtr,)
300
301 008B 0E83       mvi  c,83h     ; Poll the transmitter for completion
302 008D 1E20       mvi  e,ulctx
303 008F CD0000     call xdos
304
305 0092 3A0900     lda  tcode     ; Fetch return code
306 0095 C3CE00     jmp  exitdl
307
308                ; TWAIT: Transmission completion poll routine.
309                ;
310                ; Output:
311                ; A = 0 if not complete
312                ; 0ffh if complete
313
314                twait:
315

```

```

316 0098 CD0000      call   gettcode          ; A := GETTCODE - Xmit return code
317
318 009B 5F          mov    e,a              ; get return code processing vectore
319 009C 1600        mvi   d,0
320 009E 21A700      lxi   h,trtbl
321 00A1 19          dad   d
322
323 00A2 5E          mov    e,m              ; dispatch on return code
324 00A3 23          inx   h
325 00A4 66          mov    h,m

```

CP/M RMAC ASSEM 1.1 #008 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

326 00A5 6B          mov    l,e
327 00A6 E9          pchl
328
329                ; Return code dispatch table
330
331 00A7 B700      trtbl: dw   psxret          ; Good transmission
332 00A9 B700      dw   psxret          ; No Data Link Ack
333 00AB B700      dw   psxret          ; Too many collisions
334 00AD B700      dw   psxret          ; Transmitter is disabled
335 00AF B500      dw   tsleep          ; Transmitter is idle
336 00B1 B500      dw   tsleep          ; Transmitter is in progress
337 00B3 B500      dw   tsleep          ; Transmitter is waiting for ack
338
339                tsleep:
340
341 00B5 AF          xra   a              ; Code for continue to sleep
342 00B6 C9          ret
343
344                psxret:
345                ; Enter here if something happened
346 00B7 D2BB00      jnc   twakeup        ; Jump if no transmit error
347 00BA 2F          cma
348                ; Else-->Indicate error
349
350                twakeup:
351 00BB 320900      sta   tcode          ; Store return code
352 00BE 3EFF        mvi   a,0ffh        ; Signal poll successful
353 00C0 C9          ret
354
355
356
357                ; PSRECV: Receive a packet into buffer pointed at by DE. Length of
358                ; packet must be less than length of buffer in BC. HL is the receive
359                ; timeout count.
360                ;
361                ; Upon return clear the carry bit if a packet received and ACKed.
362                ; Set the carry flag if any error occurred.
363                ;
364
365                psrecv:
366
367 00C1 CD0000      call   receive        ; := RECEIVE(HL,DE,BC)
368
369 00C4 0E83        mvi   c,83h          ; Poll until receive complete
370 00C6 1E21        mvi   e,ulcrx
371 00C8 CD0000      call   xdos
372
373 00CB 3A0A00      lda   rcode          ; Fetch return code
374
375                ; Common exit routine for returning to the pseudo-console handler
376
377                exitdl:
378
379

```

CP/M RMAC ASSEM 1.1 #009 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

380 00CE B7          ora   a              ; Assume no error
381 00CF F0          rp
382                ; Return if no error
383 00D0 2F          cma
384 00D1 37          stc
385 00D2 C9          ret
                ; Indicate error

```

```

386
387 ; RWAIT: Poll routine to detect receive status.
388 ;
389 ; Output:
390 ; A = 0 if receive not complete
391 ; 0ffh if receive complete
392
393 rwait:
394
395 00D3 CD0000 call getrcode ; A := GETRCODE
396
397 00D6 5F mov e,a ; form dispatch vector
398 00D7 1600 mvi d,0
399 00D9 21E200 lxi h,rrtbl
400 00DC 19 dad d
401
402 00DD 5E mov e,m ; dispatch on receive completion code
403 00DE 23 inx h
404 00DF 66 mov h,m
405 00E0 6B mov l,e
406 00E1 E9 pchl
407
408 ; Receive completion code dispatch table
409
410 00E2 F000 rrrtbl: dw rgood ; Good receive
411 00E4 F600 dw rbad ; Bad receive
412 00E6 F600 dw rbad ; Disabled
413
414 if not interrupts
415 dw rbad ; Still idle after timeout
416 else
417 00E8 FA00 dw ridle ; Idle
418 endif
419
420 00EA EE00 dw rsleep ; Inprogress
421 00EC EE00 dw rsleep ; In progress and for us.
422
423 rsleep:
424
425 00EE AF xra a ; Code for continue to sleep
426 00EF C9 ret
427
428 rgood:
429 rwakeup:
430
431 00F0 320A00 sta rcode ; Store return code
432 00F3 3EFF mvi a,0ffh ; Wake up code
433 00F5 C9 ret

```

CP/M RMAC ASSEM 1.1 #010 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

434
435 rbad:
436
437 00F6 2F cma ; Code for error
438 00F7 C3F000 jmp rwakeup ; Jump to wake up receive process
439
440 if interrupts
441
442 ridle:
443
444 00FA CD0000 call rtmochk ; Check for timeout
445 00FD DAF600 jc rbad ; if timeout, signal error
446 0100 C3EE00 jmp rsleep ; Continue to wait if no timeout
447
448 0103 C9 ret
449
450 endif
451 page

```

CP/M RMAC ASSEM 1.1 #011 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

452
453
454 ; NIOD routines
455

```

```

456
457
458 ; SETBAUD: Set the baud rate based on the baud rate code in A. Do special
459 ; logic for self-clocked mode.
460 ;
461 ; 0 = 9600 baud
462 ; 1 = 19200 baud
463 ; 9 = 76800 baud self-clock
464 ; 11= 153600 baud self-clock
465 ; 13= 307200 baud self-clock
466 ;
467 ; If this station cannot handle the requested baud rate, then set
468 ; the carry flag.
469
470 setbaud:
471
472 0104 E60F      ani    0fh          ; mask all but the baud bits
473 0106 210B00   lxi    h,curbaud   ; are we at the current baud rate?
474 0109 BE       cmp    m
475 010A C8       rz          ; yes-->all done
476
477 010B 47       mov    b,a          ; else-->get baud rate generator value
478 010C E607     ani    7
479 010E 5F       mov    e,a
480 010F 1600     mvi    d,0
481
482 0111 210C00   lxi    h,btbl      ; point to vertical-to-horizontal decode
483 0114 19       dad    d            ; table
484
485             if    slfclkd
486 0115 78       mov    a,b
487 0116 E608     ani    selfmsk     ; is this a self-clocked value?
488 0118 C24001   jnz   selfclkd
489             endif
490
491 011B 3E03     mvi    a,baudsl   ; get legal baud rate mask
492 011D A6       ana    m
493 011E 37       stc
494 011F C8       rz          ; return with error if its an illegal rate
495
496             if    slfclkd
497 0120 3E05     mvi    a,5         ; else-->switch off possible self-clock mode
498 0122 D306     out   siocmd
499 0124 3E6A     mvi    a,dtroff   ; disable DTR in SIO register 5
500 0126 D306     out   siocmd
501
502 0128 3E04     mvi    a,4         ; disable sync mode in register 4
503 012A D306     out   siocmd
504 012C 3E4F     mvi    a,disslf
505 012E D306     out   siocmd

```

CP/M RMAC ASSEM 1.1 #012 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

506             endif
507
508 0130 211400   lxi    h,baudtbl   ; point to async baud rate table
509
510 outbau:
511
512 0133 19       dad    d            ; get async baud rate value
513 0134 7E       mov    a,m
514 0135 D300     out   baudgen     ; load it into the baud rate generator
515             ; NOTE: This is not a CTC
516
517 0137 210B00   lxi    h,curbaud
518 013A 70       mov    m,b         ; set current baud byte
519
520 013B CD1E02   call   wait        ; allow the system to reach equilibrium
521
522 013E A7       ana    a            ; return success
523 013F C9       ret
524
525             if    slfclkd
526 ; Throw SIO into self-clocked mode
527
528 selfclkd:
529

```

```

530 0140 3E2A      mvi    a,baudsh      ; Is this a legal rate?
531 0142 A6        ana    m
532 0143 37        stc
533 0144 C8        rz                    ; return an error if not
534
535 0145 3E04      mvi    a,4           ; enable sync mode in register 4
536 0147 D306      out    siocmd
537 0149 3E0F      mvi    a,enaslf
538 014B D306      out    siocmd
539
540 014D 3E05      mvi    a,5           ; enable DTR in register 5
541 014F D306      out    siocmd
542 0151 3EEA      mvi    a,dtron
543 0153 D306      out    siocmd
544
545 0155 211600     lxi    h,scbaudt     ; point to baud rate table for self-clock mode
546 0158 C33301     jmp    outbau        ; program the baud rate generator
547
548
549
550                ; DSBLXMIT: Disable the transmitter if in self clocked mode
551
552                dsblxmit:
553
554                if    slfclkd
555 015B 3A0B00     lda    curbaud      ; are we in self-clocked mode?
556 015E E608      ani    selfmsk
557 0160 C8        rz                    ; no-->don't bother
558
559 0161 3E05      mvi    a,5           ; disable SIO from transmitting by disabling

```

CP/M RMAC ASSEM 1.1

#013

ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

560 0163 D306      out    siocmd      ; DTR in register 5
561 0165 3E6A      mvi    a,dtroff
562 0167 D306      out    siocmd
563
564 0169 3E05      mvi    a,5           ; Enable receive by re-enabling DTR
565 016B D306      out    siocmd
566 016D 3EEA      mvi    a,dtron
567 016F D306      out    siocmd
568
569
570 0171 C9        ret
571
572
573                ; XMIT: Transmit the byte in A on network A.
574
575
576                xmit:
577
578                if    not interrupts
579                push  psw
580
581                xmit1:
582
583                in    siostat      ; don't overrun the transmitter if we're
584                ani    xrdymask    ; interrupt-driven; wait for TxReady
585                jz    xmit1
586
587                pop   psw
588                endif
589
590 0172 D304      out    sioxmit     ; blast that byte
591 0174 C9        ret
592
593
594                ; RECV: Receive a byte from Network A. Set the carry flag if there was
595                ; a receive error.
596                ;
597                ; For Z80-SIO receive errors are handled by the special receive
598                ; condition interrupts.
599
600                recv:
601
602                if    not interrupts
603                call  netidle

```



```

604          jc      rto          ; set error condition if the net went idle
605
606          in      siostat       ; else-->wait until a character is in the
607          ani     rrdymask      ;   buffer
608          jz      recv
609
610          call    chkstat       ; check for receive errors
611
612          else
613 0175 A7    ana     a            ; clear carry flag

```

CP/M RMAC ASSEM 1.1 #014 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

614          endif
615
616 0176 DB04  in      siorecv     ; input the character
617 0178 C9   ret
618
619          rto:                ; set an error
620
621 0179 AF   xra     a
622 017A 37   stc
623 017B C9   ret
624
625
626          ; CHKSTAT: Check error status bits of a receive error. If not error then
627          ; clear the carry flag and return. Otherwise figure out which
628          ; error occurred and increment its counter and set the carry flag.
629          ; Issue an error reset command to the UART.
630
631
632          chkstat:
633
634 017C 3E01  mvi     a,1          ; get error status from SIO read register 1
635 017E D306  out     siocmd
636 0180 DB06  in      siostat
637
638 0182 E670  ani     errbits
639 0184 C8   rz              ; no error occurred-->all done
640
641          if      netstats     ; gather statistics on the type of error
642 0185 47   mov     b,a
643 0186 E610  ani     pmsk
644 0188 CA9101 jz      np              ; not a parity error
645
646 018B 210000 lxi    h,parcntr      ; else-->
647 018E CD0000 call   incntr         ; increment parity error counter
648
649          np:
650
651 0191 78   mov     a,b
652 0192 E605  ani     obit
653 0194 CA9D01 jz      no              ; not an overrun
654
655 0197 210000 lxi    h,ovrcntr     ; else-->
656 019A CD0000 call   incntr         ; increment overrun counter
657
658          no:
659
660 019D 78   mov     a,b
661 019E E606  ani     fbit
662 01A0 CAA901 jz      nf              ; not a framing error
663
664 01A3 210000 lxi    h,frmcntr     ; else-->
665 01A6 CD0000 call   incntr         ; increment framing error counter
666
667          nf:

```

CP/M RMAC ASSEM 1.1 #015 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

668          endif
669
670 01A9 3E30  mvi     a,errst     ; reset error condition
671 01AB D306  out     siocmd
672 01AD 37   stc              ; signal an error
673 01AE C9   ret

```

```

674
675
676
677      ; NETIDLE: See if network A is idle. If idle then set the carry flag.
678
679      netidle:
680
681      01AF 3E10          mvi    a,10h          ; reset interrupts
682      01B1 D306          out    siocmd
683      01B3 D306          out    siocmd          ; do it twice to reject glitches on DCD
684
685      01B5 DB06          in     siostat         ; is there a data-carrier detect?
686      01B7 E608          ani    carmsk
687      01B9 C8           rz           ; yes-->net is in use-->carry flag cleared
688
689      01BA AF           xra    a
690      01BB CD0401        call   setbaud         ; net is idle-->reset to hailing rate (9600)
691      01BE 37           stc
692      01BF C9           ret          ; set net idle to true
693
694
695      if      interrupts
696
697      ; ENBLRECV: Enable the channel A receiver interrupts.
698
699      enblrcv:
700
701      01C0 3E01          mvi    a,1           ; enable interrupts on all characters
702      01C2 D306          out    siocmd
703      01C4 3E11          mvi    a,011h        ; NOTE: This mask would have to be 015h on
704      01C6 D306          out    siocmd        ; channel B
705      01C8 C9           ret
706
707      ; DSBLRECV: Disable the channel A receiver interrupts.
708
709      dsblrcv:
710
711      01C9 3E01          mvi    a,1           ; Disable interrupts on received characters
712      01CB D306          out    siocmd        ; (Keep status interrupts enabled)
713      01CD D306          out    siocmd        ; NOTE: Channel B mask is 05h
714      01CF C9           ret
715
716      endif
717
718
719      ; PGMUART: Program the Network UART channel
720
721      pgmuart:

```

CP/M RMAC ASSEM 1.1 #016 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

722
723      if      interrupts
724
725      ; The 820 already has the SIO vector address
726      ; programmed from channel B. Other
727      ; implementations will have to provide linkage
728      ; to the vector area in the main XIOS, and
729      ; load the vector offset into SIO write
730      ; register 2
731      01D0 210000        lxi    h,niisr       ; load status interrupt service routine vector
732      01D3 220AFF        shld   siov5
733      01D6 210000        lxi    h,dlissr      ; load transmit ISR vector
734      01D9 220CFF        shld   siov6
735      01DC 210000        lxi    h,reisr       ; load receiv ISR vector
736      01DF 220EFF        shld   siov7
737      endif
738
739      01E2 211C00        lxi    h,sioiblk     ; point to SIO initialization block
740      01E5 060C          mvi    b,sioilen    ; length of block
741      01E7 F3           di
742
743      pgm1:
744
745      01E8 7E           mov    a,m           ; output the block to the SIO
746      01E9 D306          out    siocmd
747      01EB 23           inx    h

```

```

748 01EC 05          dcr    b
749 01ED C2E801     jnz    pgml
750
751 01F0 FB          ei
752 01F1 AF          xra    a          ; set up hailing baud rate = 9600
753 01F2 CD0401     call   setbaud
754 01F5 C9          ret
755
756
757                ; INITUART: Initialize the uart for network A by issuing a reset command
758                ;
759
760                inituart:
761
762 01F6 3E03         mvi    a,3          ; disable the receiver through register 3
763 01F8 D306         out    siocmd
764 01FA 3EC0         mvi    a,disrcv
765 01FC D306         out    siocmd
766
767 01FE DB06         in     siostat      ; is there a garbage byte?
768 0200 E601         ani    rrdymask
769 0202 CA0A02      jz     initu        ; no-->continue initialization
770
771 0205 DB04         in     siorecv      ; else-->eat the character
772 0207 C3F601      jmp    inituart     ; try again
773
774                initu:
775

```

CP/M RMAC ASSEM 1.1 #017 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

776 020A 3E30         mvi    a,errst     ; reset error conditions
777 020C D306         out    siocmd
778
779 020E 3E03         mvi    a,3          ; re-enable the receiver
780 0210 D306         out    siocmd
781 0212 3EC1         mvi    a,enarcv
782 0214 D306         out    siocmd
783
784 0216 C9          ret
785
786                ; INITRECV: Initialize a receive operation
787
788                initrcv:
789
790 0217 CDF601      call   inituart
791
792                if    interrupts
793 021A CDC001      call   enblrcv     ; enable receiver interrupts
794                endif
795
796 021D C9          ret
797
798
799                ; WAIT - Wait 100 micro seconds
800
801                wait:
802
803 021E 3E16         mvi    a,timeval
804
805                w:
806
807 0220 3D          dcr    a          ; 04
808 0221 A7          ana    a          ; 04
809 0222 C22002     jnz    w          ; 12
810                ; ---
811 0225 C9          ret          ; 30 T-States total
812
813
814                ; RESTUART: Reinitialize the UART to the way it was in the
815                ; original BIOS after completing the network operations
816
817
818                restuart:
819 0226 C9          ret          ; UART not used except by network
820
821

```

```

822           ; CSNIOD: Do any cold start initialization which is necessary.
823           ;           Must at least return the value of BAUDS
824           ;           If the network uses the printer port then set the carry flag
825           ;           otherwise clear it.
826
827           csniod:
828
829 0227 01032A           lxi     b,bauds           ; return the legal baud rates

```

CP/M RMAC ASSEM 1.1 #018 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

830 022A B7           ora     a           ; not using a printer port
831 022B C9           ret
832
833 022C           end

```

CP/M RMAC ASSEM 1.1 #019 ULCNET DATA LINK LAYER MP/M XIOS MODULE

```

BAUDGEN      0000    70# 514
BAUDS        2A03    68# 829
BAUDSH       002A    65# 68 530
BAUDSL       0003    64# 68 491
BAUDTBL      0014   141# 508
BTBL         000C   138# 482
CARBIT       0003    80#
CARMSK       0008    81# 686
CHKSTAT      017C    39 610 632#
CONIN        0000    53 242
CONOUT       0000    53 269
CONST        0000    53 216
CSDLL        0000    50 218
CSNIOD       0227    40 827#
CURBAUD      000B   135# 473 517 555
DEVOK        0036   184 189#
DEVTABLE     0005   126# 130 195
DISRCV       00C0    95# 764
DISSLF       004F    97# 504
DLISR        0000    58 733
DLLBAU       0003    42 120#
DLLON        0000    50 219
DSBLRECV     01C9    45 709#
DSBLXMIT     015B    41 552#
DTROFF       006A    93# 499 561
DTRON        00EA    92# 542 566
ENARCV       00C1    94# 781
ENASLF       000F    96# 537
ENBLRECV     01C0    45 699# 793
ERRBITS      0070    83# 638
ERRST        0030    82# 670 776
EXITDL       00CE   306 378#
FALSE        0000    27#
FBIT         0006    88# 661
FMSK         0040    89#
FRMCNTR      0000    52 664
GETRCODE     0000    49 395
GETTCODE     0000    49 316
INCCNTR      0000    52 647 656 665
INITRECV     0217    39 788#
INITU        020A    37 769 774#
INITUART     01F6    38 760# 772 790
INTERRUPTS   FFFF   29# 44 56 155 414 440 578 602 695 723
              792
NCONIN       005C    35 238#
NCONOUT      0075    35 265#
NCONST       0043    35 212#
NETADR       0000    42 119#
NETCON       0020   106# 214 240 267
NETIDLE      01AF    39 603 679#
NETSTATS     FFFF   30# 641
NF           01A9   662 667#
NIISR        0000    58 731
NO           019D   653 658#

```

CP/M RMAC ASSEM 1.1 #020 ULCNET DATA LINK LAYER MP/M XIOS MODULE

Listing F-3: ULCnet Data-link Layer MP/M XIOS Module

Appendix G

Using CP/NET 1.2 with CORVUS OMNINET

Corvus OMNINET is an inexpensive, high-performance CSMA/CA networking system supporting up to 63 hosts on a one-megabit-per second, twisted-pair cable. OMNINET host interface adaptors are intelligent coprocessors that deal with all aspects of network communication of the host in which they are installed, up to and including the transport layer of the ISO open system model. The sample SNIOS and NETWRKIF files following this discussion show one way to use Corvus engineering transporters to implement a CP/NET system.

G.1 The Corvus Engineering Transporter

The Corvus engineering transporter is a card for evaluating Corvus OMNINET with minimum modification to an existing Z80 system. The transporter is not an end-user product, but it is similar enough in hardware design to most production systems using OMNINET to work with little modification.

General information about the Corvus transporter is presented here to help you understand the operation of the sample codes at the end of this appendix. For more information, refer to Corvus documentation.

Communication with the transporter hardware is simplified by the fact that the transporter is microprocessor-based and uses autonomous DMA to access its host computer's memory directly. All communication between host and transporter is controlled by well organized data structures existing in host memory. The only port I/O the host ever does is the transmission, to the transporter hardware, of 24-bit pointer objects (as three serial bytes, most significant byte first) via an output port. Note that all Corvus multibyte objects are in most significant byte first order. These pointer objects refer to transporter command blocks, described in [Table G-1](#).

Field	Size	Explanation
OPERATION COMMAND CODE	8 bits	sends a message.
RESULT BLOCK POINTER	24 bits	gives the address of a data structure for the transporter to update with completion information.
SOCKET CODE	8 bits	defines which of the 4 virtual communication channels to use for this operation.
DATA BUFFER POINTER	24 bits	gives the address of a message buffer for this operation.
DATA LENGTH FIELD	16 bits	gives the length of the message to be transmitted or maximum message length accepted, if this is a receive operation. The maximum length allowed for a single message packet is 2048.
CONTROL FIELD	8 bits	gives the length of an independent auxiliary message that can be sent to a special CONTROL buffer in the destination host at an address different from

LENGTH		that of the destination message buffer. In the case of a receive command, this field specifies the largest such CONTROL message acceptable.
DESTINATION HOST	8 bits	specifies network address of the target host. Legal network addresses are 0-63, or 255 for broadcast messages. A host's address is set by switches connected to the transporter hardware.

Table G-1. Transporter Command Block

Not all fields are used by all commands, but the syntax of the command block is usually consistent, except in the case of special diagnostic commands.

The result pointer in the command block must contain the address of a large enough data structure in host memory to accept the completion information that the specified command produces. Note that the result block is associated with the operation the command block describes. If more than one operation is posted to the transporter hardware, each must have its own result block available. [Table G-2](#) describes a typical result block.

Field	Size	Explanation
OPERATION STATUS CODE	8 bits	set to 254 by the transporter processor once it has read and accepted the command block. This field is later set by the transporter to a result code when it has completed the requested operation.
SOURCE HOST NUMBER	8 bits	gives the network address of the node from which this message packet came.
ACTUAL DATA LENGTH	16 bits	gives the actual length of the message in the receive buffer.
CONTROL MESSAGE BUFFER	0-255 bytes	a buffer large enough to accept any CONTROL message transmitted with the main message packet. The command block that points to this result block must allow such messages.

Table G-2. Receive Result Block

Up to four simultaneous receive operations can be in progress at any one time, waiting for messages for the four logical sockets in the host. Only one message can be posted for transmission at any one time, but this can be done even while four receive operations are pending. Messages from one node are only acceptable to another node if it has a receive command outstanding specifying the socket to which the message is directed.

In use, the host processor must build a command block, then post it to the transporter hardware by outputting one byte at a time of its 24-bit address to the transporter via an output port. The transporter uses an input ready status bit to synchronize this transfer. Command pointers can be transfers done at any time except while the transporter is processing a command block to transmit a message. That operation ties up the transporter until the message has been delivered, or the transporter has given up trying. Network latency is low, so the transporter is unavailable only briefly.

Once the transporter has read and accepted a command, it sets the operation status code in the result block to

254. It is advisable for the host to preset this byte to 255 before sending the transporter the pointer, so that the transporter can confirm that the command was accepted by checking for the change.

The host then polls all active result blocks, waiting for any operation status code to change to a value other than 0FEh. This change means the transporter has completed the operation associated with that result block, and data and result information are available. To simplify interpretation of results, all error codes are between 80h and 0FEh, and all success codes are less than 80h. Send and receive calls that succeed give the number of retries as a completion code, but this code is always less than 7Fh.

OMNINET transporter interfaces usually support generation of a host interrupt whenever the transporter writes to a result block. This relieves the host of having to poll result blocks for completion. To simplify OMNINET evaluation, the engineering transporter is not usually configured to use interrupts. The sample programs demonstrate the use of the transporter both without interrupts and with external interrupt hardware. Servers usually need interrupt hardware or an XIOS polling routine to achieve a usable throughput, but the sample drivers can be made to run without either if high throughput is not a goal.

The coprocessor interface structure the transporter uses is close to the ideal model of a perfect transport layer. The transporter hardware deals with all retries, message acknowledgments, packet sequencing checking, and error detection totally transparently to the host it serves. The data-structure based message interface between the host and transport layer is useful even in implementing non-OMNINET interrupt-driven transport layers for CP/NET.

G.2 Implementation Structure

In the sample implementation, very few OMNINET features were needed. All CP/NET traffic is on one logical channel (SOCKET 2), leaving the others free for such non-CP/NET uses as providing bootstrap channels between diskless devices and optional processes to load them, providing non-CP/NET peripheral sharing routines or even supporting a second network operating system in concurrent use.

Because CP/NET processes its own control fields (message headers), the control message options are not used and are set to zero. In the evaluation transporter, the most significant byte of the memory address is not used and is always set to zero. Other hardware implementations can use this byte for segment control to allow the message buffers to be banked out, or for a 16-bit processor.

The network node ID of an OMNINET host is set by six switches on its transporter hardware. In this implementation, the NODE number is the CP/NET network ID. Set the ID of the SERVER to 00. A requester can have any other unique OMNINET ID code except 0FF hex. This ID code freedom is achieved by a routine in the NETWRKIF module that binds requester ID codes dynamically to processes in the SERVER.RSP module by tracking login and logoff messages. Hence, up to 63 requesters can be supported, as long as no more than NSLAVES are logged in at any one time. Because the transporter handles all low-level communication concerns, the NETWRKIF module is relatively compact; and 16 requesters are easily supported in most systems.

To simplify coding the interface modules, data structure constructor macros eliminate the need for typing all the definitions again and again for each requester. This technique requires that the indices into the resulting arrays of data structures be computed at run-time, but this is easy to do and, where possible, is part of initialization.

G.3 The SNIOS Implementation

The intelligent nature of the OMNINET interface makes coding the SNIOS a simple exercise. Allocate a set of prefabricated transporter command blocks and associated result blocks. Even though the requester never has more than one operation pending at a time, it is simpler to use separate command blocks for each needed

operation type than to recycle the same command block.

Unfortunately, relocating 8080 assemblers like RMAC do not easily deal with relocation of multibyte pointers that are not in Intel® standard memory order. It is simplest to set the result block pointers at initialization; that approach is used here.

After setting up these pointers, the NTWRKINIT routine posts a prebuilt transporter command block called INITTCB to the transporter via the routine called OMNI\$STROBE. If the transporter does not accept the pointer, initialization aborts and an error returns to the NDOS. If the transporter accepts the pointer, NTWRKINIT calls OMNI\$WFDONE to poll the result block associated with INITTCB until the transporter reports a completion. If the initialization operation succeeds, the node number presently set into the transporter's switches is found as a result code. If initialization fails, a value > 80h corresponding to an error code is found and returned to NTWRKINIT, and NTWRKINIT aborts and returns an error code to the NDOS. Otherwise, the node number returned is installed in configtbl and the default message buffer's SID field, the requester ID and a banner print on the console, and a success code is returned to the NDOS.

The NTWRKERROR entry is functionally identical to NTWRKINIT except that it does not print a banner or requester ID code.

The NTWRKSTS, CNFGTBLADR, and NTWRKWBOOT routines are identical in function and operation to those used with other transport layers.

When the NDOS calls the SENDMSG routine, the BC register pair contains a pointer to the message to be sent on the network. This routine translates the CP/NET header information of that message into a form consistent with OMNINET and then puts it into a prefabricated transporter command block called TXTCB. The CP/NET DID is used as the target node physical address on the network. The address of the whole message, including the CP/NET header, is placed in the buffer field of TXTCB after the pointer is rearranged into MSB, LSB sequence. The CP/NET SIZ field is adjusted to give the total message length, including the CP/NET header, and is placed in the appropriate field of the TXTCB.

The OMNINET interface primitives OMNI\$STROBE and OMNI\$WFDONE again post the command to the transporter and, if successful, await completion of the transmission operation. The completion code is transformed into a flag the NDOS expects. Because a very busy server might not have a buffer posted when the requester sends the message, even though the transporter does multiple retries by itself, a retry loop tries to send the message again, if necessary. In practice, retries are rare, but the retry loop is useful when debugging a server.

Like SENDMSG, the RECEIVMSG routine is primarily an exercise in the translation of parameters and their transmission to the transporter. The operation of RECEIVMSG is easily understood by reading its code, with one exception; if a receive is posted, and no message ever comes in, the transporter waits forever for a message. To simplify debugging and recovery from network errors, the OMNI\$WFDONE routine times out after about 20 seconds (on a 2 mhz processor) and returns an error flag to its caller. Most servers ordinarily respond in this time, so the RECEIVMSG routine issues a cancel receive command to the transporter via a prefabricated command block called UNRXTCB. RECEIVMSG then returns to the NDOS with an error code.

If the receive call is not cancelled, an unsolicited or late message might be written into host memory at the requested address long after the host is using that memory for something else. Most autonomous transport layers support this kind of cancellation.

The implementation here is less than 280h bytes long, including the default 138-byte message buffer. If space is tight, the message printing and banner routines can be placed in the default buffer, a single transporter command block and result block can be recycled for all commands, and concessions to modularity can be

made to yield an even smaller SNIOS.

G.4 The NETWRKIF Implementation Model

This sample OMNINET NETWRKIF uses a slightly different intermodule communication model from the one usually used to implement a serial asynchronous star network. Instead of using one process per server process to implement the network input and output, a single input process and a single output process route all messages. This type of structure is far more efficient for any party-line type of network interface hardware because fewer dispatches occur per transaction. Those transactions that do occur take less time and far less code is required to implement the NETWRKIF. In addition, the structure is easier to understand and debug, and all traffic converges through one piece of code, allowing you to implement message routing extensions to your network.

This model is easily understood by studying the general function of the network receiver and transmitter process separately.

The network receiver process in this version is named SERVERX. It is responsible for collecting each incoming message as it arrives, identifying the server process it is for, and writing a pointer to the message into that process's input queue. In addition, SERVERX functions as a surrogate server process to advise requesters that are not logged in that they have no server process to use.

SERVERX uses run-time binding of requester ID codes to server processes. SERVERX does this by keeping a table of the input queue addresses of all the server processes it supports and the ID code of the requester currently logged in to each process. SERVERX examines each incoming messages SID field and searches the table to find out whether SID is presently associated with a server process. If not, an error reply message is constructed in the same buffer that the message arrived in, and SERVERX writes this message directly to the network output process for transmission back to the requester.

For this process to function properly, SERVERX must track all login and logoff messages that pass through it. Every time a login message is received, SERVERX checks its mapping table to find out whether that requester is currently associated with a server process. If it is, no action is taken. If not, SERVERX tries to find an idle server entry in the table. Idle entries are shown in this table as in use by requester 255. If a free server entry is located, SERVERX enters the requester's ID into it, and then sends the login message to that server process's input queue. If none are available, an error reply message is constructed by SERVERX and sent back to the requester.

Logoff messages are handled by finding that requester's server entry, marking it as empty (255) , and then routing the logoff message to the server's input queue. If that requester was never logged in in the first place, SERVERX sends it an error, as previously explained.

Because there is no way to know which server process an incoming message will be for at the time a buffer is posted to the transporter for a receive call, buffers are not permanently assigned to particular server processes. Instead, a list of empty buffers is kept in an MP/M II queue, and SERVERX obtains the buffers from the queue as needed and available for posting to the transporter.

The OMNINET primitives are similar to those used by the SNIOS, except that an MX queue ensures that the transporter is not in use by another process when SERVERX wants to post a command block pointer to it.

As the arrival time of the next message is unknown, SERVERX must be suspended while it waits for the next message to arrive. This can be done by an XDOS flag wait in the WF\$RXDONE OMNINET primitive or by delay-based polling. If your XIOS can be easily modified, another alternative is to add an XIOS polling routine. Using the delay call to suspend the process drastically reduces network throughput because only 60 incoming messages can arrive per second.

The SERVETX process is extremely simple. It reads messages from a single input queue and posts them, using mutual exclusion, to the transporter. Because messages are quickly disposed of by the network, there is no point in suspending SERVETX. It uses a different completion routine than SERVERX, which merely waits until a completion code is received from the transporter, and then returns to its caller. To simplify debugging, a timeout is included to prevent a hardware or software problem from locking up the system.

Once SERVETX has finished sending the message, it returns the buffer that it was in to the free buffer management queue, making it available for SERVERX. SERVETX then goes back to read its input queue to wait for another message to process.

Theoretically, such a system can function with fewer buffers than server processes. But in practice, it is best to have at least one more buffer than the number of server processes in the pool to deal with messages such as failed login attempts that never get routed to a server.

The rest of the code in each process simply initializes data structures, creates queues, initializes hardware, and performs other routine tasks.

Note that the distribution version of CP/NET 1.2 does not work with this SERVETX process without a minor patch. SERVER.RSP must be patched to create output UQCBs with the same name for all server processes instead of making each queue name unique. Once this is done, all processes in SERVER.RSP direct their output to a single SERVETX process. Instructions for installing this patch are included in *CP/NET V1.2 Application Note 02*.

G.5 Possible Improvements to NETWRKIF

This interface is by no means ideal. Little error recovery is done for registers that fail to log off. A watchdog timing process can be easily added to correct this problem. This process is not shown here, to simplify understanding of the OMNINET interface. But such a process is only needed in systems with more physical requesters than server processes to prevent their being locked up by departed users.

One possible improvement is to further reduce the number of dispatches per CP/NET transaction by using direct code to manage the buffer list and using the transporter mutual exclusion function instead of the MP/M II queue facility. The M/PM II queue facility is powerful and easy to use, but avoid using it in situations where dispatch overhead exceeds the time for which a process is likely to require suspension unless the suspension is unavoidable for process synchronization reasons.

Another worthwhile improvement is to modify the NETWRKIF to minimize the period during which the server cannot respond to incoming messages, by seeing that the next buffer is more quickly posted for the next received message after a receive completion occurs. The present version does not do this until the incoming message has been processed by SERVERX. This causes unneeded network traffic because messages sent by requesters during this time are futile.

High-performance servers can make good use of two physical sets of transporter hardware, with different node addresses, on the same loop. Using two transporters can totally bypass the need to use MX techniques because one transporter can be reserved solely for transmitting messages.

Interesting networks can be easily constructed by having more than one OMNINET loop, each with its own transporter. The SERVERX process associated with each loop can filter messages not intended for local SLVSPs to a second, third, or fourth SERVETX process associated with higher level loops. Such filtering bridges can be used to build hierarchical CP/NET systems of any degree of complexity.

Other processes can concurrently send and receive messages totally unrelated to the CP/NET context using the same transporter as long as they honor the MXomni mutual exclusion queues and do not use the same

socket for their communication as CP/NET. These processes can implement a variety of supervisory and auxiliary functions, or they can implement additional concurrent virtual circuits that cooperating requesters can use for point-to-point traffic. Such point-to-point virtual circuits can be coordinated by CP/NET mail functions.

```

CP/M RMAC ASSEM 1.1 #001 SAMPLE SLAVE NETWORK I/O SYSTEM FOR CORVUS OMNINET 20 OCT 82

1 title 'Sample Slave Network I/O System for CORVUS OMNINET 20 Oct 82'
2 page 54
3
4 ;
5 ;
6 ;
7 ; SAMPLE SLAVE NETWORK IO SYSTEM FOR CP/NET 1.2 ;
8 ; VERSION FOR CORVUS OMNINET "ENGINEERING" TRANSPORTER ;
9 ; (Requires RMAC for assembly) ;
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32 ; FRANCE ;
33 ;
34 ; who sponsored the development of one of its ancestors. ;
35 ;
36 ;
37 ;
38 ; ***** CONSTANT DECLARATIONS *****
39 ;
40 0000 = FALSE equ 0
41 FFFF = TRUE equ not FALSE
42 ;
43 ; configuration and option constants
44 0064 = TXTRIES equ 100 ;Transmit message retries
45 008A = BUFFSIZE equ 138 ;max default buffer size
46 0200 = MAXMSG equ 512 ;largest message accepted by receiver
47 0080 = SKT0 equ 80h ;legal omninet socket tokens
48 0090 = SKT1 equ 90h
49 00A0 = SKT2 equ 0a0h
50 00B0 = SKT3 equ 0b0h
51 00A0 = SOCKET equ SKT2 ;this SNIOS uses only channel 2
52 ;
53 ; OMNINET Constants
54 ; Completion/return codes

CP/M RMAC ASSEM 1.1 #002 SAMPLE SLAVE NETWORK I/O SYSTEM FOR CORVUS OMNINET 20 OCT 82

55 0000 = NOERR equ 0 ;done (no errors or retries)
56 00C0 = ETXOK equ 0c0h ;echo succeeded with no retries (not used here)
57 0080 = ETXFAIL equ 80h ;Transmit failed
58 0081 = E2LONG equ 81h ;wouldn't fit in destination socket
59 0082 = ENOSKT equ 82h ;destination socket not set up
60 0083 = EBDCTL equ 83h ;bad control field length
61 0084 = EBDSKT equ 84h ;illegal socket number
62 0085 = EBDDDES equ 85h ;invalid destination node number/socket in use
63 0086 = EBDNODE equ 86h ;bad node number in command (not 0-7fh or ffh)
64 00FE = ECMDOK equ 0feh ;command has been read by transporter

```

```

65      ; legal command tokens
66 0040 = SENDF equ 40h      ;send message
67 00F0 = RCVF equ 0f0h    ;set up receive socket
68 0010 = ENDRCVF equ 10h  ;stop receive
69 0020 = INITF equ 20h    ;initialize transporter
70      ; Transporter control ports
71 00F8 = NETBASE equ 0f8h  ;base address of transporter IO interface
72 00F9 = TSTAT equ Netbase+1 ;ready status port
73 0010 = TCRDY equ 10h    ;status mask for ready bit
74 00F8 = TDATA equ Netbase ;command block pointer port
75
76      ; Network Status Byte Constants
77      ;
78 0010 = ACTIVE equ 10h    ;slave logged in on network
79 0002 = RCVERR equ 2h    ;error in received message
80 0001 = SENDERR equ 1h   ;unable to send message
81
82      ; CP/M BDOS function constants
83 0005 = BDOS equ 5        ;absolute BDOS entry
84 0009 = PRINTF equ 9      ;print message function
85 0002 = CONOUTF equ 2     ;output char in E to console
86
87      ; General Constants
88 000A = LF equ 0ah        ;Line Feed
89 000D = CR equ 0dh        ;Carriage Return
90
91      ; ***** GENERATED CODE AND DATA BEGIN HERE *****
92
93      ; Public Jump vector for SNIOS entry points
94 0000 C3F400 jmp ntwrkinit ;network initialization
95 0003 C34801 jmp ntwrksts ;network status
96 0006 C35201 jmp cnfgtbladr ;return config table addr
97 0009 C36701 jmp sendmsg ;send message on network
98 000C C3A601 jmp receivevmsg ;receive message from network
99 000F C33801 jmp ntwrkerror ;network error
100 0012 C35601 jmp ntwrkwboot ;network warm boot
101
102      ; Public Slave Configuration Table
103 configtbl:
104 Network$status:
105 0015 00 db 0 ;network status byte
106 0016 00 slvid1: db 0 ;slave ID (from switches)
107 0017 000000000000 db 0,0, 0,0, 0,0, 0,0 ;Disk map table for units A:-P:
108 001F 000000000000 db 0,0, 0,0, 0,0, 0,0

```

CP/M RMAC ASSEM 1.1 #003 SAMPLE SLAVE NETWORK I/O SYSTEM FOR CORVUS OMNINET 20 OCT 82

```

109 0027 000000000000 db 0,0, 0,0, 0,0, 0,0
110 002F 000000000000 db 0,0, 0,0, 0,0, 0,0
111 0037 0000 db 0,0 ;console device
112 0039 0000 db 0,0 ;list device
113 003B 00 db 0 ;buffer index
114
115 003C 00 dflt: db 0 ;FMT (DEFAULT MESSAGE BUFFER)
116 003D 00 db 0 ;DID
117 003E 00 slvid2: db 0 ;SID
118 003F 05 db 5 ;FNC
119 0040 00 db 0 ;SIZ
120 0041 ds 1 ;MSG(0) List number
121 0042 ds BUFSIZE ;MSG(1) ... MSG(128)
122
123
124      ; ***** PREFABRICATED OMNINET TRANSPORTER COMMAND BLOCKS *****
125
126      ; Command block for transmitting a message
127 TXtcb:
128 00CC 40 TXtcmd: db SENDF ;command field
129 00CD 00 db 0 ;bits 16-24 of result block ptr
130 00CE 0000 TXtrslt: db 0,0 ;result block pointer (MSB,LSB)
131 00D0 A0 TXtskt: db SOCKET ;socket (channel) number
132 00D1 00 db 0 ;bits 16-24 of message buffer ptr
133 00D2 0000 TXtmsg: db 0,0 ;message buffer pointer (MSB,LSB)
134 00D4 0000 TXtdlen: db 0,0 ;data field length (MSB,LSB)
135 00D6 00 TXtclen: db 0 ;control field length
136 00D7 00 TXtdest: db 0 ;Destination address (transport layers)
137      ; Result vector for above command block
138 TXresult:

```

```

139 00D8 00      TXrcode:      db      0          ;return code
140
141              ;      Command block for setting up a receive operation
142      RXtcb:
143 00D9 F0      RXtcmd:      db      RCVF          ;command field
144 00DA 00              db      0
145 00DB 0000     RXtrslt:      db      0,0          ;result block pointer (MSB,LSB)
146 00DD A0      RXtskt:      db      SOCKET        ;socket number
147 00DE 00              db      0
148 00DF 0000     RXtmsg:      db      0,0          ;message address (MSB,LSB)
149 00E1 02      RXtdlen:     db      MAXMSG/256      ;max data field length (MSB,LSB)
150 00E2 00              db      MAXMSG and 255
151 00E3 00      RXtclen:     db      0          ;max control field length
152 00E4 00      RXtdest:     db      0          ;(not used in a receive operation)
153              ;      Result vector for receiver
154      RXresult:
155 00E5 00      RXrcode:      db      0          ;return code
156 00E6 00      RXrsrce:     db      0          ;source HOST #
157 00E7 0000     RXrdlen:     db      0,0          ;received message length (MSB,LSB)
158
159              ;      Command block for receive cancel operation
160      UNRXtcb:
161 00E9 10      UNRXtcmd:    db      ENDRCVF        ;command field
162 00EA 00              db      0

```

CP/M RMAC ASSEM 1.1 #004 SAMPLE SLAVE NETWORK I/O SYSTEM FOR CORVUS OMNINET 20 OCT 82

```

163 00EB 0000     UNRXtrslt:    db      0,0          ;result block pointer (MSB,LSB)
164 00ED A0      UNRXtskt:    db      SOCKET        ;socket number
165              ;      Result vector for receive cancel
166      UNRXresult:
167 00EE 00      UNRXrcode:   db      0          ;return code
168
169              ;      Command block for transporter initialization command
170      INITtcb:
171 00EF 20      INITtcmd:   db      INITF          ;command field
172 00F0 00              db      0
173 00F1 0000     INITtrslt:   db      0,0          ;result block pointer (MSB,LSB)
174              ;      Result vector for initialization
175      INITresult:
176 00F3 00      INITrcode:  db      0          ;return code (if valid,=ID code)
177
178
179              ;      ***** PUBLIC CODE ENTRIES BEGIN HERE *****
180
181              ;      Externally accessed routine to initialize transporter
182              ;      (RETURNS A=0 if succeeds, else 0ffh.)
183      ntwrkinit:
184 00F4 CD3801   call    ntwrkerror      ;init transporter, tcbs and id code
185 00F7 D8              rc          ;return error if init fails
186 00F8 110601   lxi    d,initmsg       ;else prinw slave ID and banner
187 00FB CDF001   call    print$msg      ;
188 00FE 3A1600   lda    slvid1          ;
189 0101 CDD601   call    prhex          ;print slave ID
190 0104 AF      xra    a              ;and return to caller with a=0
191 0105 C9      ret
192
193      initmsg:
194 0106 0D0A534E49 db    CR,LF,'SNIOS (c)1982 Vano Associates Inc.'
195 012A 0D0A534C41 db    CR,LF,'SLAVE ID = $'
196
197
198              ;      Externally accessed routine inits or re-inits module
199              ;      (RETURNS A=0 if succeeds, else 0ffh.)
200      ntwrkerror:
201 0138 AF      xra    a              ;zero network status byte
202 0139 321500   sta    Network$status  ;
203 013C CDF501   call    omni$init      ;init transporter, tcbs and id code
204 013F D8              rc          ;carry means error, A=0ffh
205 0140 321600   sta    slvid1          ;update this slaves id in table
206 0143 323E00   sta    slvid2          ;and default message
207 0146 AF      xra    a              ;and return with no error
208 0147 C9      ret
209
210
211              ;      Externally accessed routine returns Network Status Byte in A
212              ;      (also clears any error bits active)

```

```

213          ntwrksts:
214 0148 211500      lxi    h,network$status
215 014B 46          mov    b,m
216 014C 3EFC        mvi    a,not(RCVERR or SENDERR)

```

CP/M RMAC ASSEM 1.1 #005 SAMPLE SLAVE NETWORK I/O SYSTEM FOR CORVUS OMNINET 20 OCT 82

```

217 014E A0          ana    b
218 014F 77          mov    m,a
219 0150 78          mov    a,b
220 0151 C9          ret
221
222
223          ; Externally accessed routine Returns Configuration Table Ptr in HL
224          cnfgtbladr:
225 0152 211500      lxi    h,configtbl
226 0155 C9          ret
227
228
229          ; Externally accessed routine is called each time the CCP is reloaded
230          ; from disk. (Dummy procedure for now.)
231          ntwrkwboot:
232 0156 115C01      lxi    d,wboot$msg          ;return via print$msg
233 0159 C3F001      jmp    print$msg
234
235          wboot$msg:
236 015C 0D0A3C4350 db    CR,LF,'$'
237
238
239          ; Externally accessed routine sends Message BC--> on Network
240          ; (returns A=0 if succeeds, else A=0ffh.)
241          ;
242          ; NOTE that although the OMNINET transporter does its own transport
243          ; layer retries, this routine does additional retries to deal with
244          ; servers that are slow in posting receive calls since transport
245          ; level retries are exhausted in a very short real-time period.
246          sendmsg:
247 0167 61          mov    h,c          ;move buffer pointer to Transporter ctrl block
248 0168 68          mov    l,b          ;(note reversed byte order for Transporter.)
249 0169 22D200      shld  TXtmsg
250
251          ;
252 016C 210400      lxi    h,4          ;get CP/Net message length from SIZ field
253 016F 09          dad    b
254 0170 6E          mov    l,m
255 0171 2600        mvi    h,0
256 0173 110600      lxi    d,6          ;add packet header lgth to get actual size
257 0176 19          dad    d          ; of packet for transport layer purposes
258 0177 7C          mov    a,h          ;swap bytes to MSB, LSB order
259 0178 65          mov    h,l
260 0179 6F          mov    l,a
261 017A 22D400      shld  TXtdlen      ;store length in TCB data length field
262
263          ;
264 017D 03          inx    b          ;get DID from message
265 017E 0A          ldax  b
266 017F 32D700      sta  TXtdest      ;put it into TCB destination address field
267
268          ;
269 0182 116400      lxi    d,TXTRIES   ;use DE as retry counter
270
271          ;
272          send$again:
273          ;head of message transmission retry loop
274          push  d
275          lxi  b,TXtcb ;send TCB pointer to transporter hardware

```

CP/M RMAC ASSEM 1.1 #006 SAMPLE SLAVE NETWORK I/O SYSTEM FOR CORVUS OMNINET 20 OCT 82

```

271 0189 CD2E02      call  omni$strobe
272 018C D1          pop    d
273 018D DAA101      jc    snderr       ;if not accepted, goto fatal error handler
274
275          ;
276 0190 01D800      lxi    b,TXresult  ;else poll result block until completion code
277 0193 D5          push  d          ;is returned by hardware
278 0194 CD5C02      call  omni$wfdone
279 0197 D1          pop    d
280
281          ;
282 0198 E680        ani    80h        ;completion codes 80h-ffh are error codes
283 019A C8          rz              ;return 00 to caller if no errors

```



```

283 019B 1B      dcx    d      ;else decrement retry counter
284 019C 7B      mov    a,e
285 019D B2      ora    d
286 019E C28501  jnz    send$again      ;retry transmit if any retries left
287              ;
288 01A1 3E01      snderr: mvi   a,SENDERR      ;goto common exit code to update error flags
289 01A3 C3CE01    jmp    nerr      ;(part of receivemsg routine)
290
291
292              ;
293              ;   Externally accessed routine waits for a message directed to this node
294              ;   and returns it in the buffer BC-->. To aid debugging, a timeout of
295              ;   about 20 seconds (2 Mhz processor) is implemented that will return an
296              ;   error if no message is received. That is long enough for most normal
297              ;   servers to respond.
298              ;
299              ;   (RETURNS A=0 if good msg, =0ffh if bad msg or timeout.)
300 01A6 68      receivemsg: mov    l,b      ;swap buffer pointer bytes to MSB,LSB order
301 01A7 61      mov    h,c
302 01A8 22DF00  shld   RXtmsg      ;put buffer ptr to its TCB field
303              ;
304 01AB 01D900  lxi    b,RXtcb
305 01AE CD2E02  call   omni$strobe ;post control block address to hardware
306 01B1 DACC01  jc     rxerr      ;fatal error if hardware won't accept it
307              ;
308 01B4 01E500  lxi    b,RXresult
309 01B7 CD5C02  call   omni$wfdone ;else wait for a completion from hardware
310 01BA E680    ani    80h
311 01BC C8      rz          ;return 00 to caller if no error reported
312              ; the rest is the fatal error handler for receive calls
313 01BD 01E900  lxi    b,UNRXtcb  ;otherwise cancel the receive call
314 01C0 CD2E02  call   omni$strobe ; (using prefabricated cancel command block)
315 01C3 D2CC01  jnc   rxerr      ;If won't accept this command either, quit here
316              ;
317 01C6 01EE00  lxi    b,UNRXresult ;else wait for completion of cancel command
318 01C9 CD5C02  call   omni$wfdone ;ignore result (always fatal error return)
319 01CC 3E02      rxerr: mvi   a,RCVERR ;exit via code that updates status byte
320
321              ; This is also used by sendmsg to update Network$status and return 0ffh
322 01CE 211500  nerr:  lxi    h,Network$status
323 01D1 B6      ora    m
324 01D2 77      mov    m,a      ;update status

```

CP/M RMAC ASSEM 1.1 #007 SAMPLE SLAVE NETWORK I/O SYSTEM FOR CORVUS OMNINET 20 OCT 82

```

325 01D3 3EFF      mvi    a,0ffh
326 01D5 C9      ret          ;return 0ffh to caller
327
328
329              ; ***** UTILITY ROUTINES CALLED BY ABOVE BEGIN HERE *****
330
331              ; prints A in hex on console
332 01D6 F5      prhex:  push   psw
333 01D7 07      rlc
334 01D8 07      rlc
335 01D9 07      rlc
336 01DA 07      rlc
337 01DB CDDF01  call   nibl      ;print high nibble
338 01DE F1      pop     psw      ;and fall through to print low nibble
339
340 01DF E60F      nibl:   ani    0fh
341 01E1 C630      adi    '0'
342 01E3 FE3A      cpi    '9'+1
343 01E5 DAEA01  jc     printa
344 01E8 C607      adi    7
345 01EA 5F      printa: mov    e,a
346 01EB 0E02      mvi    c,CONOUTF
347 01ED C30500  jmp    BDOS      ;print ascii and return
348
349
350              ; print message DE--> until $ on console device
351 01F0 0E09      print$msg: mvi   c,PRINTF      ;prints $ delimited string DE-->
352 01F2 C30500  jmp    BDOS      ;bdos(sprintf,wboot$msg)
353
354
355
356              ; ***** LOW LEVEL OMNINET TRANSPORTER DRIVERS BEGIN HERE *****

```

```

357
358           ;      Initialize transporter and return its ID code in A or 0ffh if can't.
359           ;      Carry is also set if error, clear if no error.
360   omni$init:           ;initialize pointers in our control blocks
361   01F5 11D800         lxi      d,TXresult      ;NOTE: this is done at run time to avoid
362   01F8 63             mov      h,e             ; relocation problems caused by the need to
363   01F9 6A             mov      l,d             ; have pointers for CORVUS transporter use
364   01FA 22CE00         shld    TXtrslt         ; in MSB, LSB form instead of 8080 format.
365
366           ;
366   01FD 11E500         lxi      d,RXresult
367   0200 63             mov      h,e
368   0201 6A             mov      l,d
369   0202 22DB00         shld    RXtrslt
370
371           ;
371   0205 11EE00         lxi      d,UNRXresult
372   0208 63             mov      h,e
373   0209 6A             mov      l,d
374   020A 22EB00         shld    UNRXtrslt
375
376           ;
376   020D 11F300         lxi      d,INITresult
377   0210 63             mov      h,e
378   0211 6A             mov      l,d

```

CP/M RMAC ASSEM 1.1 #008 SAMPLE SLAVE NETWORK I/O SYSTEM FOR CORVUS OMNINET 20 OCT 82

```

379   0212 22F100         shld    INITtrslt
380
381           ;
381   0215 01EF00         lxi      b,INITtcb      ;send init command block pointer to transporter
382   0218 CD2E02         call    omnistrobe      ;to reset it and get its ID code
383   021B 9F             sbb     a               ;in case of error, preset return code 0 or ff
384   021C D8             rc              ;fatal error if hardware won't accept pointer
385
386           ;
386   021D 01F300         lxi      b,INITresult   ;else wait for result of operation
387   0220 CD5C02         call    omni$wfdone     ;wait for done
388   0223 321600         sta     slvid1          ;result code should be ID code so put in table
389   0226 323E00         sta     slvid2          ;and in default message SID
390
391           ;
391   0229 07             rlc              ;set CY=bit 7 of return code
392   022A 1F             rar              ;so CY=1 if error
393   022B D0             rnc              ;return with ID code if no error
394   022C 9F             sbb     a         a     ;else set carry=1 and A=0ffh and return
395   022D C9             ret
396
397
398           ;
398           ;      Sends the 16 bit POINTER in BC to the transporter hardware as
399           ;      a 24 bit pointer (MSB first). Returns CY set if hardware will
400           ;      not accept any byte in a reasonable time else CY clear.
401   omni$strobe:
402   022E 210200         lxi      h,2           ;Find address of rslt block from TCB BC-->
403   0231 09             dad     b             ;pre-set result code in block to ff (busy)
404   0232 7E             mov     a,m
405   0233 23             inx     h
406   0234 6E             mov     l,m
407   0235 67             mov     h,a
408   0236 36FF         mvi     m,0ffh
409
410           ;
410   0238 AF             xra     a             ;MSB is always 0
411   0239 CD4302         call    omni$st        ;send bits 23-16 of pointer to hardware
412   023C D8             rc              ;(abort if timeout)
413
414           ;
414   023D 78             mov     a,b           ;send bits 15-8 of pointer to hardware
415   023E CD4302         call    omni$st        ;(abort if timeout)
416   0241 D8             rc
417
418           ;
418   0242 79             mov     a,c           ;send bits 7-0 of pointer to hardware
419                           ; (fall into omni$st)
420
421           ;
421           ;      called by omni$strobe to send one byte from A to transporter hardware
422           ;      returns CY set if hardware doesn't come ready in a reasonable time.
423   omni$st:
424   0243 F5             push    psw           ;save data for now
425   0244 1150C3         lxi     d,50000       ;set timeout
426
427   omni$st0:
427   0247 DBF9             in      TSTAT         ;read status port and check busy bit
428   0249 E610             ani     TCRDY
429   024B CA5302         jz     omni$st1      ;if busy, go increment and test timeout
430

```

431	024E F1	pop	psw	;else output the byte
432	024F D3F8	out	TDATA	;to the transporter TCB pointer input register

CP/M RMAC ASSEM 1.1 #009 SAMPLE SLAVE NETWORK I/O SYSTEM FOR CORVUS OMNINET 20 OCT 82

433	0251 B7	ora	a	
434	0252 C9	ret		;and return with no error shown (CY=0)
435				;
436		omni\$st1:		;else
437	0253 1B	dcx	d	
438	0254 7A	mov	a,d	
439	0255 B3	ora	e	
440	0256 C24702	jnz	omni\$st0	;loop back if not timed out yet
441				;
442	0259 F1	pop	psw	;else
443	025A 37	stc		
444	025B C9	ret		;return error flag (CY=1)
445				
446				
447				; waits till timeout (about 20 secs) for result block BC--> to show done
448				; returns A=returned status code. If timeout occurs, the returned
449				; status will still be 0FEH or 0FFH.
450		omni\$wfdone:		
451	025C 11FFFF	lxi	d,0ffffh	;setup timeout counters
452	025F 2E14	mvi	l,20	
453				;
454		omni\$wfdone1:		
455	0261 0A	ldax	b	;is the result code still > 0f0h?
456	0262 FEF0	cpi	0f0h	
457	0264 D8	rc		;no, return to caller
458				;
459	0265 1B	dcx	d	;else decrement timeout
460	0266 7B	mov	a,e	
461	0267 B2	ora	d	
462	0268 C26102	jnz	omni\$wfdone1	;timeout yet?
463	026B 2D	dcr	l	
464	026C C26102	jnz	omni\$wfdone1	;no, go back and check again
465				;
466	026F 0A	ldax	b	;yes, timeout
467	0270 C9	ret		;return with completion code in A
468				
469				
470	0271	end		

CP/M RMAC ASSEM 1.1 #010 SAMPLE SLAVE NETWORK I/O SYSTEM FOR CORVUS OMNINET 20 OCT 82

ACTIVE	0010	78#		
BDOS	0005	83#	347	353
BUFFSIZE	008A	45#	121	
CNFGTBLADR	0152	96	224#	
CONFIGTBL	0015	103#	225	
CONOUTF	0002	85#	346	
CR	000D	89#	194	195 236
DFLT	003C	115#		
E2LONG	0081	58#		
EBDCTL	0083	60#		
EBDDDES	0085	62#		
EBDNODE	0086	63#		
EBDSKT	0084	61#		
ECMDOK	00FE	64#		
ENDRCVF	0010	68#	161	
ENOSKT	0082	59#		
ETXFAIL	0080	57#		
ETXOK	00C0	56#		
FALSE	0000	40#	41	
INITF	0020	69#	171	
INITMSG	0106	186	193#	
INITRCODE	00F3	176#		
INITRESULT	00F3	175#	376	386
INITTCB	00EF	170#	381	
INITTCMD	00EF	171#		
INITTRSLT	00F1	173#	379	
LF	000A	88#	194	195 236
MAXMSG	0200	46#	149	150
NERR	01CE	289	322#	
NETBASE	00F8	71#	72	74

NETWORKSTATUS	0015	104#	202	214	322		
NIBL	01DF	337	340#				
NOERR	0000	55#					
NTWRKERROR	0138	99	184	200#			
NTWRKINIT	00F4	94	183#				
NTWRKSTS	0148	95	213#				
NTWRKWBOOT	0156	100	231#				
OMNIINIT	01F5	203	360#				
OMNIST0	0243	411	415	423#			
OMNIST0	0247	426#	440				
OMNIST1	0253	429	436#				
OMNISTROBE	022E	271	305	314	382	401#	
OMNIWFDONE	025C	277	309	318	387	450#	
OMNIWFDONE1	0261	454#	462	464			
PRHEX	01D6	189	332#				
PRINTA	01EA	343	345#				
PRINTF	0009	84#	352				
PRINTMSG	01F0	187	233	351#			
RCVERR	0002	79#	216	319			
RCVF	00F0	67#	143				
RECEIVMSG	01A6	98	299#				
RXERR	01CC	306	315	319#			
RXRCODE	00E5	155#					
RXRDLN	00E7	157#					

CP/M RMAC ASSEM 1.1	#011	SAMPLE SLAVE NETWORK I/O SYSTEM FOR CORVUS OMNINET 20 OCT 82					
RXRESULT	00E5	154#	308	366			
RXRSRCE	00E6	156#					
RXTCB	00D9	142#	304				
RXTCLEN	00E3	151#					
RXTCMD	00D9	143#					
RXTDEST	00E4	152#					
RXTDLN	00E1	149#					
RXTMSG	00DF	148#	302				
RXTRSLT	00DB	145#	369				
RXTSKT	00DD	146#					
SENDAGAIN	0185	268#	286				
SENDERR	0001	80#	216	288			
SENDF	0040	66#	128				
SENDMSG	0167	97	246#				
SKT0	0080	47#					
SKT1	0090	48#					
SKT2	00A0	49#	51				
SKT3	00B0	50#					
SLVID1	0016	106#	188	205	388		
SLVID2	003E	117#	206	389			
SNDERR	01A1	273	288#				
SOCKET	00A0	51#	131	146	164		
TCRDY	0010	73#	428				
TDATA	00F8	74#	432				
TRUE	FFFF	41#					
TSTAT	00F9	72#	427				
TXRCODE	00D8	139#					
TXRESULT	00D8	138#	275	361			
TXTCB	00CC	127#	270				
TXTCLEN	00D6	135#					
TXTCMD	00CC	128#					
TXTDEST	00D7	136#	264				
TXTDLN	00D4	134#	260				
TXTMSG	00D2	133#	249				
TXTRIES	0064	44#	266				
TXTRSLT	00CE	130#	364				
TXTSKT	00D0	131#					
UNRXRCODE	00EE	167#					
UNRXRESULT	00EE	166#	317	371			
UNRXTCB	00E9	160#	313				
UNRXTCMD	00E9	161#					
UNRXTRSLT	00EB	163#	374				
UNRXTSKT	00ED	164#					
WBOOTMSG	015C	232	235#				

Listing G-1. Sample Slave Network I/O System for Corvus OMNINET

CP/M RMAC ASSEM 1.1	#001	SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82					
---------------------	------	--	--	--	--	--	--

```

1          title 'Sample Server Network I/F for CORVUS OMNINET 20-Oct-82'
2          page 54
3
4          ;
5          ;
6          ;
7          ; SAMPLE MASTER NETWORK IO SYSTEM FOR CP/NET 1.2
8          ; VERSION FOR CORVUS OMNINET "ENGINEERING" TRANSPORTER
9          ; (Requires RMAC for assembly)
10         ;
11         ; COPYRIGHT (C) 1982 by VANO ASSOCIATES, INC.
12         ; P.O. BOX 12730
13         ; New Brighton, MN 55112
14         ; U.S.A.
15         ; (612) 631-1245
16         ; ALL RIGHTS RESERVED
17         ;
18         ; ANY USE OF THIS CODE without the imbedded copyright notice
19         ; is hereby strictly prohibited.
20         ;
21         ; Permission is hereby granted to Digital Research Inc. to use
22         ; this source file for educational and illustrative purposes in
23         ; conjunction with CP/Net 80 documentation. Any other use of
24         ; this code without the EXPRESS WRITTEN PERMISSION of VANO
25         ; ASSOCIATES INC. is hereby strictly prohibited.
26         ;
27         ; This file is provided courtesy of:
28         ;
29         ; R2E (Realisations Etude Electroniques)
30         ; Z.A.I. de Courtaboeuf
31         ; BP 73 91942 Les Ulis
32         ; FRANCE
33         ;
34         ; who sponsored the development of one of its ancestors.
35         ;
36         ; Note that this version requires that the CP/NET SLAVESP
37         ; process be properly patched to send all output traffic
38         ; to output queue 0. For the current (1.2) beta release, the
39         ; following patch is enough:
40         ;
41         ; Make this change in unrelocated server.rsp module.
42         ; -a543
43         ; 0543 mvi a,30
44         ; 0545 jmp 34f
45         ; Then resave the module and its bit map.
46         ;
47         ;
48         ;
49         ;
50         FFFF = YES equ 0ffffh
51         0000 = NO equ not YES
52
53         ; assembly mode switches
54         0000 = DEBUG equ NO ;assemble for debugging with rdt

```

```

CP/M RMAC ASSEM 1.1 #002 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82
55 FFFF = RSP equ YES ;assemble as a resident process
56 0000 = INTERRUPT equ NO ;transporter can interrupt (advisable)
57
58 ; Logical Configuration constants
59 0002 = NSLAVES equ 2 ;maximum number of slaves supported
60 0096 = SRVR$STK$SIZ equ 150 ;stack size needed by SLVSPs
61 0034 = SRVR$PD$SIZ equ 52 ;PD size for SLVSPs
62 0118 = BUFFSIZE equ 280 ;maximum message buffer size
63 0003 = NMSG$BUFFS equ 1+NSLAVES ;number of message buffers allocated
64 0040 = RX$PRIORITY equ 64 ;receive process priority
65 003F = TX$PRIORITY equ 63 ;usually higher than rx
66
67 ; Physical configuration constants (FOR OUR INSTALLATION)
68 00F8 = OMNI$BASE equ 0F8h ;transporter base address
69 00A0 = OMNI$SOCKET equ 0a0h ;omninet transporter socket code (2)
70 0008 = OMNI$FLAG equ 8 ;XDOS flag for int. driven transporter
71 0007 = RST$NUM equ 7 ;interrupt level if interrupt driven
72 0038 = INT$VCTR equ RST$NUM * 8
73

```

```

74      ;      transporter IO PORT constants for CORVUS "ENGINEERING" transporter
75 00F8 =      OMNI$DATA      equ      OMNI$BASE      ;TCB pointer data port
76 00F9 =      OMNI$STAT      equ      OMNI$BASE + 1  ;status port
77 0010 =      OMNI$RDY      equ      10h      ;ready bit (=1) in OMNI$STAT
78      ; the rest are not part of standard CORVUS "ENGINEERING" transporter
79 00FA =      OMNI$ACK      equ      OMNI$BASE + 2  ;int ack port (any data write)
80 00FB =      OMNI$MASK      equ      OMNI$BASE + 3  ;int mask port (b0, 1= enbl)
81 0001 =      OMNI$PENDING  equ      1          ;int pending (=1) in " "
82 0001 =      OMNI$ENABLE   equ      1          ;int enable mask command
83 0000 =      OMNI$DISABLE  equ      0          ;int disable mask command
84
85      ;      BDOS and XDOS Equates
86 0009 =      PRINTF      equ      9          ;message to console
87 0084 =      FLAGWAITF    equ      132         ;flag wait
88 0085 =      FLAGSETF    equ      133         ;flag set
89 0086 =      MAKEQ      equ      134         ;make queue
90 0089 =      READQ      equ      137         ;read queue
91 008B =      WRITEQ     equ      139         ;write queue
92 008D =      DELAY      equ      141         ;delay
93 008E =      DSPTCH     equ      142         ;dispatch
94 0090 =      CREATEP    equ      144         ;create process
95 0091 =      SET$PRIORITY equ      145         ;set caller's priority
96 0093 =      DETACH     equ      147         ;detach console
97 009A =      SYDATAD    equ      154         ;get system data page address
98
99      ;      MISC useful constants
100 000D =      CR          equ      0dh        ;carriage return
101 000A =      LF          equ      0ah        ;line feed
102
103
104      codeseg:
105      if not RSP
106      ;      .PRL Initialization entry point for whole module
107      lxi      sp,ServerxSTKTOP      ;switch to rx process stack
108      mvi      c,SET$PRIORITY

```

CP/M RMAC ASSEM 1.1 #003 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

109      mvi      e,RX$PRIORITY
110      call     bdos
111      if not DEBUG
112      mvi      c,DETACH
113      call     bdos      ;detach console
114      endif ; DEBUG
115      ret
116
117      bdosadr:
118      dw      codeseg - 100h + 5      ;bdos entry pointer
119      else ; not RSP
120      ;      in an rsp, this is filled in by GENSYS and the rx process is created
121      ;      automatically
122      bdosadr:
123      0000 0000      dw      0000h
124      endif ; not RSP
125
126      page

```

CP/M RMAC ASSEM 1.1 #004 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

127
128      ;
129      ;
130      ;
131      ;      This is the network receiver server process module
132      ;
133      ;      The receive server obtains a buffer from FreeBuff and gives it
134      ;      to the transporter hardware for receive use. It then waits
135      ;      for a message completion by calling the wf$rx$done routine
136      ;      Once a return from that routine occurs, the receiver server
137      ;      checks the slave number and sends a pointer to that message
138      ;      buffer to the SLVSP support process corresponding to that
139      ;      slave's server. Once the message pointer has been passed, the
140      ;      process loops back for the next message and continues in this
141      ;      fashion forever.
142      ;
143      ;      At present, receive errors are considered to be the Slave's

```

```

144 ; problem since normal error recovery is allegedly handled by the ;
145 ; transporter firmware. Only error free messages are passed on, ;
146 ; the rest are ignored unless the error is the absence of a free ;
147 ; support process in which case a "NOT LOGGED IN" error is sent ;
148 ; by the receiver process to the offending slave. ;
149 ; ;
150 ; In order to prevent clobbering the transporter when it is busy ;
151 ; transmitting, the receiver must be synchronized with the ;
152 ; transmit server. In this implementation, this is handled by ;
153 ; an MX Queue. ;
154 ; ;
155 ; _____ ;
156 ; ;
157 ;
158 ; receiver server process descriptor (position dependent if RSP)
159 ServerxPD:
160 0002 0000 dw 0 ;link
161 0004 0040 db 0,RX$PRIORITY ;status,priority
162 0006 6400 dw $ + 94 ;stack pointer
163 0008 5365727665 db 'ServerX ' ;name
164 0010 00FF db 0,0ffh ;console, memseg
165 0012 ds 82 ;reserved for MP/M use and stack
166 ServerxSTKTOP:
167 0064 9800 dw InitRX ;startup PC for process
168
169 ; User queue control block array used by this module for message queues.
170 ; Each element is 3 words long and is one UQCB followed by its message.
171 0006 = UQCBLN equ 6 ;constant used to index array
172 0004 = XQCBMSG equ 4 ;subindex for message word
173
174 INUQCB: ;array name
175 0000 # ??xx set 0
176 rept NSLAVES
177 dw (inqcb$array + ??xx) ;;Q pointer, msg addr, message word
178 dw $+2
179 dw 0
180 ??xx set ??xx + INQCB$SIZE

```

CP/M RMAC ASSEM 1.1 #005 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

181 endm
182 0066+AC04 DW (INQCB$ARRAY + ??XX)
183 0068+6A00 DW $+2
184 006A+0000 DW 0
185 006C+C604 DW (INQCB$ARRAY + ??XX)
186 006E+7000 DW $+2
187 0070+0000 DW 0
188
189 ; UQCB used by ServerX to get free buffers from Q
190 0072 1E057600 gbuf$uqcb: dw buffQCB,newbuff
191 0076 0000 newbuff: dw 0 ;message is a free buffer ptr from pool
192
193 ; UQCB used by ServerX to get transporter from MX Q
194 0078 A8087C00 omnirx$uqcb: dw omniQ,rx$mx$msg
195 007C 0000 rx$mx$msg: dw 0
196
197 ; UQCB used by ServerX to send error messages to outQ
198 007E E0048200 err$out$uqcb: dw outQCB,err$out$msg ;pointer, msgadr
199 0082 err$out$msg: ds 2 ;used to send error messages
200
201 ; receiver transporter control block
202 0084 F0 rxtcb: db 0f0h ;post read command
203 0085 00 db 0 ;result hi (always 0)
204 rxrslt:
205 0086 0000 db 0,0 ;result middle and low (NOT 8080 order)
206 0088 A0 db OMNI$SOCKET ;transporter message socket code
207 0089 00 db 0 ;data pointer high (always 0)
208 008A 0000 db 0,0 ;data pointer middle, low
209 008C 01 db BUFFSIZE/256 ;data max length hi
210 008D 18 db BUFFSIZE and 255 ;data max length lo
211 008E 0000 db 0,0 ;ctrl lgth (0 for now), host (not used)
212
213 0090 0000000000rxrslt: db 0,0,0,0,0,0,0,0 ;result block for rx
214
215 ; _____ ;
216 ; ;
217 ; Receiver server process initialization entry point ;

```

```

218 ; (initializes all of module) ;
219 ;
220 0098 CDCD08 InitRX: call omni$init ;init hardware & get ID code from its switches
221 009B 32FB02 sta configtbl+1 ; store ID in config table as master ID
222 ;
223 009E 0E86 mvi c,MAKEQ ;create the free buffer Q
224 00A0 111E05 lxi d,buffQCB
225 00A3 CDA408 call bdos
226 ;
227 00A6 11AC04 lxi d,inqcb$array
228 00A9 0E02 mvi c,NSLAVES ;create input Qs (1/slave supported)
229 make$inQs:
230 00AB D5 push d
231 00AC C5 push b
232 00AD 0E86 mvi c,MAKEQ
233 00AF CDA408 call bdos
234 00B2 C1 pop b

```

CP/M RMAC ASSEM 1.1 #006 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

235 00B3 D1 pop d
236 00B4 211A00 lxi h,INQCB$SIZE
237 00B7 19 dad d
238 00B8 EB xchg EB
239 00B9 0D dcr c
240 00BA C2AB00 jnz make$inQs
241 ;
242 00BD 11E004 lxi d,outQCB ;create the output Queue (only 1)
243 00C0 0E86 mvi c,MAKEQ
244 00C2 CDA408 call bdos
245 ;
246 00C5 11B901 lxi d,ServetxPD ;create the network output process
247 00C8 0E90 mvi c,CREATEP
248 00CA CDA408 call bdos
249 ;
250 00CD 0E9A mvi c,SYDATAD ;get system data page address
251 00CF CDA408 call bdos
252 00D2 110900 lxi d,9
253 00D5 19 dad d ;install config table address at sysdat(9)
254 00D6 11FA02 lxi d,configtbl
255 00D9 73 mov m,e
256 00DA 23 inx h
257 00DB 72 mov m,d
258 ;
259 00DC 219000 lxi h,rxrslt ;initialize transporter command block result
260 00DF 55 mov d,l ;field to point to receive result block
261 00E0 5C mov e,h ; (done at run time because of reversed byte
262 00E1 EB xchg EB ; order used by CORVUS.)
263 00E2 228600 shld rxrsltp
264
265
266 ; Receiver server process loop head
267 00E5 0E89 RXloop: mvi c,READQ
268 00E7 117200 lxi d,gbuf$uqcb
269 00EA CDA408 call bdos ;get a free message buffer from Q
270 ;
271 RXretry:
272 00ED 2A7600 lhld newbuff
273 00F0 5C mov e,h
274 00F1 55 mov d,l
275 00F2 EB xchg EB ;swap bytes for CORVUS command block
276 00F3 228A00 shld rxtcb+6 ;put buffer address pointer in rx tcb
277 ;
278 00F6 117800 lxi d,omnirx$uqcb ;read MX message from OMNINET HARDWARE MX Q
279 00F9 0E89 mvi c,READQ
280 00FB CDA408 call bdos
281 ;
282 00FE 018400 lxi b,rxtcb ;send TCB pointer to hardware
283 0101 CDF508 call omni$strobe
284 ;
285 0104 F5 push psw ;return MX message
286 0105 117800 lxi d,omnirx$uqcb
287 0108 0E8B mvi c,WRITEQ
288 010A CDA408 call bdos

```

CP/M RMAC ASSEM 1.1 #007 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82


```

289 010D F1          pop    psw          ;restore return code from omni$strobe routine
290                ;
291 010E DAED00      jc     RXretry        ;no choice except to retry if not accepted
292                ;
293 0111 019000      lxi    b,rxrslt      ;wait for a completion from hardware
294 0114 CD2309      call   wfrxdone      ;
295 0117 E680        ani    80h           ;if error on message, re-post buffer
296 0119 C2ED00      jnz   RXretry
297                ;
298                ; buffer contains a valid message at this point, so process it
299 011C 2A7600      lhld  newbuff       ;get FMT to A
300 011F 7E          mov    a,m
301 0120 23          inx   h              ;get SID to C
302 0121 23          inx   h
303 0122 4E          mov    c,m
304                ;
305 0123 E6FE        ani    0feh         ;look for login/logoff messages
306 0125 C24601      jnz   RXL2          ;message type 0 or 1?
307 0128 23          inx   h              ;yes, check FNC
308 0129 7E          mov    a,m
309 012A FE40        cpi    40h          ;login?
310 012C C23801      jnz   RXL1          ;not login, go on
311 012F CDA301      call  logiton       ;ELSE try to find a free SLVSP in table
312 0132 C26C01      jnz   RXL3          ;found one (or already logged in), go on
313 0135 C34C01      jmp   RX$send$err   ;sorry,no free processes, go advise slave
314                ;
315 0138 FE41        RXL1: cpi    41h          ;logoff?
316 013A C24601      jnz   RXL2          ;not logoff, go on
317 013D CD9A01      call  logitoff      ;ELSE try to remove that slave from table
318 0140 C26C01      jnz   RXL3          ;if successful, go on
319 0143 C34C01      jmp   RX$send$err   ;otherwise go tell slave it wasn't logged in
320                ;
321 0146 CD8001      RXL2: call   get$slvsp  ;not login/logoff so get slvsp msg address
322 0149 C26C01      jnz   RXL3          ; for that slave if it is logged in and go
323                ; send message to its Q else fall through
324                ;
325                ; this code sends a "NOT LOGGED IN" error message back to requester
326                RX$send$err:
327 014C 2A7600      lhld  newbuff       ;build an error message in the same buffer
328 014F 228200      shld  err$out$msg   ;
329 0152 3601        mvi    m,1          ;FMT=1
330 0154 23          inx   h
331 0155 7E          mov    a,m          ;swap DID and SID
332 0156 23          inx   h
333 0157 46          mov    b,m
334 0158 77          mov    m,a
335 0159 2B          dcx   h
336 015A 70          mov    m,b
337 015B 23          inx   h              ;leave FNC field alone
338 015C 23          inx   h
339 015D 23          inx   h
340 015E 3601        mvi    m,1          ;SIZ=1
341 0160 23          inx   h
342 0161 36FF        mvi    m,0ffh      ;message = 0FFH (extended error flag)

```

CP/M RMAC ASSEM 1.1 #008 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

343 0163 23          inx   h
344 0164 360C        mvi    m,12         ;"NOT LOGGED IN" code
345 0166 117E00      lxi    d,err$out$uqcb ;post to network transmitter process
346 0169 C37801      jmp   rxl4          ;using common write Q code
347                ;
348                ; this code sends the message address to the appropriate SLVSP Q
349 016C 2A7600      RXL3: lhld  newbuff       ;DE--> msg field of correct UQCB here
350 016F EB          xchg                          ;put message ptr in UQCB message field
351 0170 73          mov    m,e
352 0171 23          inx   h
353 0172 72          mov    m,d
354 0173 11FBFF      lxi    d,-(XQCBMSG + 1);index back to UQCB base address
355 0176 19          dad   d
356 0177 EB          xchg
357                ;
358 0178 0E8B        rxl4: mvi    c,WRITEQ
359 017A CDA408      call  bdos          ;send it to Queue
360 017D C3E500      jmp   RXloop        ;go back and get another buffer and continue
361

```

```

362
363 ; routine dynamically maps physical slave number passed in C
364 ; to a slave support process and returns its INUQCB message buffer addr
365 ; in DE and A = 0 with flags set if no room or not found, else NZ
366 get$slvsp:
367 0180 79 mov a,c ;A= requester ID
368 0181 0602 mvi b,NSLAVES ;set up for table search
369 0183 21B301 lxi h,idtbl
370 find$match: ;search till match or table end
371 0186 BE cmp m
372 0187 C29101 jnz not$match ; goto not$match if not this one
373 018A 23 inx h ;else match found, get ptr to SLVSP message
374 018B 5E mov e,m
375 018C 23 inx h
376 018D 56 mov d,m ;its slvsp msg addr
377 018E 37 stc
378 018F 9F sbb a
379 0190 C9 ret ;and return TRUE in A to caller
380 not$match:
381 0191 23 inx h ;no match, skip to next entry
382 0192 23 inx h
383 0193 23 inx h
384 0194 05 dcr b ;any more entries?
385 0195 C28601 jnz find$match ;loop back until all searched
386 0198 AF xra a ;else return failure (A=00)
387 0199 C9 ret
388
389
390 ; removes entry (C=SID) from map table (but still returns msg ptr)
391 logitoff:
392 019A CD8001 call get$slvsp
393 019D C8 rz ;not in table, just exit
394 019E 2B dcx h ;else mark entry as free and then exit
395 019F 2B dcx h
396 01A0 36FF mvi m,0ffh

```

CP/M RMAC ASSEM 1.1 #009 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

397 01A2 C9 ret
398
399 ; installs entry (C=SID) in first free entry of map table and returns
400 ; msg address. RETURNS A=0 if no space, else non-zero.
401 logiton:
402 01A3 CD8001 call get$slvsp ;see if already in table
403 01A6 C0 rnz ;if so, just use old entry
404 01A7 C5 push b ;else look for a free entry (CODE=FF)
405 01A8 0EFF mvi c,0ffh
406 01AA CD8001 call get$slvsp
407 01AD C1 pop b
408 01AE C8 rz ;no free entries, exit
409 01AF 2B dcx h ;else enter SID in table and return success
410 01B0 2B dcx h
411 01B1 71 mov m,c
412 01B2 C9 ret ;PSW is still correct from search
413
414 ; Slave mapping table has one entry per SLVSP. First byte = SID
415 ; of the requester currently using SLVSP (0ffh if none). Next word is
416 ; the address of the message field of that SLVSP's input UQCB.
417 idtbl:
418 0000 # ??xx set 0
419 rept NSLAVES
420 db 0ffh
421 dw (INUQCB + XQCBMSG + ??xx)
422 ??xx set ??xx + UQCBLEN
423 endm
424 01B3+FF DB 0FFH
425 01B4+6A00 DW (INUQCB + XQCBMSG + ??XX)
426 01B6+FF DB 0FFH
427 01B7+7000 DW (INUQCB + XQCBMSG + ??XX)
428
429 page

```

CP/M RMAC ASSEM 1.1 #010 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

430
431 ;_____

```

```

432 ;
433 ;
434 ; This is the network transmitter server process module. ;
435 ; NOTE THAT THE OMNINET TRANSPORTER MUST NOT BE DISTURBED ONCE ;
436 ; A TRANSMIT HAS BEEN POSTED UNTIL IT RETURNS A COMPLETION. ;
437 ; An MX Queue is used in this version to protect the transporter ;
438 ; from other processes. ;
439 ;
440 ; This process reads a message from the SLVSP output Q and when ;
441 ; awakened by one posts that buffer for transmission via the ;
442 ; transporter to the requester. This process then waits until ;
443 ; the transporter reports a completion as determined by the ;
444 ; wf$txdone routine. The buffer pointer from that message is ;
445 ; then sent back to the FreeBuff Q and the process loops back for ;
446 ; another message from the SLVSP output Q. Transmitter errors ;
447 ; are considered the Transporter's problem and are ignored here. ;
448 ;
449 ;
450 ; Transmitter server process descriptor
451 ServetxPD:
452 01B9 0000 dw 0 ;link
453 01BB 003F db 0,TX$PRIORITY ;status,priority
454 01BD 1B02 dw $ + 94 ;stack pointer
455 01BF 5365727665 db 'ServeTX ' ;name
456 01C7 00FF db 0,0ffh ;console, memseg
457 01C9 ds 82 ;reserved for MP/M use and as stack
458 021B 4302 dw InitTX ;stack top has startup PC
459
460 ; There is only one output queue (SLVSP --> NTRWKIF)
461 OUTUQCB:
462 021D E0042102 UQCBNtrwkQ00: dw outQCB,outQMSG ;pointer, msgadr
463 0221 outQMSG: ds 2 ;used to receive msg pointer from SLVSP
464
465 ; used by ServeTX to return them to Q when done (used at init also)
466 0223 1E052702 pbuf$uqcb: dw buffQCB,olddbuff
467 0227 0000 oldbuff: dw 0 ;msg is a freed buff ptr back to pool
468
469 ; UQCB used by ServeTX to get transporter from MX Q
470 0229 A8082D02 omnitx$uqcb: dw omniQ,tx$mx$msg
471 022D 0000 tx$mx$msg: dw 0
472
473 ; transmitter transporter control block
474 022F 40 txtcb: db 40h ;command
475 0230 00 db 0 ;result hi
476 txrslt:
477 0231 0000 db 0,0 ;result middle and low
478 0233 A0 db OMNI$SOCKET ;transporter message socket code
479 0234 000000 db 0,0,0 ;data ptr (MSB,SB,LSB)
480 0237 0000 db 0,0 ;length (MSB,LSB)
481 0239 00 db 0 ;control length
482 023A 00 db 0 ;dest host
483

```

CP/M RMAC ASSEM 1.1 #011 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

484 023B 0000000000txrslt: db 0,0,0,0,0,0,0,0 ;result block for tx
485
486 ;
487 ;
488 ; ServeTX initialization entry point ;
489 ;
490 InitTX:
491 0243 215C05 lxi h,msgbuffs ;preload the Free buffer Q with buffer ptrs
492 0246 0E03 mvi c,NMSG$BUFFS ;from start of buffer space
493
494 0248 222702 shld oldbuff
495 024B E5 push h
496 024C C5 push b
497 024D 0E8B mvi c,WRITEQ
498 024F 112302 lxi d,pbuf$uqcb
499 0252 CDA408 call bdos
500 0255 C1 pop b
501 0256 E1 pop h
502 0257 111801 lxi d,BUFFSIZE
503 025A 19 dad d
504 025B 0D dcr c
505 025C C24802 jnz freeloop

```

```

506 ;
507 025F 213B02 ; lxi h,txrslt ;initialize TX Transporter Command Block
508 0262 5C mov e,h ;to point to TX Result Block
509 0263 55 mov d,l
510 0264 EB xchg
511 0265 223102 shld txrsltp
512
513 ; ServeTX process loop
514 TXloop:
515 0268 0E89 mvi c,READQ ;wait for a message in network output Q
516 026A 111D02 lxi d,outuqcb
517 026D CDA408 call bdos
518 ;
519 0270 2A2102 lhld outQMSG
520 0273 5C mov e,h
521 0274 55 mov d,l ;put message buffer address in TX TCB
522 0275 EB xchg ;(NOTE, NOT (8080 byte order)
523 0276 223502 shld txtcb+6
524 ;
525 0279 13 inx d
526 027A 1A ldax d ;set transport layer destination addr=DID
527 027B 323A02 sta txtcb + 11
528 ;
529 027E 210300 lxi h,3
530 0281 19 dad d ;calculate physical message length
531 0282 6E mov l,m ;from SIZ field
532 0283 2600 mvi h,0
533 0285 110600 lxi d,6 ;put in TCB length field
534 0288 19 dad d
535 0289 55 mov d,l
536 028A 5C mov e,h
537 028B EB xchg

```

CP/M RMAC ASSEM 1.1 #012 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

538 028C 223702 shld txtcb+8
539 ;
540 028F 112902 lxi d,omnitx$uqcb ;get transporter hardware MX message
541 0292 0E89 mvi c,READQ
542 0294 CDA408 call bdos
543 ;
544 TXretry:
545 0297 012F02 lxi b,txtcb ;send TCB pointer to hardware
546 029A CDF508 call omni$strobe ;if can't, not much else to do but try again
547 029D DA9702 jc TXretry ; (ALTHOUGH THIS IS A FATAL HARDWARE ERROR)
548 ;
549 02A0 013B02 lxi b,txrslt ;wait for transmit completion
550 02A3 CD3409 call wftxdone ;ignore errors here as no recovery possible
551 ;
552 02A6 112902 lxi d,omnitx$uqcb
553 02A9 0E8B mvi c,WRITEQ
554 02AB CDA408 call bdos ;release MX msg
555 ;
556 02AE 2A2102 lhld outQMSG ;send the buffer back to FREEBUFF Q
557 02B1 222702 shld oldbuff
558 02B4 0E8B mvi c,WRITEQ
559 02B6 112302 lxi d,pbuf$uqcb
560 02B9 CDA408 call bdos
561 ;
562 02BC C36802 jmp txloop ;and go back and do it all with next msg
563
564
565 page

```

CP/M RMAC ASSEM 1.1 #013 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

566
567 02BF 4E5457524Bcnote: db 'NTWRKIF (c)1982 VANO ASSOCIATES, INC. - ALL RIGHTS RESERVED'
568 ;
569 ;
570 ;
571 ; GLOBAL Master Configuration Table and storage ;
572 ; (address must be installed on SysData page(9,10) at init.) ;
573 ;
574 ;
575 configtbl:

```

```

576 02FA 00      db      0          ;Master status byte
577 02FB 00      db      0          ;Master processor ID
578 02FC 02      db      NSLAVES    ;Maximum number of slaves supported
579 02FD 00      db      0          ;Number of logged in slaves
580 02FE 0000    dw      0          ;16 bit vector of logged in slaves
581 0300          ds      16          ;Slave processor ID array
582 0310 5041535357 db      'PASSWORD'    ;login password
583
584              ;          builds Server stacks and initializes them with PD storage pointers
585 0000 #        ??xx    set      0
586              rept     NSLAVES
587              ds      SRVR$STK$SIZ - 2
588              dw      srvr$pd$base + ??xx
589              ??xx    set ??xx + SRVR$PD$SIZ
590              endm
591 0318+        DS      SRVR$STK$SIZ - 2
592 03AC+4404    DW      SRVR$PD$BASE + ??XX
593 03AE+        DS      SRVR$STK$SIZ - 2
594 0442+7804    DW      SRVR$PD$BASE + ??XX
595
596              ;          allocates PD storage
597  srvr$pd$base:
598 0444          ds      NSLAVES * SRVR$PD$SIZ
599
600              ;
601              ;
602              ;
603              ;          INTERPROCESS QUEUES (both local and global) and COMMON data ;
604              ;
605              ;
606
607              ;          ServeRX --> SLVSP message queues (INPUT), 1/slave support proc.
608 001A =        INQCB$SIZE equ      26          ;constant used for index calculation
609              inqcb$array:          ;ARRAY BASE NAME
610              ;
611              ;          generate INQCBs as required
612 0030 #        ??xx    set      '0'
613              rept     NSLAVES
614              ds      2          ;;link
615              db      4eh,74h,77h,72h ;;common name is NtwrkQI
616              db      6bh,51h,49h    ;;(macro can't do lower case)
617              db      ??xx          ;;slave ID
618              dw      2,1          ;;msglen, nmbmsgs
619              ds      12          ;;MP/M pointers and buffers

```

CP/M RMAC ASSEM 1.1 #014 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

620              ??xx    set ??xx + 1
621              if (??xx EQ ('9'+1))
622              ??xx    set ??xx + 7
623              endif
624              endm
625 04AC+        DS      2
626 04AE+4E747772 DB      4EH,74H,77H,72H
627 04B2+6B5149  DB      6BH,51H,49H
628 04B5+30      DB      ??XX
629 04B6+02000100 DW      2,1
630 04BA+        DS      12
631 04C6+        DS      2
632 04C8+4E747772 DB      4EH,74H,77H,72H
633 04CC+6B5149  DB      6BH,51H,49H
634 04CF+31      DB      ??XX
635 04D0+02000100 DW      2,1
636 04D4+        DS      12
637
638              ;          SLVSP --> NETWRKIF queue (OUTPUT)
639 04E0          outQCB: ds      2          ;link
640 04E2 4E7477726B db      'NtwrkQ00'    ;name
641 04EA 02001000  dw      2,16          ;msglen, nmbmsgs
642 04EE          ds      48          ;Used by MP/M
643
644              ;          free buffer list management queue
645  buffQCB:
646 051E          ds      2          ;link
647 0520 4672656542 db      'FreeBuff'    ;name
648 0528 02001000  dw      2,16          ;msglen, nmbmsgs
649 052C          ds      48          ;reserved for MP/M

```

```

650
651
652         ;      global message buffer pool
653 005C      msgbufs:      ds          MSG$BUFFS * BUFFSIZE
654
655         ;      Utility Procedure to allow indirect BDOS/XDOS access as needed by RSP
656 00A4 2A0000 bdos:      lhld      bdosadr
657 00A7 E9          pchl
658
659         page

```

CP/M RMAC ASSEM 1.1 #015 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

660
661         ; _____
662         ; _____
663         ; _____ ;
664         ;      low level omninet support routines
665         ; _____
666         ; _____
667
668         ;      Transporter mutual exclusion QUEUE
669 00A8      omniQ:      ds          2
670 00AA 4D586F6D6E db          'MXomniQ '
671 00B2 00000100      dw          0,1          ;msglen, nmsgs
672 00B6          ds          12          ;dqph,nqph,msgin,msgout,msgcnt, buff
673
674         ;      UQCB used by omni$init to load MX Q
675 00C2 A808C608 omni$init$uqcb: dw          omniQ,init$mx$msg
676 00C6 0000      init$mx$msg:      dw          0
677
678
679         ;      Initialization transporter control block
680 inittcb:
681 00C8 20          db          20h          ;command
682 00C9 00          db          0          ;result hi
683
684 00CA 0000      initrslt:      db          0,0          ;result middle and low
685
686         ;
687 00CC 00      initrslt:      db          0          ;result block for init
688
689
690         ;      initializes transporter hardware and return its network ID code in A
691 omni$init:
692 00CD 11A808      lxi          d,omniQ
693 00D0 0E86      mvi          c,MAKEQ
694 00D2 CDA408      call         bdos          ;create hardware MX Q
695 00D5 11C208      lxi          d,omni$init$uqcb ;send it one message
696 00D8 0E8B      mvi          c,WRITEQ
697 00DA CDA408      call         bdos
698         if INTERRUPT
699             call int$init          ;(optional) setup interrupt system
700         endif
701 00DD 21CC08      lxi          h,initrslt      ;install result block pointer in initialization
702 00E0 55          mov          d,l          ;TCB
703 00E1 5C          mov          e,h          ;NOTE: NOT 8080 order, MSB,LSB
704 00E2 EB          xchg
705 00E3 22CA08      shld         initrslt
706
707         ;
707 00E6 01C808      lxi          b,inittcb      ;post initialization command block to
708 00E9 CDF508      call         omnistrobe     ;hardware
709 00EC D8          rc          ;cy=1 means can't talk to hardware
710
711         ;
711 00ED 01CC08      lxi          b,initrslt      ;wait for a completion from operation
712 00F0 CD2309      call         omni$wfdone
713 00F3 B7          ora          a

```

CP/M RMAC ASSEM 1.1 #016 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

714 00F4 C9          ret          ;return ID/result code to caller with flags set
715
716
717         ;      sends the command block pointer in BC to transporter hardware
718 omni$strobe:
719 00F5 210200      lxi          h,2          ;first preset result code byte in

```

```

720 08F8 09      dad      b          ;result block TCB result field --> to 0ffh
721 08F9 7E      mov      a,m
722 08FA 23      inx      h
723 08FB 6E      mov      l,m
724 08FC 67      mov      h,a
725 08FD 36FF    mvi      m,0ffh
726
727 08FF AF      xra      a          ;send bits 23-16 of ptr to hardware (always 0)
728 0900 CD0A09  call    omni$st
729 0903 D8      rc          ;carry means can't talk to hardware
730
731 0904 78      mov      a,b          ;send bits 15-8 of ptr to hardware
732 0905 CD0A09  call    omni$st
733 0908 D8      rc
734
735 0909 79      mov      a,c          ;send bits 7-0 of ptr to hardware
736                          ;fall into omni$st to send last byte
737
738 ;          called by omni$strobe to send one byte to transporter when ready
739 ;          (waits a reasonable time for transporter to come ready and if
740 ;          it doesn't, returns with carry set; this is a fatal error) returns
741 ;          cy=0 if succeeds
742
743 090A F5      omni$st:  push    psw          ;save data for now
744 090B 1150C3  lxi      d,50000     ;set timeout
745
746 090E DBF9    omni$st0: in      OMNI$STAT    ;see if transporter will accept byte
747 0910 E610    ani      OMNI$RDY
748 0912 CA1A09  jz      omni$st1     ;if busy, go decrement timeout and retry
749 0915 F1      pop      psw          ;else output the byte and return with CY=0
750 0916 D3F8    out     OMNI$DATA
751 0918 B7      ora      a
752 0919 C9      ret
753
754 091A 1B      omni$st1: dcx      d          ;loop back if not timeout yet
755 091B 7B      mov      a,e
756 091C B2      ora      d
757 091D C20E09  jnz     omni$st0
758 0920 F1      pop      psw
759 0921 37      stc
760 0922 C9      ret          ;else return CY=1 as error flag
761
762
763 ;          routine waits for a completion to occur on the result block
764 ;          pointed to by BC. This routine is used by the initialization
765 ;          and receiver processes. If there is no interrupt hardware in
766 ;          the system, ONLY ONE MESSAGE CAN BE RECEIVED PER CLOCK TICK of
767 ;          the system clock. This will considerably reduce server throughput

```

CP/M RMAC ASSEM 1.1

#017 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

768 ;          in most systems.
769 omni$wfdone:
770 wfrxdone:
771 0923 0A      ldax     b          ;all completion codes are < 0f0h
772 0924 FEF0    cpi      0f0h       ;see if already done before suspending caller
773 0926 D8      rc          ;yes, return immediately
774 ;          else suspend caller until a completion occurs
775 0927 C5      push    b
776 if INTERRUPT
777 lxi      d,OMNI$FLAG ;wait for ISR to set flag
778 mvi      c,FLAGWAITF
779 call    bdos
780 else
781 0928 110100  lxi      d,1        ;if no ISR, poll result block once/tick
782 092B 0E8D    mvi      c,DELAY
783 092D CDA408  call    bdos
784 endif
785 0930 C1      pop      b
786 0931 C32309  jmp     omni$wfdone
787
788 ;          As above but instead polls continually to give transmitter priority
789 ;          since transmitter usually unloads messages in less time than MP/M
790 ;          dispatch overhead, it is not worth suspending it.
791 ;          A timeout routine is included to avoid locking up system if hardware
792 ;          fails so diagnosing the problem is possible with RDT.
793 wftxdone:

```

```

794 0934 1150C3      lxi    d,50000      ;initialize hardware fail timeout
795 0937 0A          wftxd0: ldax   b      ;done yet?
796 0938 FEF0        cpi    0f0h
797 093A 3F          cmc
798 093B D0          rnc      ;set up carry properly in case of return
799 093C 1B          wftxd1: dcx   d      ;yes, return to caller with result in A, CY=0
800 093D 7B          mov    a,e          ;if not timeout, loop back
801 093E B2          ora    d
802 093F C23709      jnz   wftxd0
803 0942 37          stc
804 0943 C9          ret      ;else return to caller with CY=1 as error flag
805
806                page

```

CP/M RMAC ASSEM 1.1 #018 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

807
808                if INTERRUPT
809                ;
810                ;       Since the CORVUS "ENGINEERING" transporter has no interrupt hardware
811                ;       associated with it, the details of the interrupt initialization and
812                ;       service routines will vary from system to system. The skeleton of
813                ;       our code is provided here as a guide to understanding what is needed.
814                ;
815                ;       Routine initializes interrupt hardware and attaches ISR to XIOS
816                ;       at run-time (in somewhat bizarre fashion.) It would be better
817                ;       to make your ISR a permanent part of your XIOS since if not
818                ;       used it does no harm to the system.
819                int$init:
820                    di
821                    mvi    a,(jmp)      ;build jump in vector
822                    sta    (INT$VCTR)
823                    lxi    h,omni$isr
824                    shld   (INT$VCTR + 1) ;install new isr
825                    out    OMNI$ACK     ;clear interrupt latch
826                    mvi    a,OMNI$ENABLE ;unmask transporter interrupt
827                    out    OMNI$MASK
828                ; this code does an extremely Klugey run-time linkage to needed XIOS routines
829                    lhld   1            ;find CBOOT in MPM-II BIOS simulation table
830                    mvi    l,1
831                    mov    e,m
832                    inx    h
833                    mov    d,m
834                    push   d            ;save to find exit$reg.
835                ;
836                    xchg
837                    inx    h            ;need to go one more level to find real entry
838                    mov    e,m
839                    inx    h
840                    mov    d,m          ;this is address of real CBOOT entry in XIOS
841                ;
842                    lxi    h,9          ;calculate PDISP entry from CBOOT address
843                    dad    d
844                    shld   pdisp        ;and save it in local vector
845                ;
846                    lxi    d,3          ;XDOS address is 3 bytes above PDISP
847                    dad    d
848                    shld   xd$adr       ;save it in a local vector
849                ;
850                    pop    h            ;get XIOS branch table address back
851                    mvi    l,40h        ;calculate address of EXIT$REGION entry
852                    mov    e,m
853                    inx    h
854                    mov    d,m
855                    xchg
856                    shld   exit$region  ;save it for later use in pre-empt routine
857                    ei
858                    ret
859                ;
860                ;       omninet isr sets the appropriate XDOS flag and causes a dispatch

```

CP/M RMAC ASSEM 1.1 #019 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

```

861                omni$isr:
862                    shld   svhl
863                    pop    h

```



```

864      push    psw          ;save PSW and HL
865      shld   svret        ;save return address
866      lxi    h,0          ;swap stacks
867      dad    sp
868      shld   svstk
869      lxi    sp,isr$stk
870      push   d            ;save the other registers on new stack
871      push   b
872      ;
873      out    OMNI$ACK      ;clear interrupt latch
874      ;
875      lhld   exit$region   ; do a PRE-EMPT by patching a RET into table
876      mov    a,m          ; (Very KLUGEY but there's no other way.)
877      push   psw          ; save what was in XIOS branch table entry
878      push   h            ; and put a RET there to prevent XDOS from
879      mvi    m,(RET)      ; re-enabling interrupts
880      ;
881      mvi    c,FLAGSETF   ;call XDOS to set isr flag
882      mvi    e,OMNI$FLAG
883      call   xdos
884      ;
885      pop    h
886      pop    psw
887      mov    m,a          ;restore XIOS table entry
888      ;
889      pop    b            ;pop interrupted registers
890      pop    d
891      lhld   svstk        ;restore interrupted stack
892      sphl                      ;restore other regs. and exit
893      pop    psw
894      lhld   svret
895      push   h
896      lhld   svhl
897      db    (JMP)          ; via dispatcher
898      pdisp: dw    0        ;(link to dispatcher)
899
900      xdos:  db    (JMP)      ;special XDOS entry
901      xd$adr: dw    0        ;for ISR use
902
903      ;      ISR data areas
904      exit$region:
905      dw    0              ;address of XDOS critical region exit routine
906      ds    64            ;isr stack space
907      isr$stk:
908      svhl:  dw    0        ;temporary reg storage
909      svret: dw    0
910      svstk: dw    0        ;careful, make sure all of .RSP is reserved
911
912      endif ; of if INTERRUPT
913
914      0944      end

```

CP/M RMAC ASSEM 1.1 #020 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82

BDOS	08A4	110	113	225	233	244	248	251	269	280	288
		359	499	517	542	554	560	656#	694	697	779
		783									
BDOSADR	0000	117#	122#	656							
BUFFQCB	051E	190	224	466	645#						
BUFFSIZE	0118	62#	209	210	502	653					
CNOTE	02BF	567#									
CODESEG	0000	104#	118								
CONFIGTBL	02FA	221	254	575#							
CR	000D	100#									
CREATEP	0090	94#	247								
DEBUG	0000	54#	111								
DELAY	008D	92#	782								
DETACH	0093	96#	112								
DSPTCH	008E	93#									
ERROUTMSG	0082	198	199#	328							
ERROUTQCB	007E	198#	345								
FINDMATCH	0186	370#	385								
FLAGSETF	0085	88#	881								
FLAGWAITF	0084	87#	778								
FREELoop	0248	493#	505								
GBUFUQCB	0072	190#	268								
GETSLVSP	0180	321	366#	392	402	406					

TXMMSG	022D	470	471#						
TXPRIORITY	003F	65#	453						
TXRETRY	0297	544#	547						
TXRSLT	023B	484#	507	549					
TXRSLTP	0231	476#	511						
TXTCB	022F	474#	523	527	538	545			
UQCBLEN	0006	171#	422						
UQCBNTWRKQ00	021D	462#							
WFRXDONE	0923	294	770#						
WFTXD0	0937	795#	802						
WFTXD1	093C	799#							
CP/M RMAC ASSEM 1.1 #022 SAMPLE SERVER NETWORK I/F FOR CORVUS OMNINET 20-OCT-82									
WFTXDONE	0934	550	793#						
WRITEQ	008B	91#	287	358	497	553	558	696	
XQCBMSG	0004	172#	354	421	425	427			
YES	FFFF	50#	51	55					

Listing G-2. Sample Server Network I/O for Corvus OMNINET